# AI Problem Solving by Searching & Robot Path Planning

Lecture by Margrit Betke

Reading: Russell and Norvig, Chapter 3
Winston



# What you need to do to specify an agent-based AI Problem:

- Initial state that the agent starts in
- Actions available to agent
- Transition model & state space: Path through state space = sequence of states = sequence of actions
- Goal test
- Path cost (e.g. sum of step costs)



## "Al Toy Problems" useful for learning concepts

• 8 Puzzle: Start state:

Goal state:

12 345 678

8 Queens Problem: 8x8 chess board

Place queens so that none attacks any other queen.

"Attack state" = 2 queens are on the same row, column or diagonal.

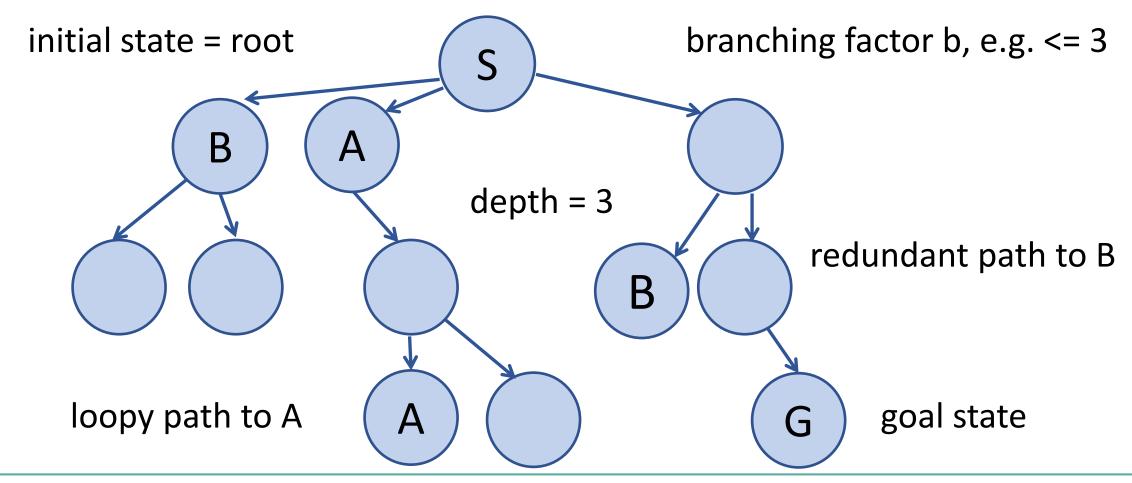


#### Problem Solving by Searching: "Real-World Problems"

- Route finding problems
- Touring problems, e.g., Traveling Salesperson (efficient path for visiting every city once)
- VLSI layout
- Robot navigation
- Drone navigation
- Protein design
- Cancer detection



# Solution: Sequence of Actions = Path through a Search Tree





#### Evaluation of Search Algorithm Performance

- Completeness: Is it guaranteed to find a solution?
- Optimality:
  - ➤ Shortest path?
  - **≻**Lowest cost?
- Time Complexity
- Space Complexity



### Path-based Search Algorithms

Task: Find shortest path through a graph

Applications: Games, robot path planning

Lots of algorithms!



• Exhaustive search: Explore all paths



#### **Exhaustive Search**

- Winston calls this strategy the "British Museum Procedure"
- Find all paths and select the shortest

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• Search tree: root level 1 node
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2<sup>nd</sup> level b nodes

3<sup>rd</sup> level b\*b nodes

4<sup>th</sup> level b\*b\*b nodes

• • • •

dth level bd nodes

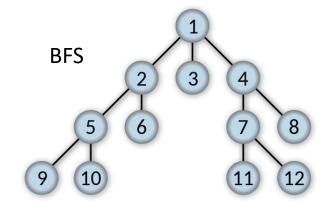
If b=10, d=10:  $10^{10} = 10$  billion paths Too many to test!

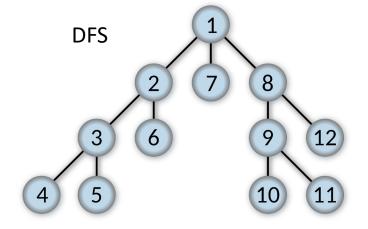


- Exhaustive search: Explore all paths
- Breadth-first search (BFS): Expand shallowest node
- Depth-first search (DFS): Expand deepest node

Cost per edge: 1

BFS & DFS typically covered in your previous classes. If not, please read Wikipedia pages







- Exhaustive search: Explore all paths
- Breadth-first search (BFS): Expand shallowest node
- Depth-first search (DFS): Expand deepest node
- Uniform cost search: Expand node with smallest cost
- Beam search: BFS but only keep limited number of best nodes
- Depth-limited search: Predetermined depth limit
- Iterative Deepening: Gradually increasing depth d = 0, 1, 2, ...
- Bidirectional search: Search from start S and goal G nodes, hoping to meet

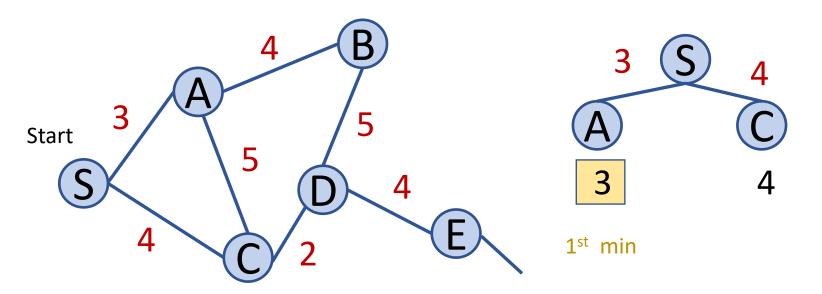


- Exhaustive search: Explore all paths
- Breadth-first search (BFS): Expand shallowest node
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- Uniform cost search: Expand node with smallest cost
   Used in AI when graphs are extremely large or infinite



#### Example for Uniform Cost Search

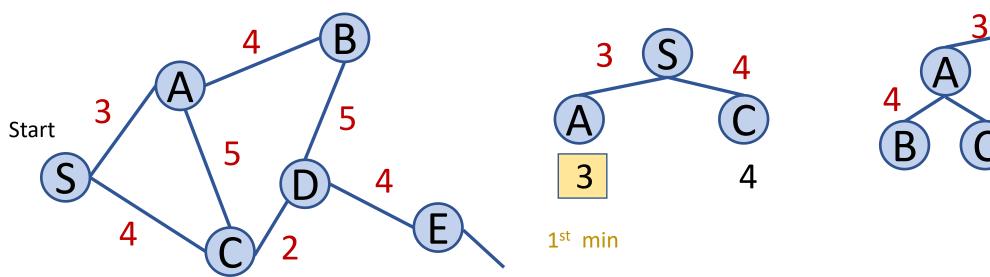
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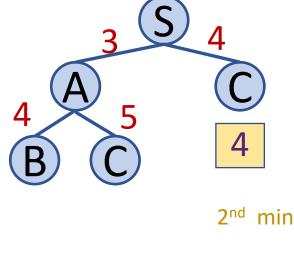




### Example for Uniform Cost Search

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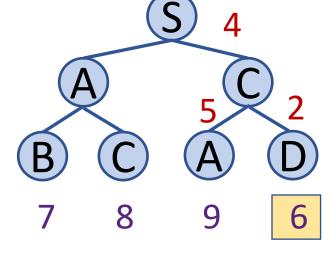




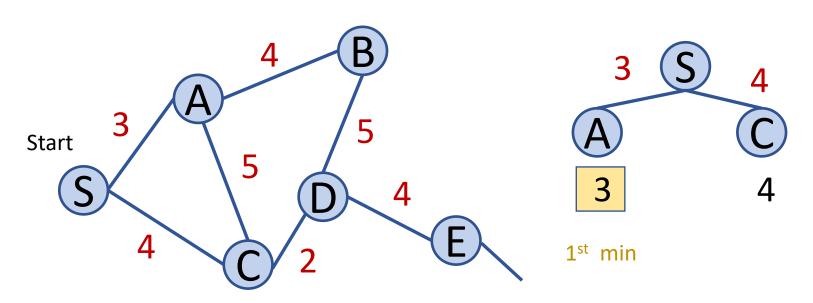


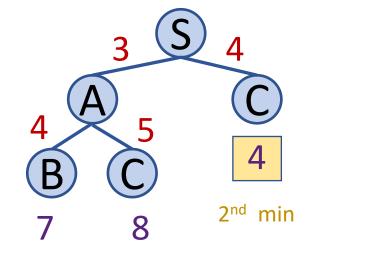
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- Uniform cost search: Expand node with smallest cost



3<sup>rd</sup> min





#### Beam Search

- Exhaustive search: Explore all paths
- Breadth-first search (BFS): Expand shallowest node
- Depth-first search (DFS): Expand deepest node
- Uniform cost search: Expand node with smallest cost
- Beam search: BFS but only keep limited number w of best nodes at each level, the beam width w

Same as BFS with w = infinite

The greater w is the fewer states are pruned

Useful in AI if BFS search tree is too large to fit in memory

Not guaranteed to find optimal solution



- Exhaustive search: Explore all paths
- Breadth-first search (BFS): Expand shallowest node
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- Uniform cost search: Expand node with smallest cost
- Beam search: BFS but only keep limited number of best nodes
- Depth-limited search: Predetermined depth limit



- Exhaustive search: Explore all paths
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- Uniform cost search: Expand node with smallest cost
- Beam search: BFS but only keep limited number of best nodes
- Depth-limited search: Predetermined depth limit
- Progressive Deepening (also called Iterative Deepening):

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Gradually increasing depth d = 0, 1, 2, ...
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We have seen this algorithm used for adversarial game playing.



- Exhaustive search: Explore all paths
- Breadth-first search (BFS): Expand shallowest node
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- Uniform cost search: Expand node with smallest cost (same as BFS with cost =1)
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- Bidirectional search: Search from start S and goal G nodes, hoping to meet
- Greedy search = branch & bound search
- Greedy search with pruning
- A\* = Greedy search with pruning and underestimates of remaining distance



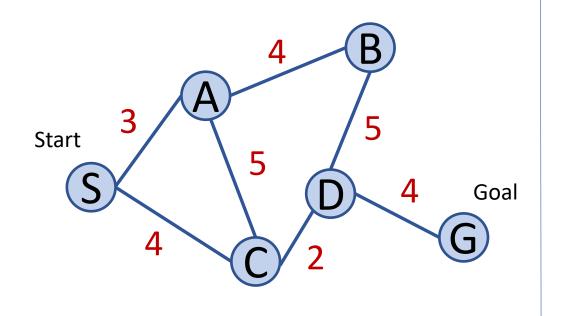
#### Greedy Search = Branch & Bound Search

Phase 1: Extend shortest partial path until goal is reached.

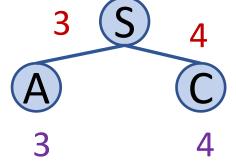
Reject loops.

Phase 2: Extend all partial paths until their length >= complete path to goal



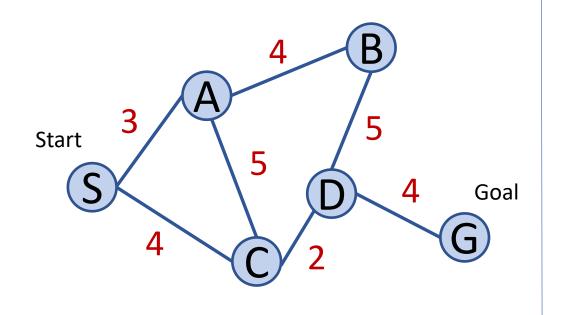


Phase 1:



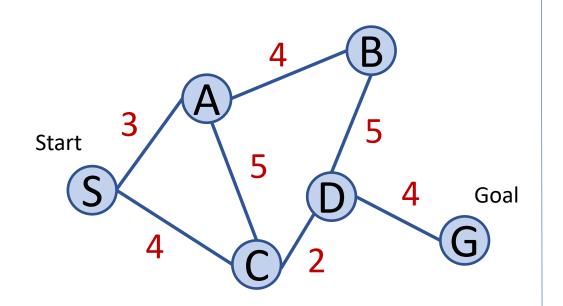
Phase 1: Extend shortest partial path until goal is reached. Reject loops.



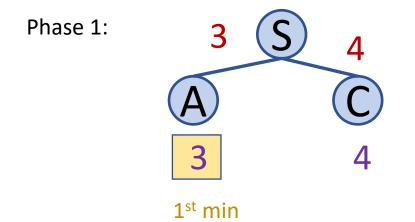


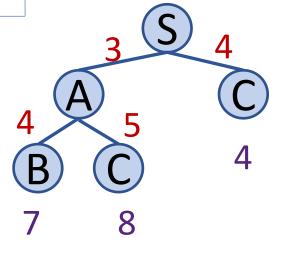
Phase 1: 3 S 4
A C
3 4
1st min

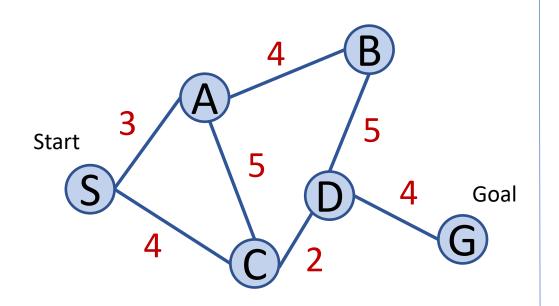
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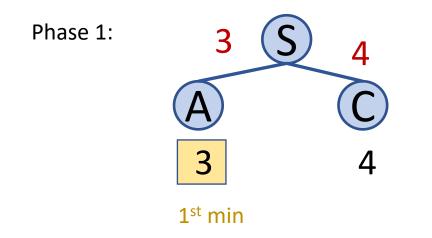
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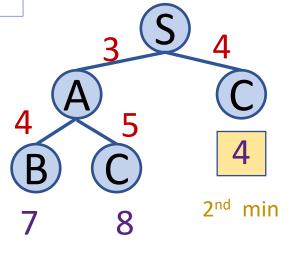


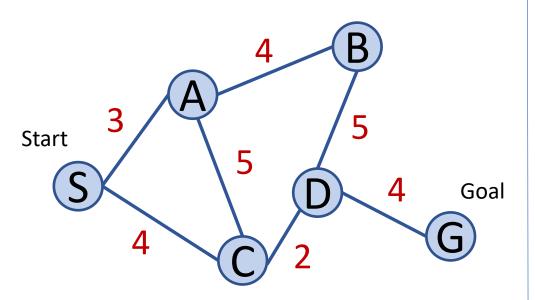




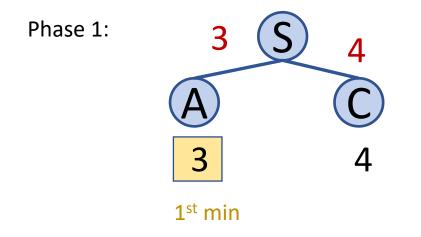
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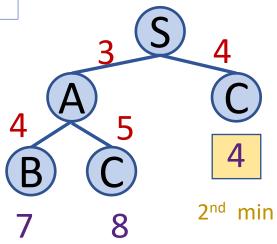


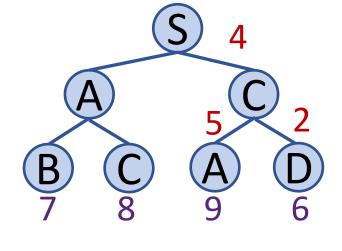


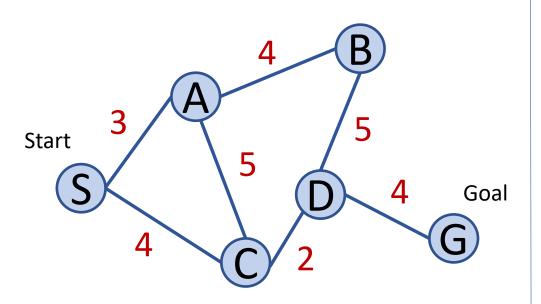


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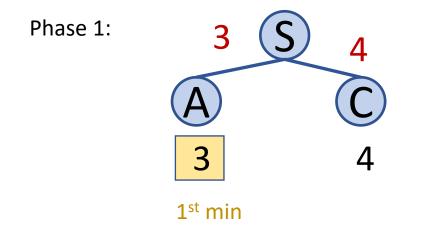


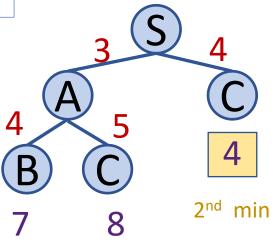


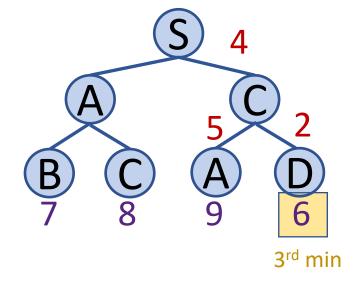


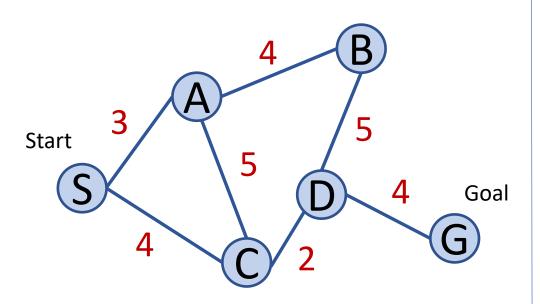


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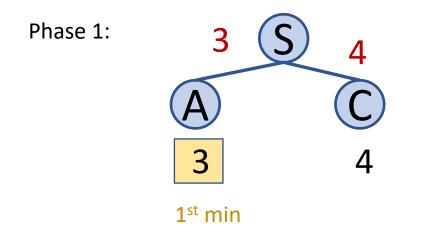


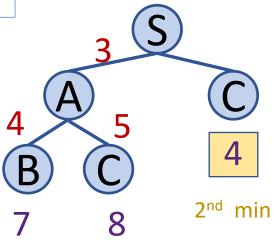


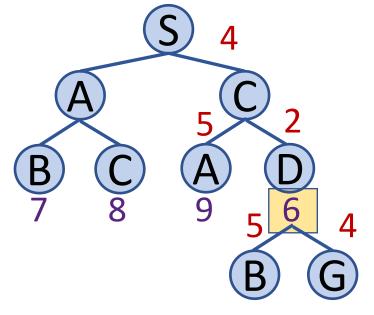




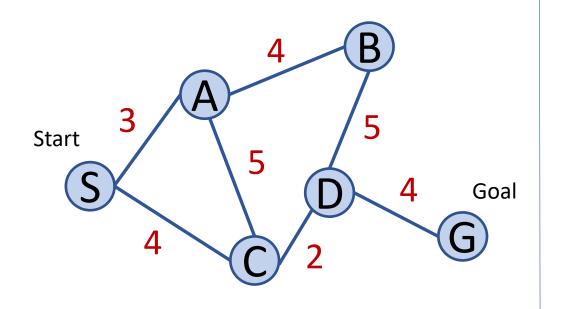
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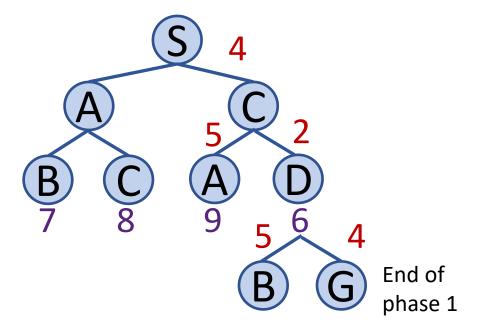




End of phase 1: G reached

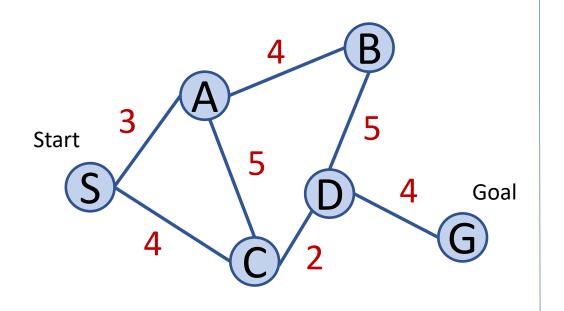


Phase 2:

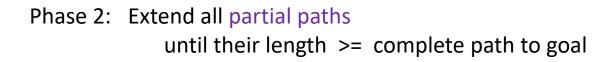


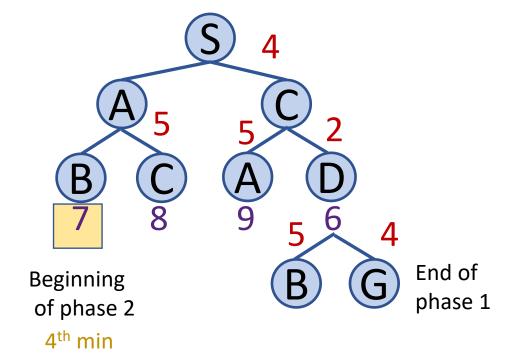
Phase 2: Extend all partial paths
until their length >= complete path to goal

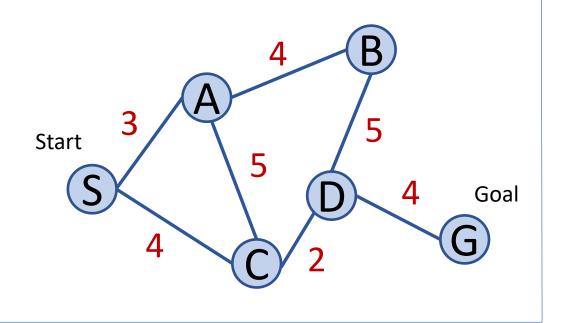




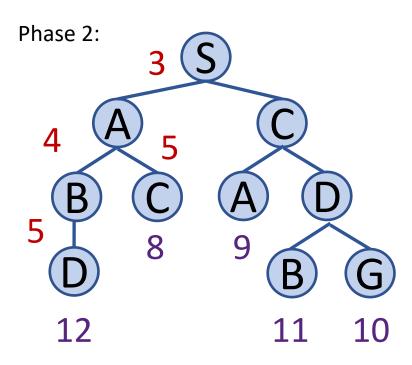
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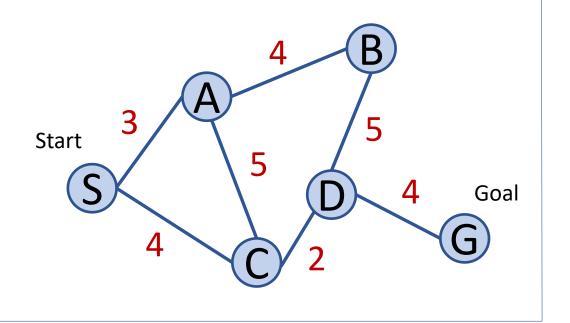




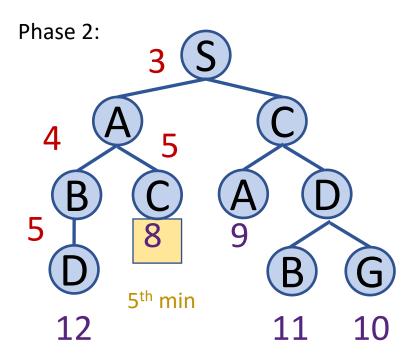


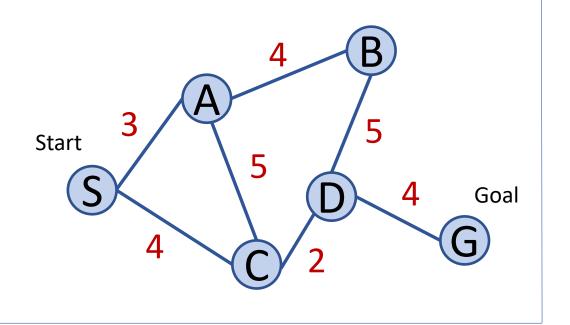
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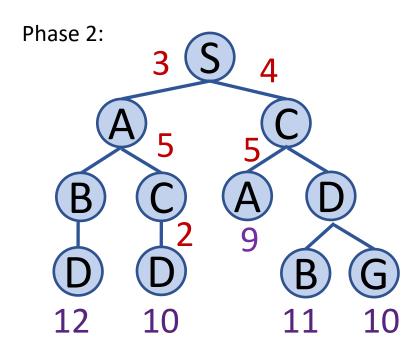


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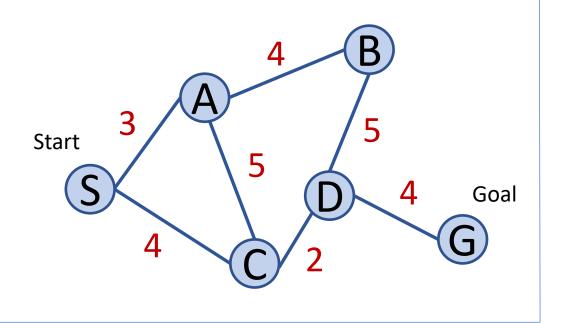




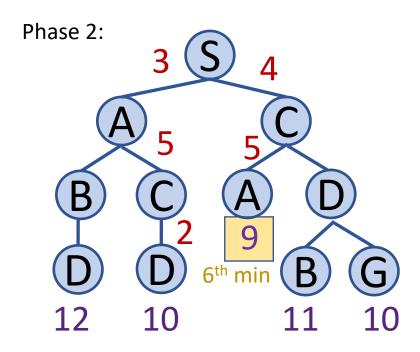
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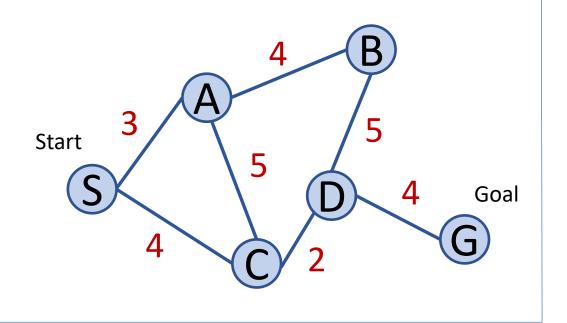




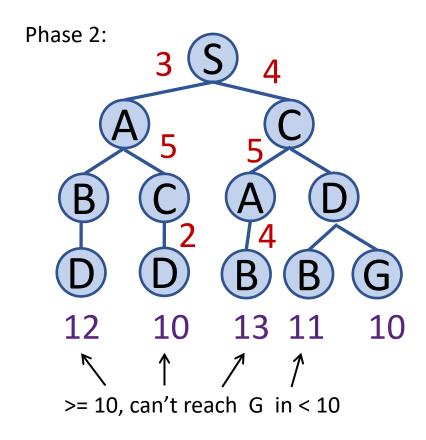
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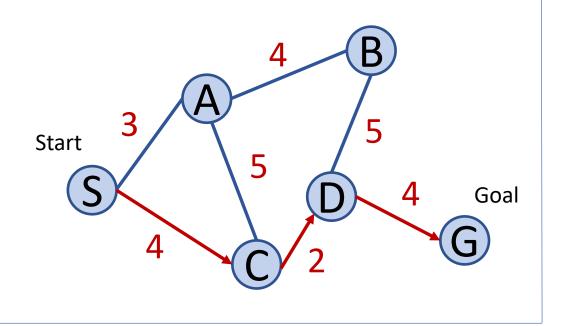




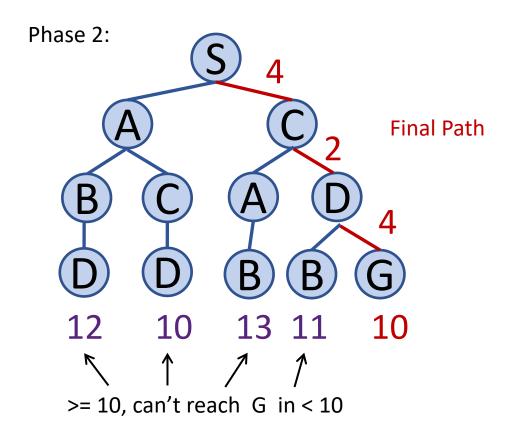


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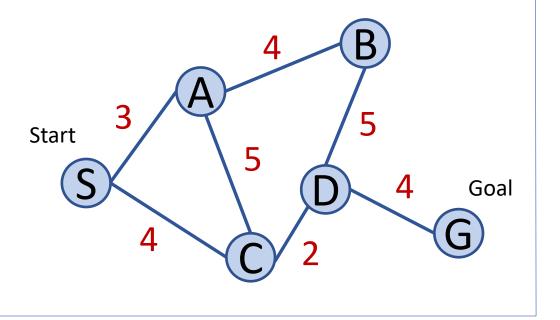
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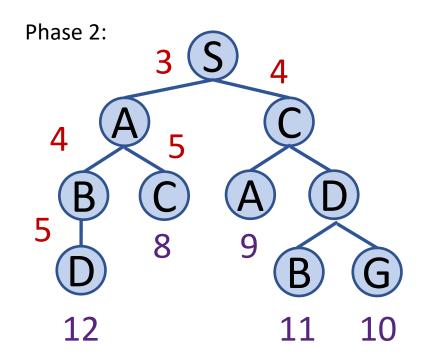


**Dynamic Programming Principle:** 

If two or more paths reach a common node, delete all paths except the minimum cost path.

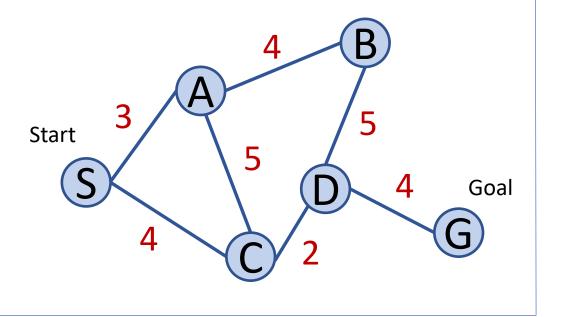


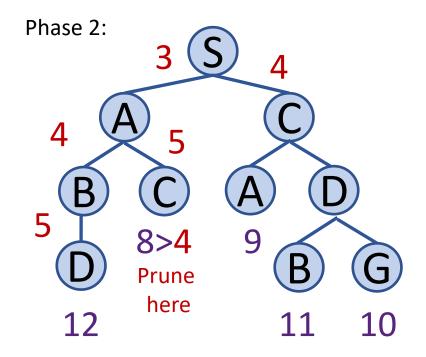




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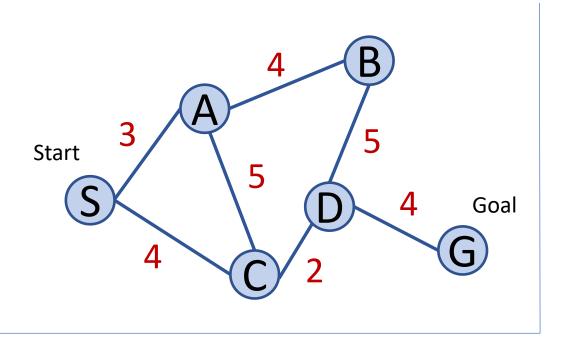




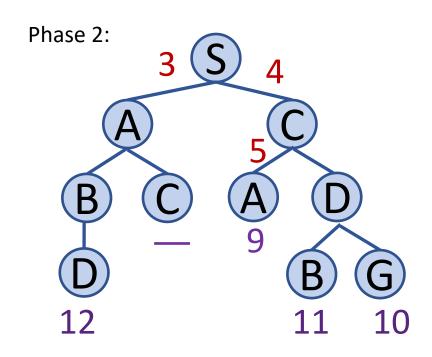


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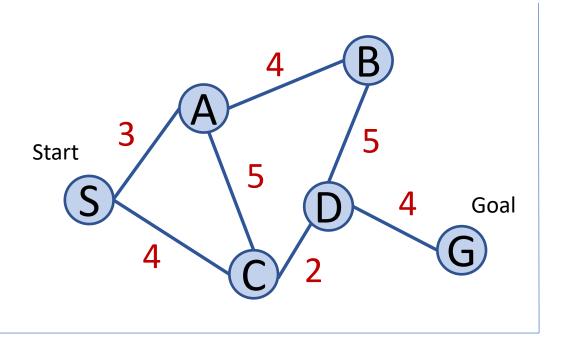




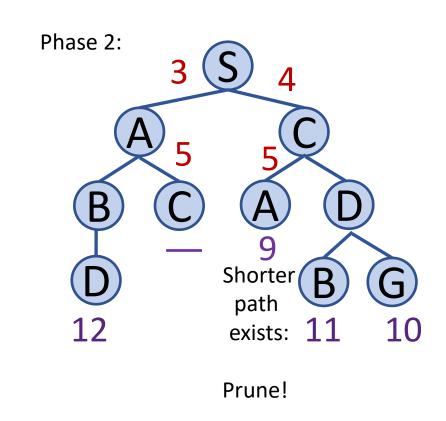
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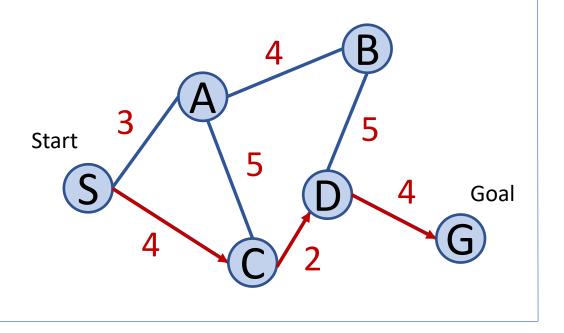




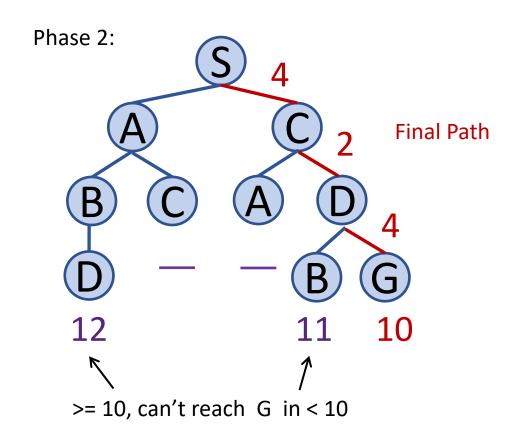
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Phase 2: Extend all partial paths
until their length >= complete path to goal



### Implementation of Greedy Search with Pruning

Data Structure: Queue

Elements of queue: Partial paths

Initialize: Place start node in queue

Until a path in queue reaches goal node or queue is empty:

Remove 1<sup>st</sup> queue element & extend path to its children

Reject loops

Add new paths to queue

Prune

Sort



### Connection to Dijkstra's Algorithm

#### Same as Greedy Search with Pruning except

Dijkstra's algorithm computes the all-pairs shortest paths while "Greedy Search with Pruning" computes the shortest path between a single start state and a single goal state.



# A\* Algorithm = Greedy Search with Pruning and Underestimates of Remaining Distance

- Remaining distance = e.g., straight-line distance on a highway map
- In each step:

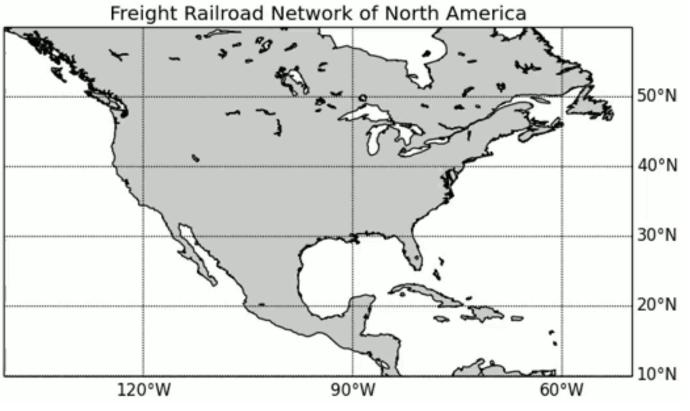
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Estimate of total path length = length of partial path + underestimate of remaining
```

"This path is at least this bad."

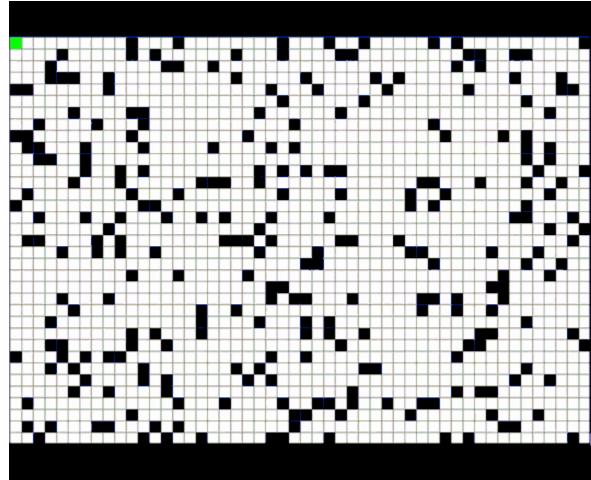


An animation of the A\* algorithm as it

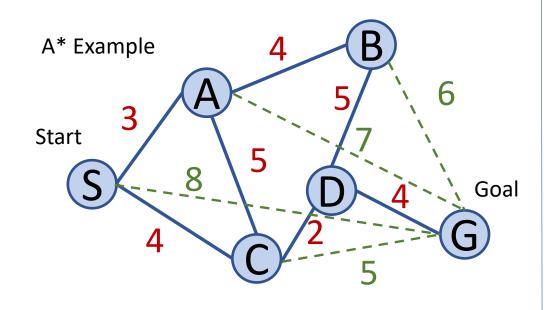
explores the North American freight train network to find the optimum path between Washington, D.C. and Los Angeles.



A\* pathfinding algorithm navigating around a randomly-generated maze







Phase 1: Extend shortest estimated partial path

(= length of partial path + underestimate of remaining)

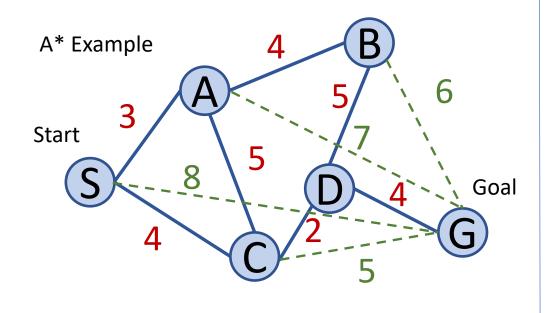
until goal is reached.

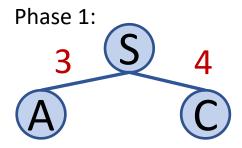
Reject loops.

Phase 2: Extend all estimated partial paths until their length >= complete path to goal



**A**\*





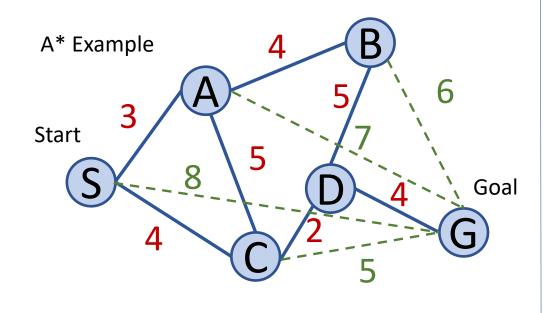
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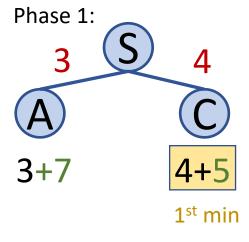
(= length of partial path + underestimate of remaining)

until goal is reached.

Reject loops.





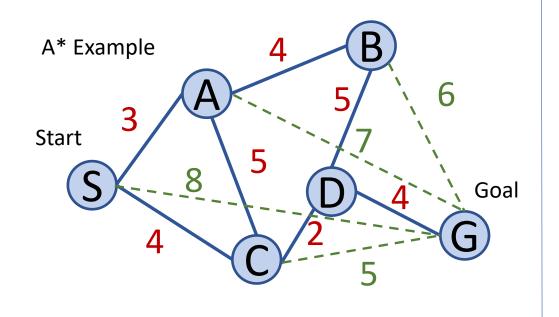


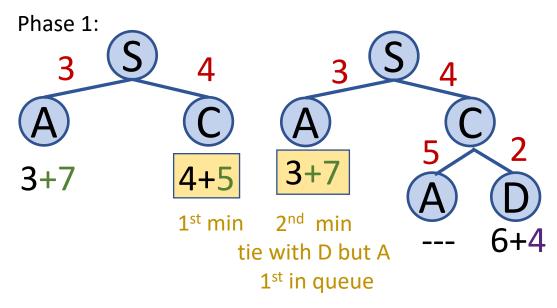
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Reject loops.





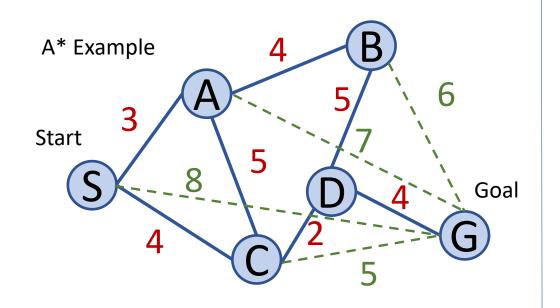
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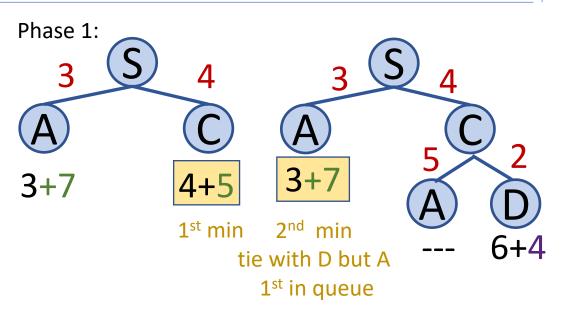
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Reject loops.





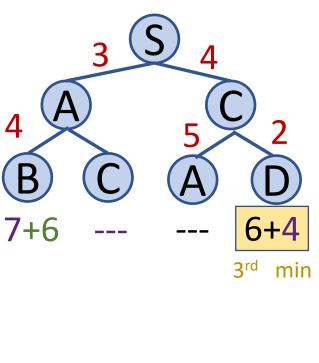


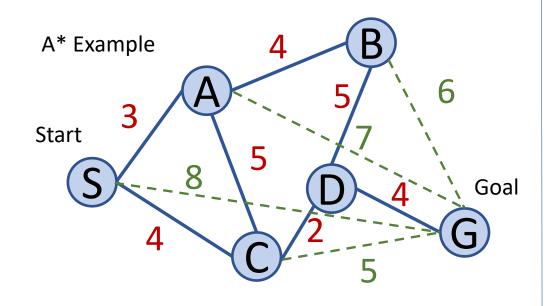
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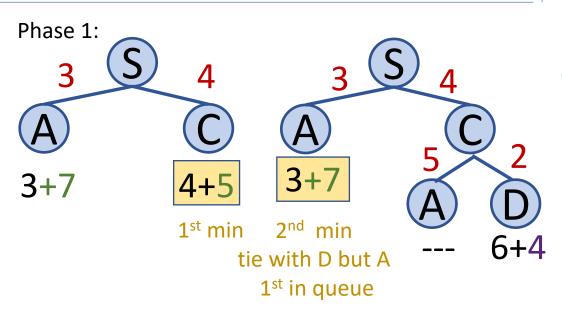
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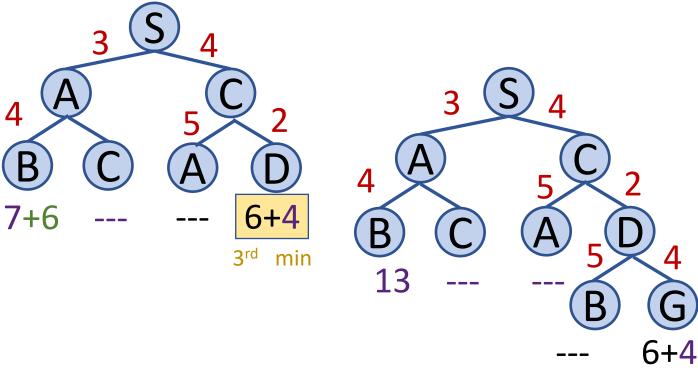


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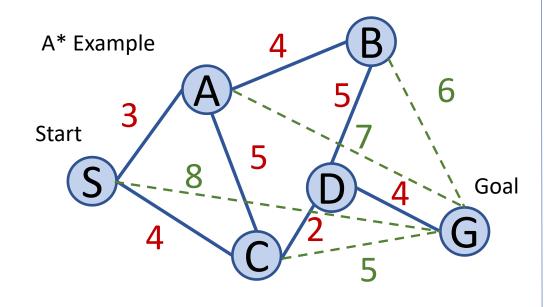
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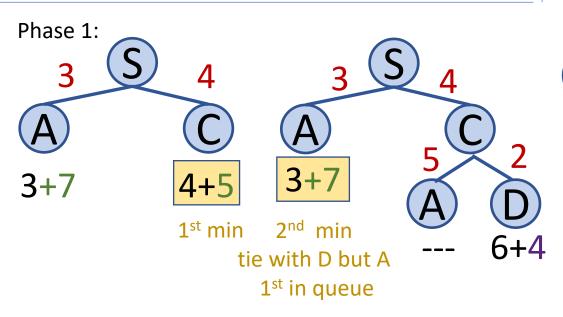
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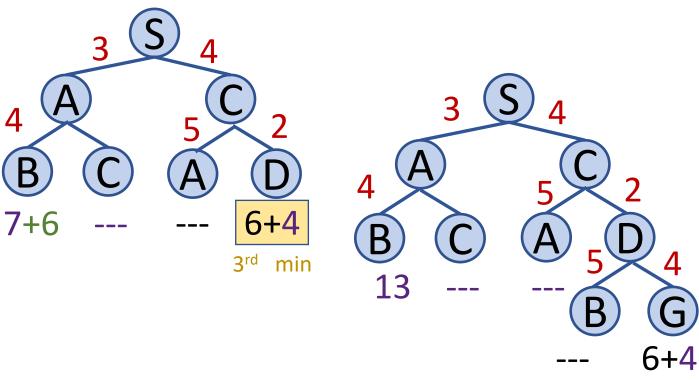


Phase 1: Extend shortest estimated partial path

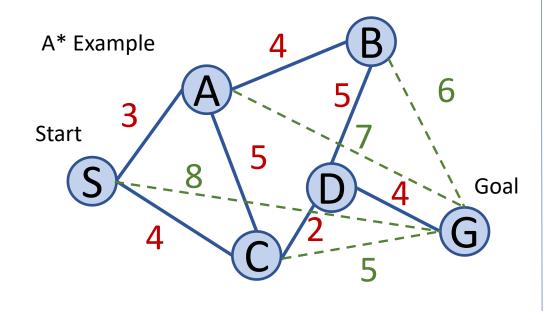
(= length of partial path + underestimate of remaining)

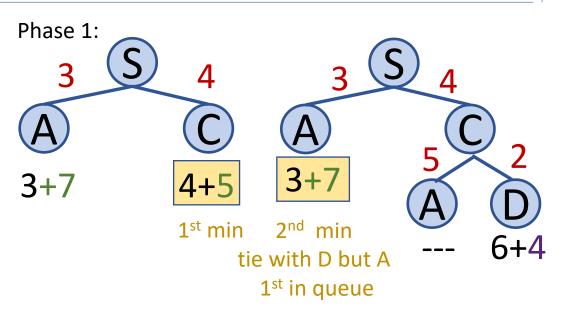
until goal is reached.

Reject loops.



End of phase 1: G reached





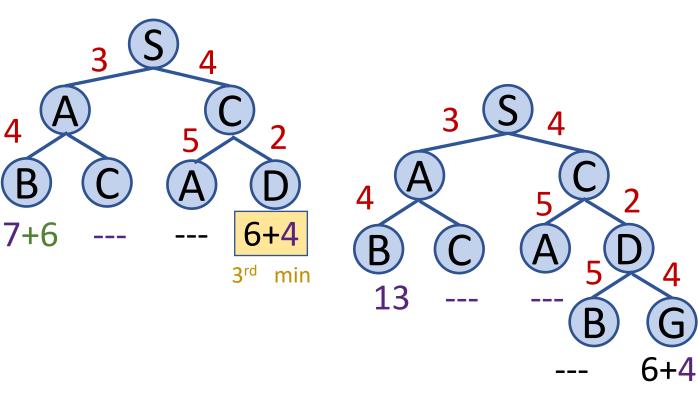
Phase 1: Extend shortest estimated partial path

(= length of partial path + underestimate of remaining)

until goal is reached.

Reject loops.

Phase 2: Extend all estimated partial paths until their length >= complete path to goal



End of phase 1: G reached

End of phase 2: 13>10



#### Admissible Heuristic

• If the heuristic function never overestimates the actual cost to get to the goal, then A\* is guaranteed to return a least-cost path from start to goal.



#### Admissible Heuristic

• If the heuristic function never overestimates the actual cost to get to the goal, then A\* is guaranteed to return a least-cost path from start to goal.

#### Time Complexity

- If goal state exists and is reachable from start state:
   Worst case O(b<sup>d</sup>) where d = depth(start,goal), b branching factor
- Otherwise, A\* will not terminate
- A good heuristic function allows A\* to prune away many of the b<sup>d</sup> nodes.



#### Monotone Underestimates of Remaining Distance

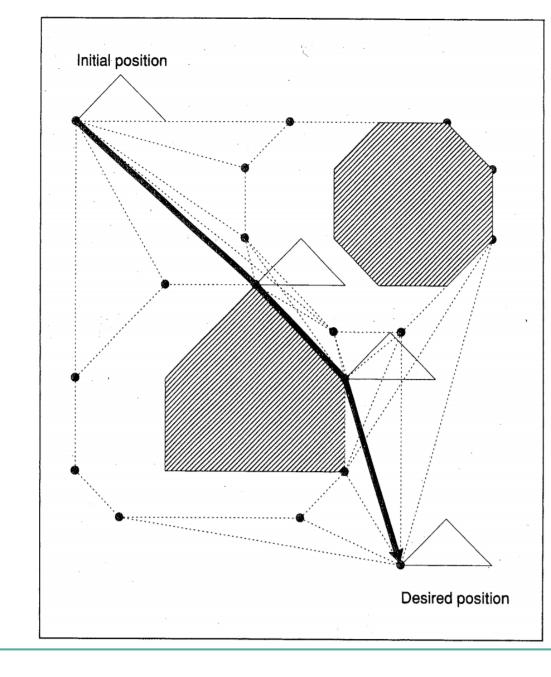
Heuristic function h(n) is "monotone" if and only if it satisfies the triangle inequality:

 $\Box$  h(n) = 0 if n=goal state

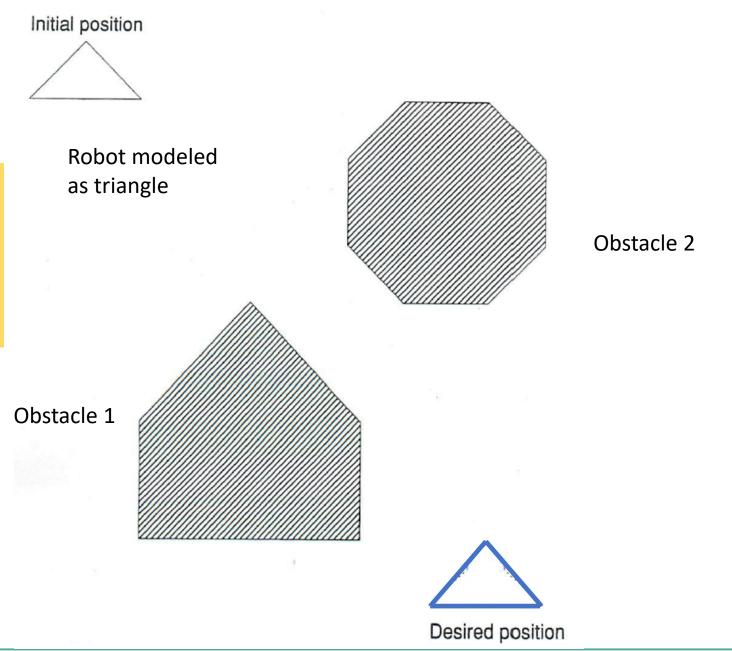
With a monotone heuristic, A\* is guaranteed to find an optimal path without processing any node more than once.

### Robot Path Planning with A\*

- Convert 2D Map into
   "Configuration Space" C
  - Robot represented as point
  - Obstacles represented as obstacles + fence
     = O
- Run A\* on Visibility Graph in
   "Free Space" F = C O









Al System Task:

Find shortest path for

robot to move from initial

robot hitting the obstacles

to goal position without

#### Create a fence around each obstacle:

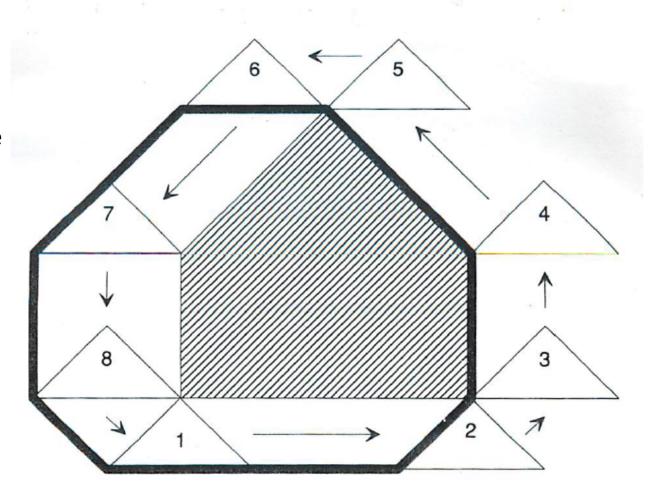
- 1. Select robot reference point
- 2. Slide robot shape around obstacle
- 3. Mark locations of reference point as fence

#### HERE:

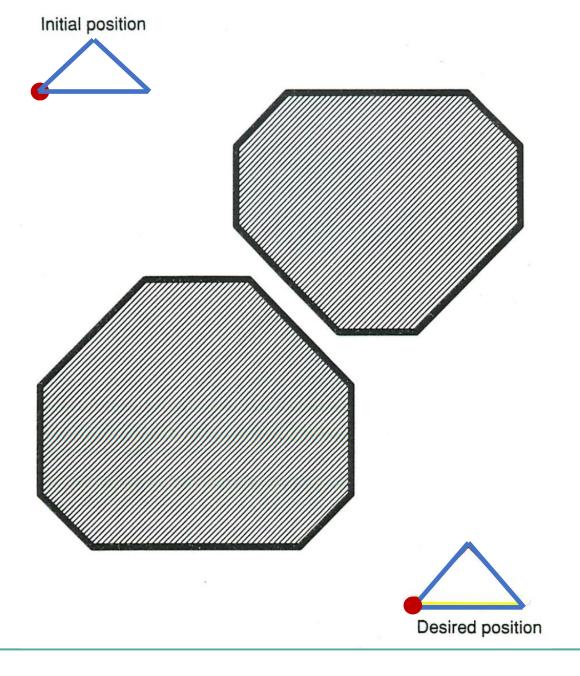
1. Reference point:



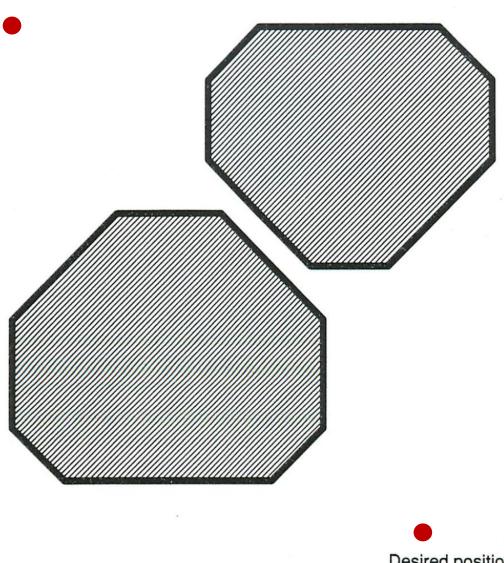
- 2. Eight unique shape positions
- 3. Fence: Thick black line





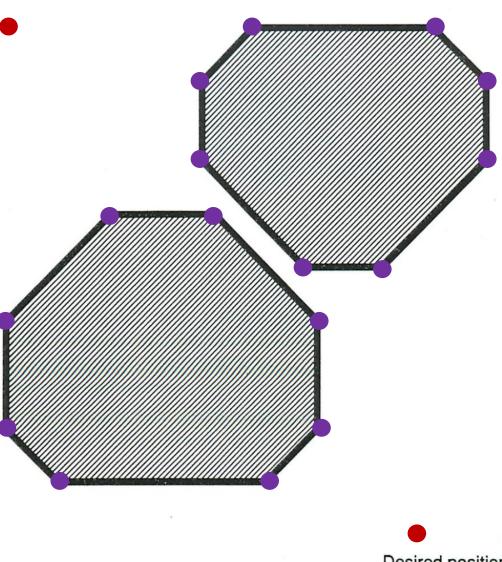


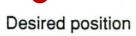






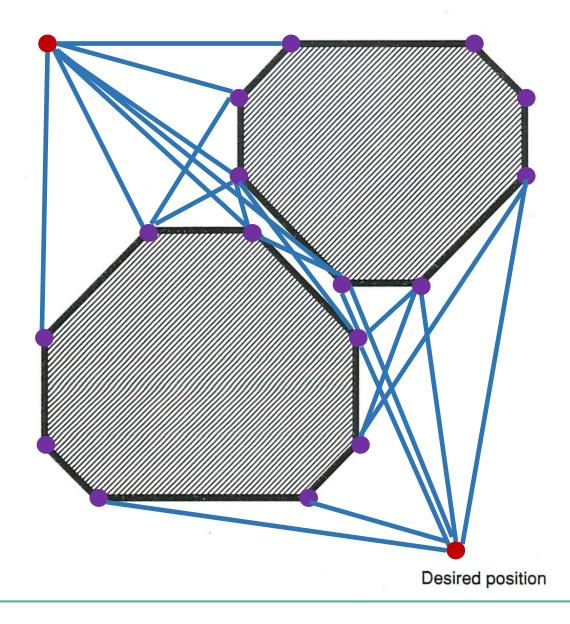






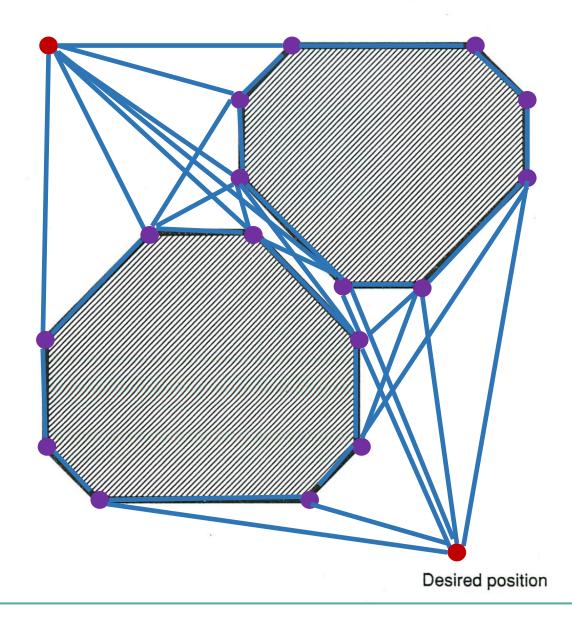


**Visibility Graph** 



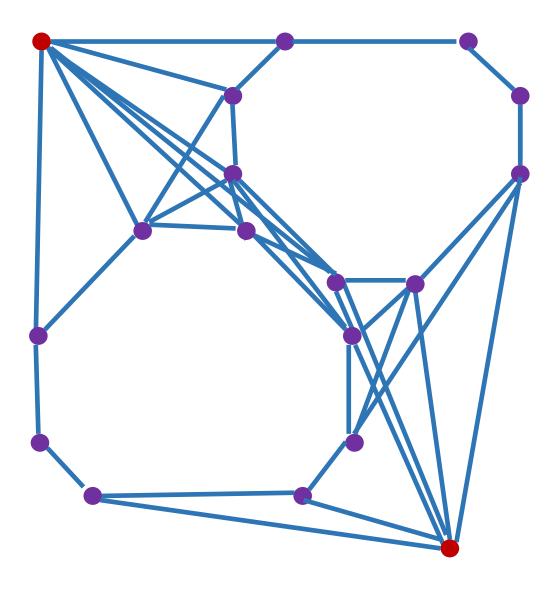


**Visibility Graph** 

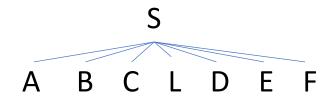




#### **Visibility Graph**



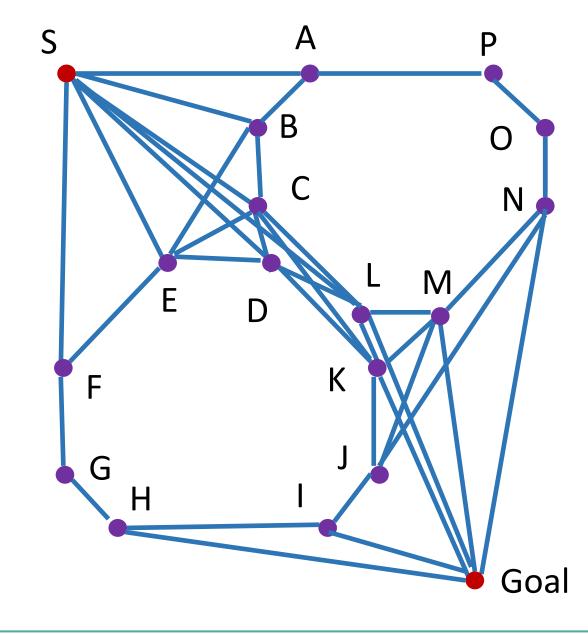




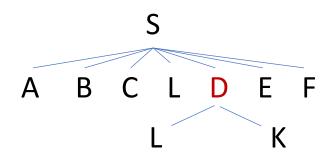
Node-to-node distance

+

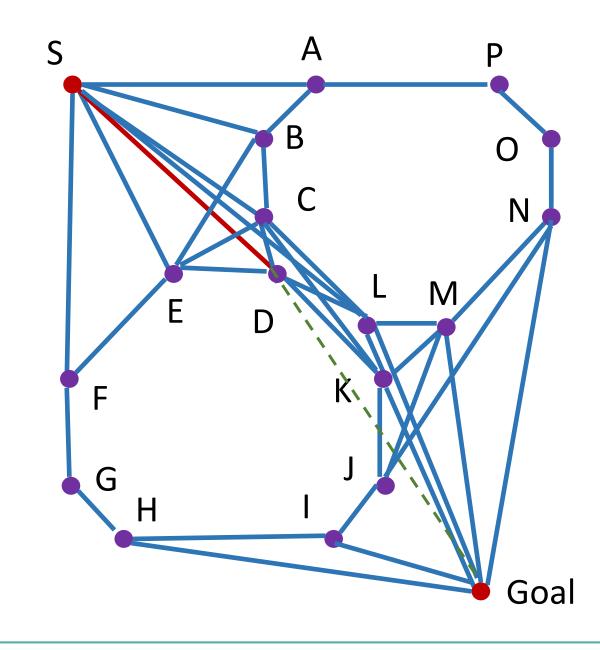
Underestimate to Goal



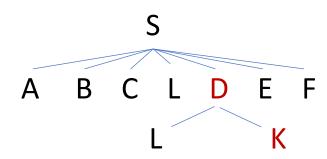




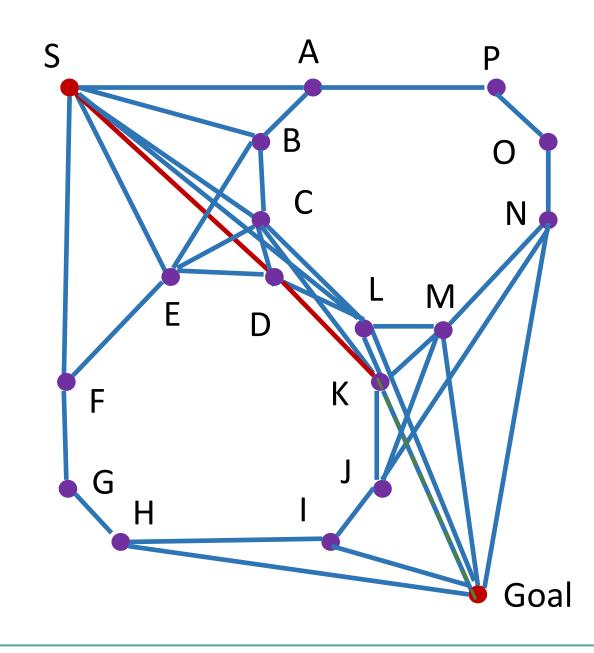
Node-to-node distance + Underestimate to Goal



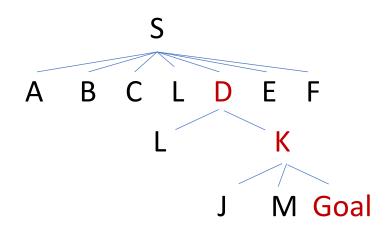




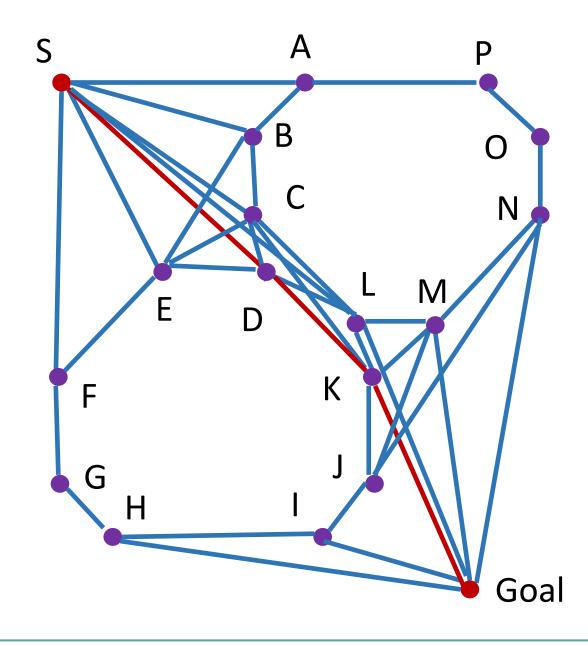
Node-to-node distance + Underestimate to Goal





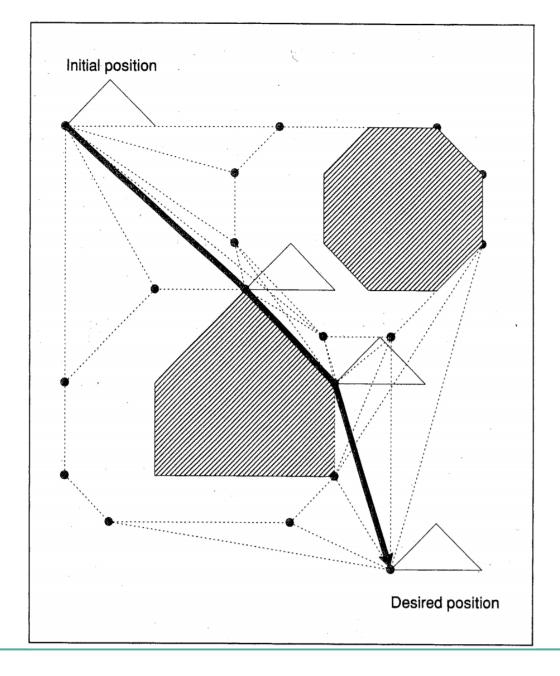


Node-to-node distance + Underestimate to Goal





Shortest path computed by A\*





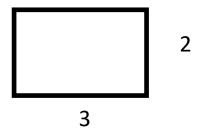
#### Your task:

Define Fence for the following robots & obstacles

1. Robot:



Obstacle:

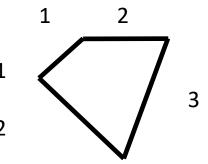


2. Robot:



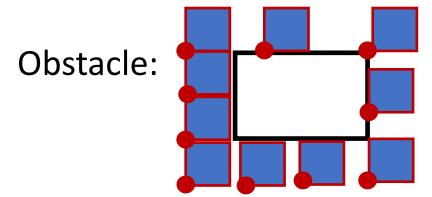
1

Obstacle:





1. Robot: 1

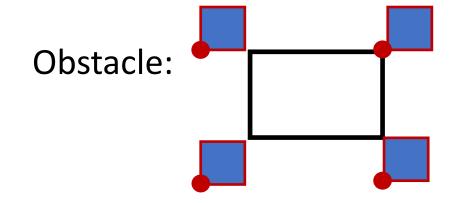


2. Robot: 1

Obstacle: 1 2

1 3

1. Robot: 1



2. Robot: 1

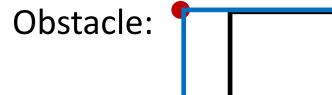
Obstacle: 1 2
1 2

1. Robot: Obstacle:

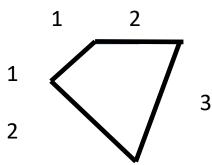


Obstacle: 1 2
1 2
2

1. Robot:

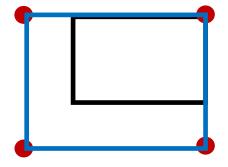


2. Robot: 1

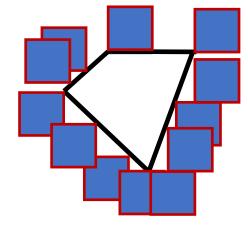


1. Robot: 1

Obstacle:

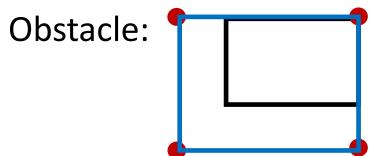


2. Robot: 1

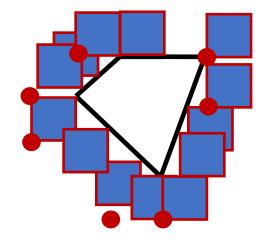




1. Robot: 1



2. Robot: 1

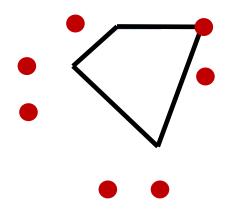




1. Robot: 1

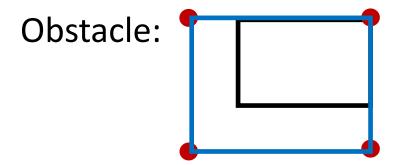
Obstacle:

2. Robot: 1

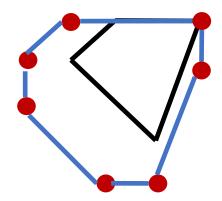




1. Robot: 1



2. Robot: 1



# General Problem of Robot Motion Planning

- Robot with k degrees of freedom: State or configuration of robot:  $(q_1, q_2, ..., q_k)$
- So far  $(q_1, q_2)$  for two-dimensional position

PUMA robot: 6 joint angles: (q<sub>1</sub>, q<sub>2</sub>, ..., q<sub>6</sub>)
 6D configuration space



# General Problem of Robot Motion Planning

Given initial point  $c_1$  and destination point  $c_2$ , in configuration space C: Robot can safely move between corresponding points in physical space if and only if

There exists a continuous path between  $c_1$  and  $c_2$  that lies entirely in the free space.



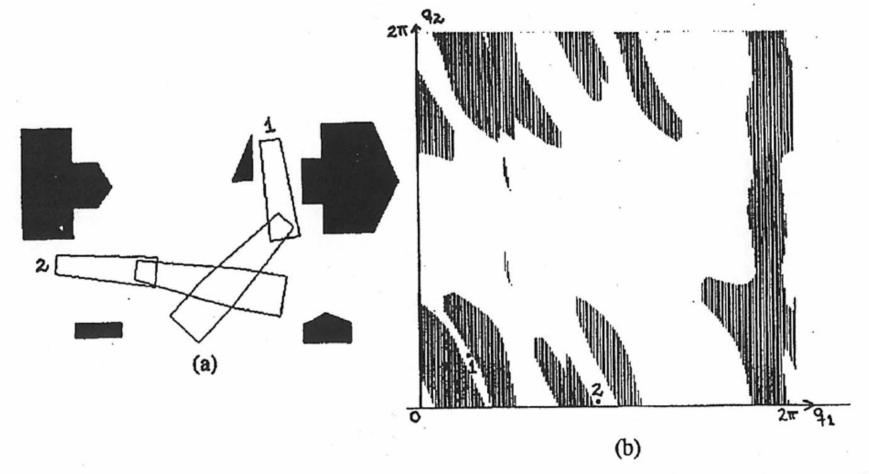
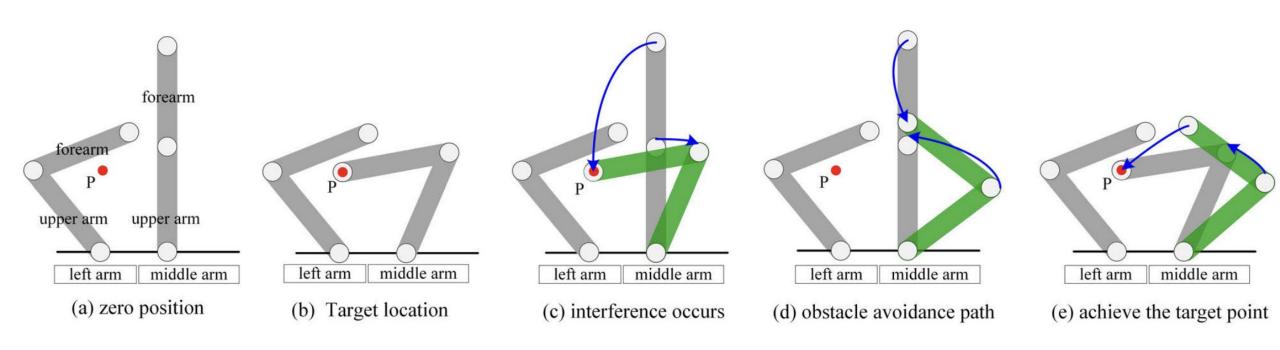


Fig. 4. (a) Two-link revolute manipulator and obstacles. (b) Two-dimensional C space with obstacles approximated by list of one-dimensional slice projections (shown dark). Initial and final position of manipulator are shown in input space and C space.

Lozano-Perez, 1987

#### Robot Obstacle Avoidance



Zhao et al., 2020



## Learning Outcomes of this Lecture

- Understand how search algorithms are evaluated
- Understand the unique properties of AI searching tasks (versus general search algorithms)
- Can explain 11 path-based search algorithms and run them on an example
- Can explain the dynamic programming principle

- Know what an admissible and a monotone heuristic function is for the A\* algorithm
- Can design a configuration space from a
   2D obstacle map & a translating robot
- Can design a visibility graph in free space
- Can run A\* on a visibility graph for robot path planning
- Understand configuration spaces of robot arms

