Tuned Pipes: End-to-end Throughput and Delay Guarantees for USB Devices

Ahmad Golchin, Zhuoqun Cheng and Richard West Boston University

Motivations

- Cyber-physical systems
- Ubiquity of USB
- Sensor-actuator loops
- Need for predictable I/O communication
 - Between device & application tasks
- Avoid manually fine-tuning system parameters for control & data flow

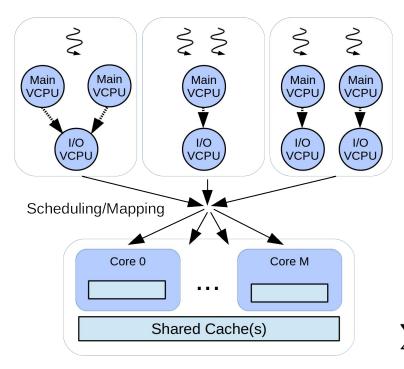
Contributions

- Tuned Pipes system framework
 - Guarantees end-to-end latency and throughput requirements between USB devices and host tasks
- A host controller driver with early demultiplexing
 - Allows USB bottom-half handler to run with the right priority and in a timely manner as opposed to Linux
- Extended our previous USB bus scheduling algorithm to comply with xHCI

Quest RTOS

- Real-time OS supporting multicore x86 platforms
 - Intel's Aero, UP, UP2, Skull Canyon, Edison, Minnowmax,...
- Dual-mode kernel
- Unified task and I/O (bottom-half) scheduling through time-budgeted virtual CPUs (VCPUs)
 - Tasks scheduling: Main VCPUs
 - Interrupt bottom-half scheduling: I/O VCPUs
- More info: www.questos.org

VCPU Scheduling in Quest RTOS

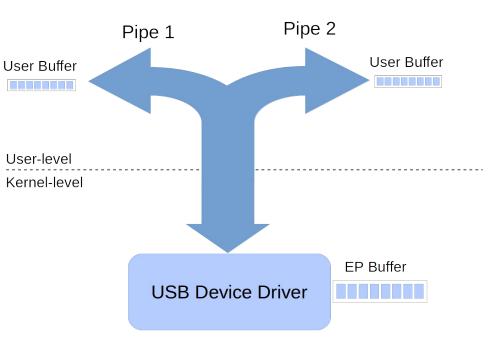


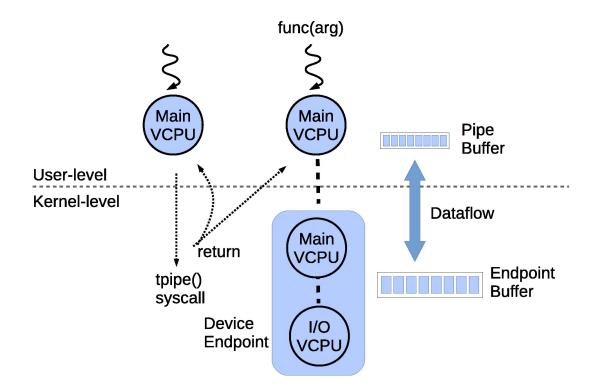
- Main VCPUs
 - Sporadic Server + RMS
 - Guarantees budget C every period T for tasks
- I/O VCPUs
 - PIBS
 - BW limited by utilization factor Uj
 - Inherits T from the task
- Temporal isolation condition:

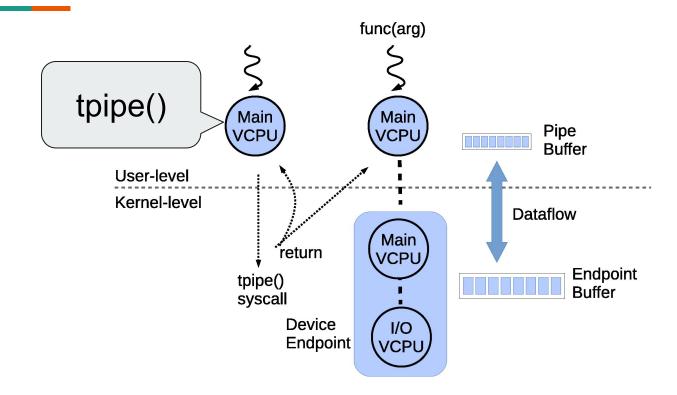
 $\sum_{i=0}^{n-1} rac{C_i}{T_i} + \sum_{j=0}^{m-1} (2-U_j). \, U_j \leq n(\sqrt[n]{2}-1)$

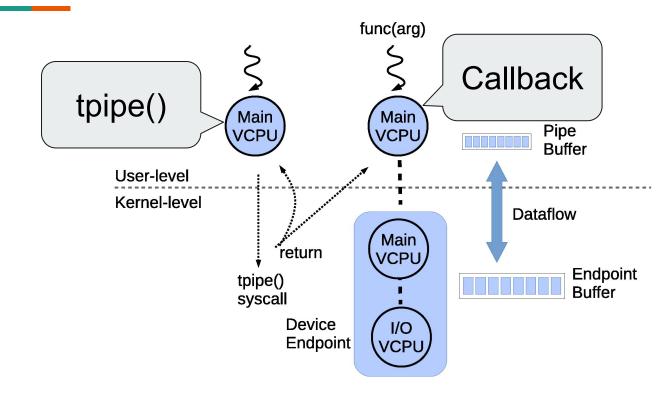
Tuned Pipes

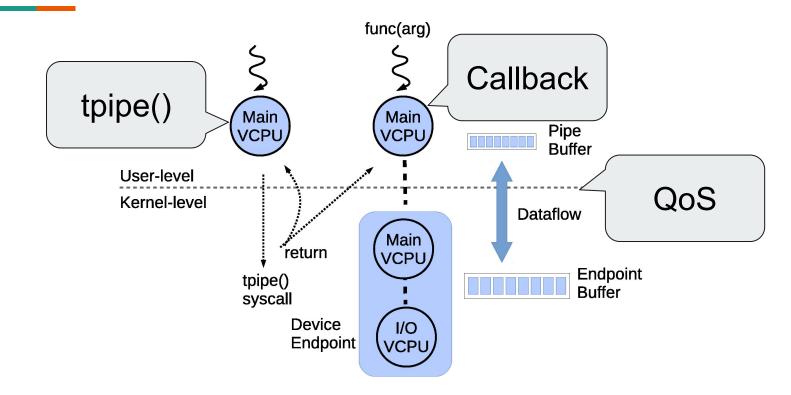
- Host-to-device
 communication channel
- Throughput and delay bounds (QoS)
- Temporal isolation
- Endpoint-pipe: 1:N registered by drivers











QoS Specification:

- Execution Time (C)
- Throughput (λ)
- IO Buffer Size (B)

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Main VCPU Parameters
$$\begin{cases} C = 1ms \\ T = 128*8 / 512000 = 2ms \end{cases}$$

Tuned Pipes - Kernel API

Endpoint:

- Endpoint attributes
- IOVCPU & sched param
- MainVCPU & sched param

Endpoint Attributes:

- Max # of Channels
- Max Throughput
- Min Latency
- Min/Max EP Buffer Size
- Min/Max Packet Size

Tuned Pipes - Kernel API

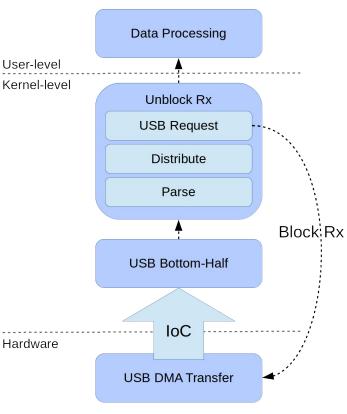
Example

- 4 channels at 500Kbps
- 1 channel at 250Kbps
- max_tput = 2.25Mbps
- ebuf_sz = 4KB
- Driver applies Little's law to set proper budget and period for it's I/O thread:



4 Delay contributors

- User thread
- Driver thread
- DMA of data
- USB bottom-half



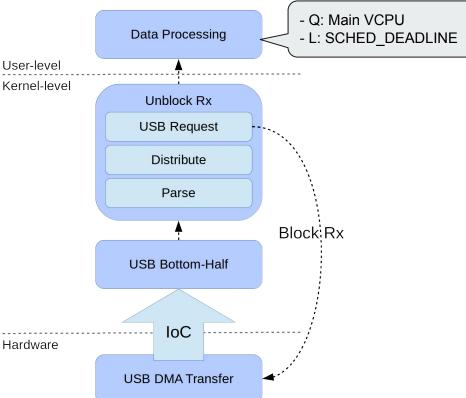
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Question:

• USB bottom-half

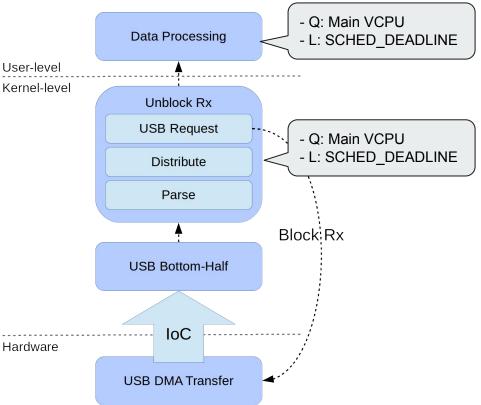
How to enforce QoS?



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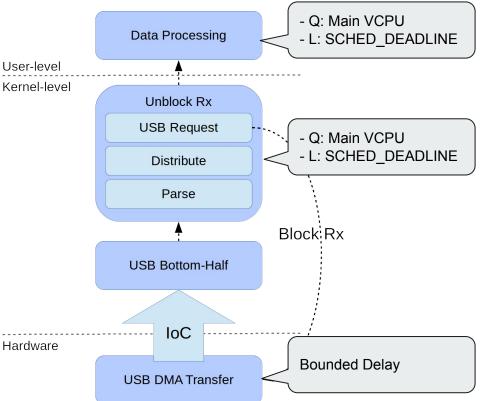
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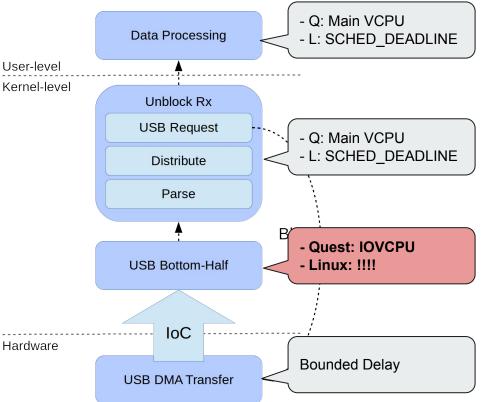
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End-to-end Data Path - Challenges

Challenges with Linux:

- USB BUS scheduling
- USB bottom-half handler priority mismatch!

What currently happens:

- Soft-IRQs

Highest priority until MAX_SOFTIRQ_RESTART \rightarrow Low priority

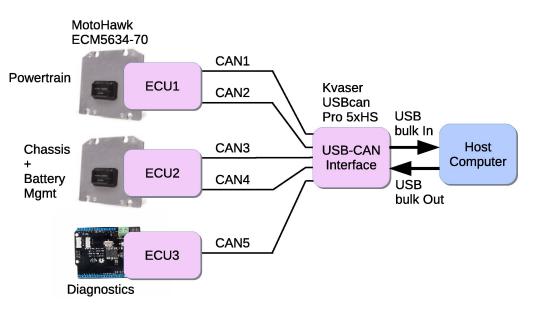
- Threaded-IRQs (e.g. PREEMPT_RT)

Fixed SCHED_FIFO priority (Default: 50)

Experimental Environment

CAN Interface

- Kvaser USBcan Pro 5xHS
- 5 channels: up to 1Mbps w/ 4KB buffer
- 2 ECUs: each exposing 2 channels
- 1 Arduino UNO + CAN-BUS Shield



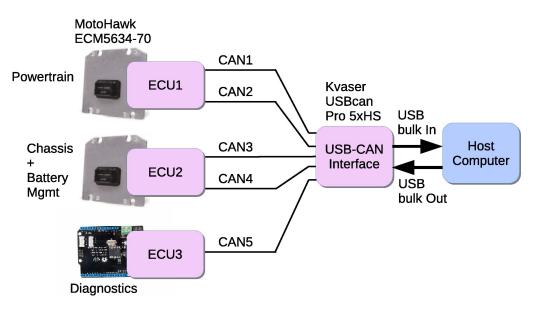
Experimental Environment

UPSquared SBC

- Dual-core Celeron N3350 @ 1.1 GHz
- xHCI 1.1 Interface

Quest RTOS

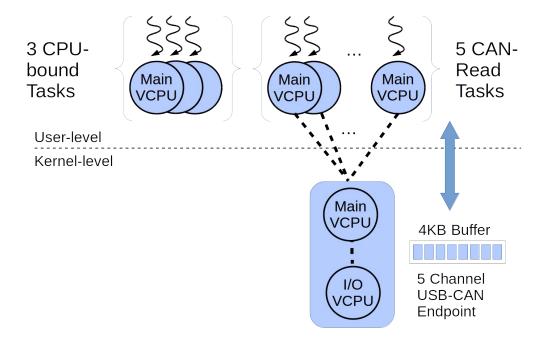
- VCPU Scheduling Ubilinux (PREEMPT_RT)
- SCHED_DEADLINE

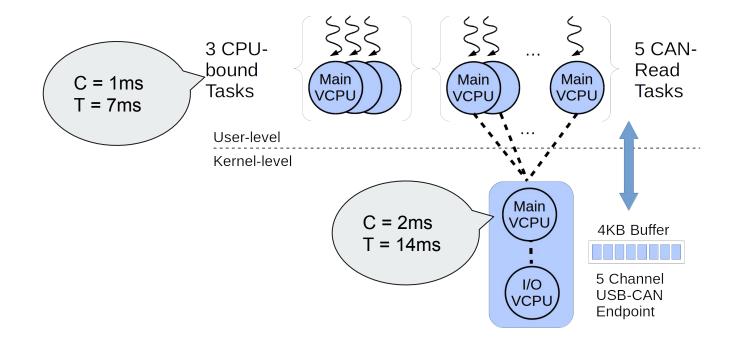


Objective: Receiving frames without:

- Loss of CAN packets
- Intervening with other tasks of higher priority Generated data traffic:

Bus	CAN1	CAN2	CAN3	CAN4	CAN5
Bandwidth (bps)	500K	250K	500K	500K	500K
Throughput %	10	20	30	40	69





Observations:

- Quest:
 - No buffer overrun
 - Negligible interference
- Linux:
 - 230 overruns over 30 seconds
 - 405 overruns over 60 seconds
 - More interference

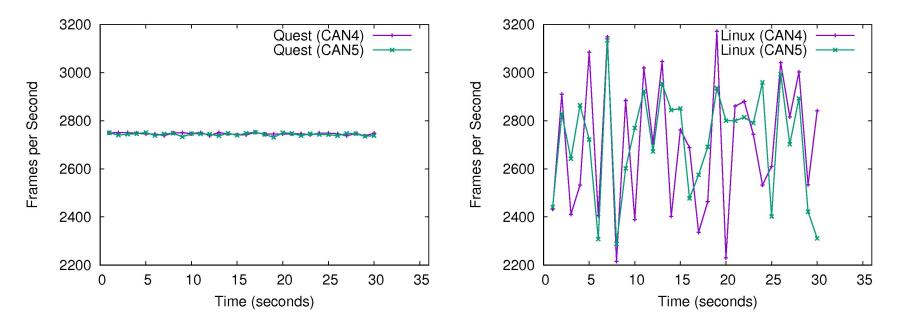
Test 2 - End-to-end Guarantees - Rx

Objective: Guaranteeing throughput using tuned pipes

- 5 Tuned pipes receiving data
- CAN 4 & 5 Throughput: 2730 to 2752 fps
- QoS: tput=2752, IObufsz=128, exec_time=2ms

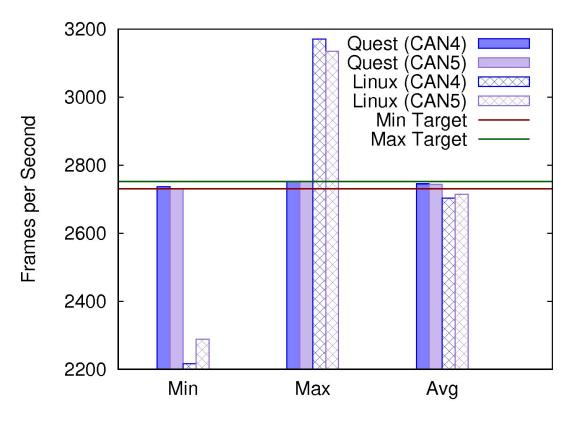
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Test 2 - End-to-end Guarantees - Rx



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Conclusions

- Tuned pipes abstraction
- Auto-tuning of system parameters
- Guarantee of throughput and delay constraints
 Not solved with SCHED_DEADLINE in Linux
- Early demultiplexing of entities waiting for INT
- Handling BH with the RIGHT priority (IOVCPU)
 - Not solved with PREEMPT_RT Linux patch

Thank you!

Comments or Questions?

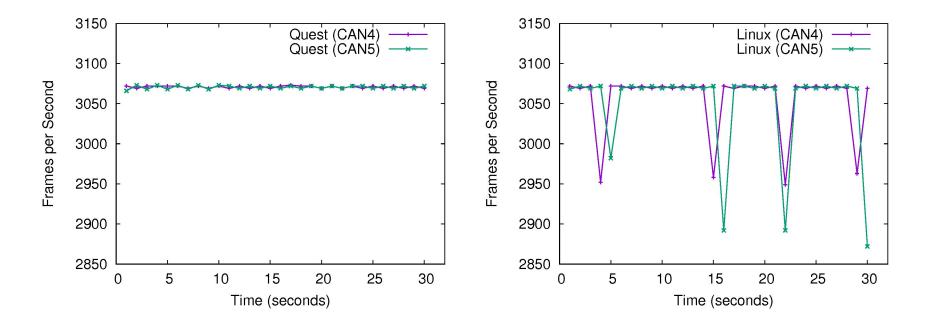
Test 3 - End-to-end Guarantees - Tx

Objective: Guaranteeing throughput using tuned pipes

Similar to the previous test, except:

- CAN 4 & 5 Receiving data every 325.4 to 327.5 uS
- Arrival rate: 3053 to 3073
- QoS: tput=3073, IObufsz=128, exec_time=2ms

Test 3 - End-to-end Guarantees - Tx



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