

# Dynamic Layer Instantiation as a Service

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**Motivation:** one of many fundamental problems with the current Internet architecture is the inability of a higher layer to request the service it requires from the lower layer. This inability, together with the desire to offer virtualized network services, have compounded existing network service management challenges. For example, in *Virtual Private Clouds* [1], some of these challenges include network isolation, custom addressing, and dynamic acquisition and release of virtual resources, with different Service Level Objectives (SLOs) and policies.

In response to these and other requirements, *e.g.* security or mobility, many *clean-slate architectures* arose. For example, XIA [2] provides a mechanism to securely support different address spaces (principals), NDN [3] offers a content distribution application architecture, Mobility-First [4] and Serval [5] focus on the mobility problem.

**Background:** our Recursive InterNetwork Architecture (RINA) [6] is based on the fundamental principle that *networking is Inter-Process Communication (IPC)*. RINA recurses the IPC service over different scopes, *i.e.*, range of operation. Specifically, a collection of distributed IPC processes with shared states is called a *Distributed Inter-process communication Facility (DIF)*. A *DIF layer* provides data transport services over a certain scope. By *DIF layer instantiation*, we mean that a layer is built by customizing management and data transfer control policies, to deliver predictable services to applications. Stacking DIF layers on top of each other allows networks to be built from smaller and more manageable layers of limited scope. The DIF layer instantiation is *dynamic* as layers can be discovered and created on the fly, and their mechanisms are able to respond and support policy adaptation as network states (*e.g.* quality of underlying services) change. **Our Contributions:** RINA enables private (*e.g.*, virtual or content delivery) networks to be dynamically instantiated, by customizing network management policies, *e.g.*, authentication, routing, addressing, and resource allocation, into a single layer, without the shortcomings of the TCP/IP architecture. To this end, we demonstrate with our prototype how such dynamic instantiation over lower-level service layers is a natural service that RINA can provide.

RINA subsumes the mechanisms and policies of all other architectures. For example, by layer isolation and explicit authentication of processes to join a layer, RINA supports secure address spaces as in XIA [2]; each DIF layer is a secure container. By naming application processes and not interfaces, RINA supports mobility and multihoming as in Serval [5] or Mobility-First [4]. Thanks to its directory mechanism, RINA supports content discovery applications as in NDN [3]. By

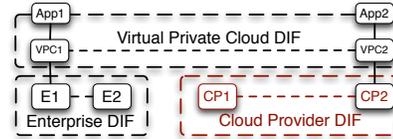


Fig. 1. Dynamic layer (DIF) formation: a higher level DIF (Virtual Private Cloud) is dynamically formed.

instantiating DIF layers with different scopes and policies, a DIF layer can be thought of as a virtual “slice” of a physical network, as envisioned by the Nebula Data Plane [7]. A DIF layer could also be a slice/layer over lower level slices. RINA references on mobility, security, and layer discovery can be found at [8].

**Demonstration:** we demonstrate the dynamic layer instantiation feature of RINA by creating on the fly a new, higher level *Virtual Private Cloud DIF*. The demonstration includes two IPC processes, VPC1 and VPC2, that initially use two separate private DIFs — an *Enterprise DIF* and a *Cloud Provider DIF* — to communicate with their respective local processes. Later on, an enterprise application process App1 asks for a flow service, so as to communicate with App2, a remote application process on the Cloud Provider DIF. Such request, handled by the underlying communication process VPC1, cannot occur unless there is a common underlying *Virtual Private Cloud DIF* to which both VPC1 and VPC2 subscribe.

Such dynamically instantiated DIF layers overcome many TCP/IP limitations. For example, since our DIF layer identifies applications (services) by location-independent names, they can move, migrate, or be multi-homed; since we support explicit (authenticated) enrollment into the DIF layer, communication through the layer is secure; and since DIF services are recursively built, they can be better coordinated and managed. Flavio and Yuefeng are both PhD students and will be setting up a demo over the GENI testbed [9], along with the poster, if accepted.

## REFERENCES

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