### A SOFTWARE AND HARDWARE ARCHITECTURE FOR NEXT-GENERATION AUTOMOTIVE SYSTEMS

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# Background







### Vehicle Growth in Electronics

- Electric vehicles, ADAS, IVI, V2X driving up cost and complexity of electronics
- Modern luxury vehicles have 50-150 ECUs source: Strategy Analytics, IHS Markit
- Global ECU market \$63.6 billion (2018) source: grandviewresearch.com
- Electronic share of total vehicle cost is rising exponentially



source: Statista 2017

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### Automotive Software Complexity

Growth in automotive electronics has given rise to growth in software complexity



Source: https://informationisbeautiful.net/visualizations/million-lines-of-code/

## Software Explosion

Software growth driven by increased vehicle functionality + increased ECU count



### ADAS – SAE 6 Levels of Driving Automation



Human monitors the driving environment

Based on: https://www.synopsys.com/automotive/autonomous-driving-levels.html



#### ADAS – SAE 6 Levels of Driving Automation



Human monitors the driving environment

Automated system monitors the driving environment

#### AUTOMOTIVE DOMAIN

•  $8 \rightarrow 16 \rightarrow 32$  bit microcontrollers

#### PC DOMAIN

• 64-bit CPUs, integrated GPUs



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#### Simple RTOS

• OSEK, FreeRTOS, Tresos, ECOS ...

Complex General Purpose OS

• Windows, Mac OS, Linux



#### **Reduce electronic costs**

- Replace ECUs with fewer hardware components
  - e.g., multicore industrial PC
- Consolidate ECU functions as software tasks
  - Easier to update, reconfigure, extend
- => Need for functional consolidation



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#### Address emerging real-time I/O needs

- Combined low-latency & high bandwidth data processing
- Google's self-driving car (2013) ~ 1GB/s data
- A.D. Angelica: http://www.kurzweilai.net/googles-self-driving-car-gathers-nearly-1-gbsec



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Functional safety and security (e.g., ISO26262, ISO21434)



#### Functional Consolidation => Need new vehicle OS

- Manage 100s of tasks on multiple cores
- Handle real-time low & high bandwidth I/O
- Provide safety, security and predictability
- Support mixed-criticality, fast boot, power management

Prohibitive complexity to write new OS from scratch

- Combine real-time with legacy code
- e.g. small RTOS + Linux
- Symbiotic solution



# Vehicle Vulnerabilities

Functional Safety (e.g., ISO26262) + Cybersecurity (e.g., ISO21434)

- ASIL classification based on Hazard Analysis and Risk Assessment
- ASIL = Exposure [E0-4] x Controllability [C0-3] x Severity [S0-3]



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Remote Surface Attacks Wi-Fi, Cellular, FM/AM radio, TPMS, Remote Keyless Entry, Bluetooth



# Moving Forward: DriveOS



### **DRAKO DriveOS**

DriveOS supports traditional hardware functions as software tasks running on a multicore virtualized platform



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### DRAKO DriveOS I/O

**USB-centric solution**: works with legacy devices + supports higher bandwidth future needs



\*Secure access to USB + CAN mediated by trusted I/O sandbox in DriveOS

### Reference Design: DRAKO GTE DriveOS



# DRAKO DriveOS Reference Stack



### **DRAKO DriveOS Functional Overview**





# Quest-V Separation Kernel (VEE'14, ACM TOCS'16)

- Monitors partition CPU cores, RAM, I/O devices among sandboxed guests
- Monitors have small trusted compute base no runtime resource management



### Quest-V Separation Kernel (VEE'14, ACM TOCS'16)

- Partitioning hypervisor statically partitions resources
- Separation kernel distributed collection of sandboxed components, indistinguishable from separate private machines for each component





### Quest RTOS (RTAS'11)

VCPUs are first-class entities within the RTOS

- Budgeted real-time execution of threads and interrupts
- Tasks → Main VCPUs (Sporadic servers: budget & period)
- Interrupts → IO VCPUs (PIBS: derive budget & period from Main VCPU)







# VCPU Control Flow



### VCPU Scheduling (RTAS'11)

Sandbox with 1 PCPU, n Main VCPUs (SS), and m IO VCPUs (PIBS)

- Ci = Budget Capacity of Main VCPU, Vi
- Ti = Replenishment Period of Vi
- Uj = Utilization factor for I/O VCPU, Vj
- Utilization bound feasibility test (with rate-monotonic scheduling of VCPUs):

$$\sum_{i=0}^{n-1} \frac{Ci}{Ti} + \sum_{j=0}^{m-1} (2 - Uj) \cdot Uj \le n \cdot (\sqrt[n]{2} - 1)$$





### Tuned Pipes (RTSS'18)

Like POSIX pipes but guarantee throughput and delay on communication

• Simpson's 4-slot (asynchronous) & FIFO (synchronous) buffering



**Boomerang** I/O subsystem in Quest-V supports real-time pipelines across Quest RTOS and legacy OSes

- Rate match tasks in pipeline to avoid blocking or missed data
- Quest appears as a real-time virtual device interface to Linux/Android

### DriveOS: Example OpenPilot ADAS + IC + IVI (EMSOFT'21)





### DriveOS: OpenPilot Control Loop Latency (EMSOFT'21)

ADAS Control Loop End-to-end Latency in presence of background Linux tasks
 Target bound = 10ms



\*Both Linux cases use PREEMPT\_RT. Optimized Linux maps USB interrupts to separate core

### Conclusions

Now is the time to look to alternative hardware + OS automotive solutions

DriveOS uses hardware virtualization for real time temporal and spatial isolation of software functions

- + Multicore PC-class platform replaces ECUs with software tasks
- + USB-centric I/O control
- + Symbiosis between RTOS & legacy OS
- + Real-time I/O & task pipeline processing

Fast startup of critical services on PC-class hardware (RTAS'22)





- Functional consolidation to drive down costs of electronics
- Centralized software stack to reduce hardware + code complexity
- Must consider OS challenges
  - More than just supporting driverless & connected cars
  - ML is great for object detection but an RTOS is needed to avoid objects!
- Real-time I/O is critical



## **Related Work**

Automotive Company / System	Operating System	Features		
DRAKO DriveOS™	Quest RTOS, Quest-V Separation Kernel + Yocto Linux / Android	<b>Centralized</b> ; Quest/Linux/Android sandboxes IC, IVI, HVAC, Powertrain,ADAS, etc Simulink Multi-OS Support		
Toyota Entune 3.0 (Future: Arene)	Automotive Grade Linux (Arene: Apex.OS)	Infotainment (Arene will support autonomy)		
BMW OS7 and OS8 (iX)	Greenhills Integrity RTOS + Linux	[Linux] Infotainment, IC [RTOS] RT vehicle control functions		
Polestar + Google	Automotive Android	Infotainment		
Nvidia Drive OS	Nvidia Hypervisor, QNX Neutrino RTOS, Linux	ADAS, Linux + QNX SDK		
Ford Sync 3	QNX (current); Android (future)	Microkernel, RTOS, Infotainment		
TTTech	Car.OS	Supports AUTOSAR, Linux, QNX + others <b>Centralized</b> ; IC, IVI, ADAS, HVAC, Powertrain		
Mercedes Benz	MB.OS	<b>Centralized</b> ; RTOS + Linux support IVI, Powertrain, ADAS, Body Control, HVAC		
Tesla	Linux + FSD (Full Self Driving)	Infotainment (AMD for Model 3 & Y), ADAS		









# System Software Safety

Temporal and Spatial Isolation

• Ensure critical tasks are free from interference from less critical tasks

Timing and Functional Safety

- Ensure timing-critical tasks meet deadlines
- Functionally correct output values for given inputs

Correct Information Exchange

No loss, duplication or corruption of data

Memory Safety

No buffer overruns, stack under/overflow, invalid memory addressing

IO Safety

Controlled access to IO devices



### System Security

Integrity

- Avoid attacker compromizing critical functionality
  - e.g., Miller & Valasek, 2014 Jeep Cherokee CAN attack via remote access to IVI
- Resource partitioning and access only via secure interfaces
- Validate arguments to functional interfaces

Confidentiality

- Avoid leaking sensitive data (CAN packets, personal information, app data,...)
- Encrypt data or enforce information flow policies
- Eliminate side channels (e.g., via caches possibly use cache/page coloring)
- Use containerization for critical components

Access Rights

- Avoid user gaining elevated accesses to resources beyond allowed rights
  - e.g., CVE-2019-5736 Breaking out of Docker via RunC
- Enforce a capability mechanism on access to resources
- Digitally sign software images



### Today's ECU Vehicle Network





### **DRAKO DriveOS**

Leverage the Quest-V separation kernel

- Open Source
- Partitions CPU cores, RAM, I/O devices among guests

Co-locate Quest RTOS with Linux and Android guests on same hardware

Real-time interface for device I/O

- + Processing moved to PC
- + I/O via e.g. USB-CAN or custom control-class interface





### Example: Quest-V for DriveOS





# Cache Partitioning (Spatial and Temporal Isolation)

- Shared caches controlled using color-aware memory allocator [COLORIS PACT'14]
- Quest-V uses EPTs to map guest physical to machine physical addresses



- Last-level cache occupancy prediction based
- on h/w performance counters
  - local (core) + global (all core) cache hits and misses between scheduling points [Book Chapter, OSR'11, PACT'10]



# Quest RTOS – USB Scheduling (RTAS'13)

USB 2/3.x Bus scheduler

Each periodic request represented as a tuple (wi, ti)

- wi time to send transaction i
- ti time interval of transaction i

Given set of n tuples {(w1,t1), (w2,t2),...,(wn,tn)}, is there an assignment of USB transactions to 125uS microframes, such that no frame is over-committed?

A request assigned to microframe f is also assigned to microframe f+n\*ti,  $n \in N$ 

Using variant of first-fit decreasing packing algorithm, shown to outperform Linux

- Sort by decreasing wi (largest first)
- First pick request based on smallest ti, breaking ties with largest wi



# Quest RTOS – USB Scheduling (RTAS'13)

- Consider all permutations of 1 to 5 requests
- Intervals: 2, 4, 8, 16 microframes
- Packet Sizes: 32, 64,...,1024 bytes
- Quest  $\approx$  150 thousand failed schedules
- Linux  $\approx$  95 million failed schedules





### Boomerang Inter-OS Task Pipeline Example (RTAS'20)

Boomerang tuned pipe path (1) spans Quest + Linux + USB-CAN

Boomerang tuned pipe path (2) spans Quest + USB-CAN



DriveOS: Boomerang Results



Boomerang sub-system in DriveOS meets communication timing guarantees A Linux SMP (multicore) OS with real-time extensions cannot perform I/O predictably

#### Jumpstart Power Management (RTAS'22)

- PC hardware requires Firmware POST, bootloader, device & service initialization to boot OS
- DriveOS uses Jumpstart ACPI S3 suspend-to-RAM & resume-from-RAM for low latency restart of critical tasks (e.g., CAN gateway services)



#### Jumpstart Power Management (RTAS'22)

- Jumpstart services span all guests
  - RTOS coordinates suspension but enables parallel reboot
- Potential for ACPI S4 suspend-to-disk using non-volatile memory (e.g., Intel Optane)
  - Eliminates system power usage during suspension



### DriveOS: Screenshot 1/4





### DriveOS: Screenshot 2/4





### DriveOS: Screenshot 3/4



#### DriveOS: Screenshot 4/4



# Simulink Multi-OS Modeling and Code Generation

- Model-based design for Multi-OS target
- Automatic support for nested ELF binaries with inter-sandbox RPC bindings





	ŕ		
thread1: 50/1100			
		LinuxThr	ead



QuestSandbox

LinuxSandbox



# Quest(-V) Simulink Blocks

Shared memory intertask/sandbox communication







#### CAN-bus Management





Example: Quest HVAC CAN - Shared Memory Logic



# Mapping a Function to a Quest VCPU

Configuration Darameters

Configurable Parameters:

- 1. Target Sandbox:
- 2. Task Budget (C)
- 3. Execution Period (T)

	configuration Pa	arameters				
neters: box: (C) eriod (T)	Block Parameters: vcpusetup1 A Level-2 C S-function for Quest VCPU s For code generation, it creates a pthrea the pthread. For simulation it does not implement an some basic setup. Parameters Name of the function of the thread hva Target OS Sandbox • Quest • Linux Worst-case Execution Time (us) 500 Period (us) 1000 OK Cancel	Setup (mask) (link) d and binds a main VCPU to ything significant, other that c_control	three	vcpusetup ad1: 500/2000	A VCPU is bo automotive fu output signal <b>vcpusetup</b> bl	ound to an Inction via the link of the lock
Set up ne connect t	ew channel or o an existing one		itialize sender_ch		fun channel_no channel_no1	

