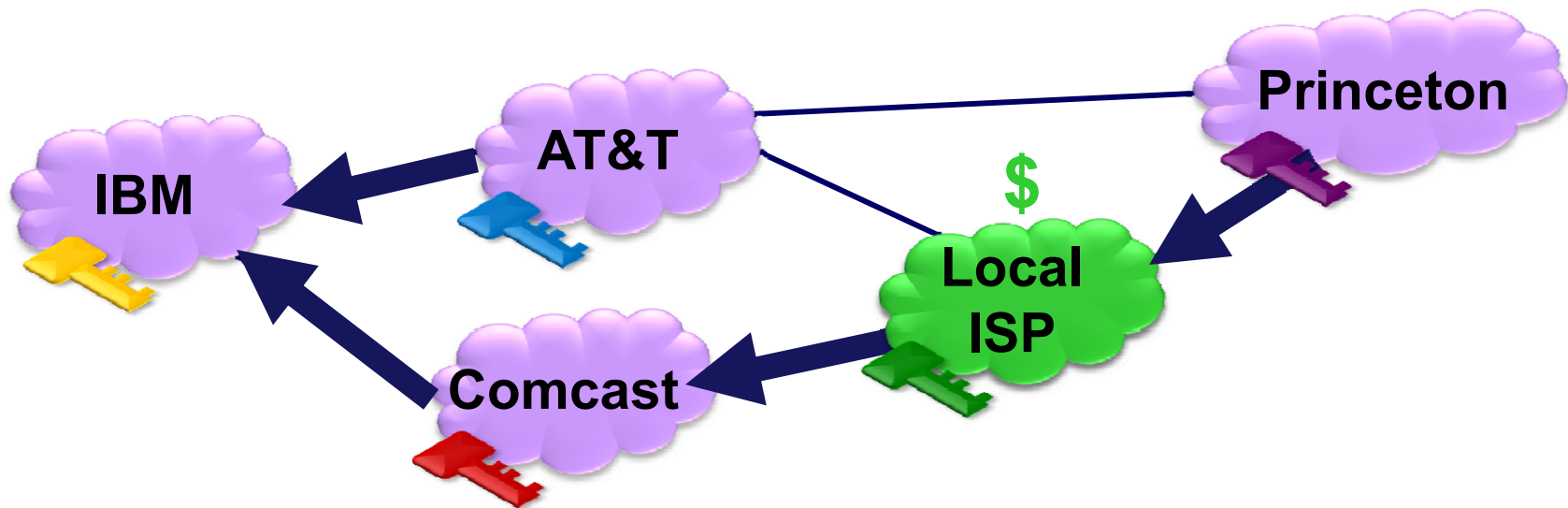


# Rationality and Traffic Attraction

## Incentives for Honest Path Announcement in BGP



**Sharon Goldberg**

Shai Halevi

Aaron D. Jaggard

Vijay Ramachandran

Rebecca N. Wright



Princeton University

SIGCOMM 2008



# Incentives and Security

We use game theory to understand which secure protocols should be deployed in the Internet.

**We ask:** Does traffic on the Internet actually follow the paths announced in **BGP**?

**Approach:** Assume that nodes are economic entities

- They are **rational** -- try to maximize utility.



**Our Results:** Mostly **bad news**.

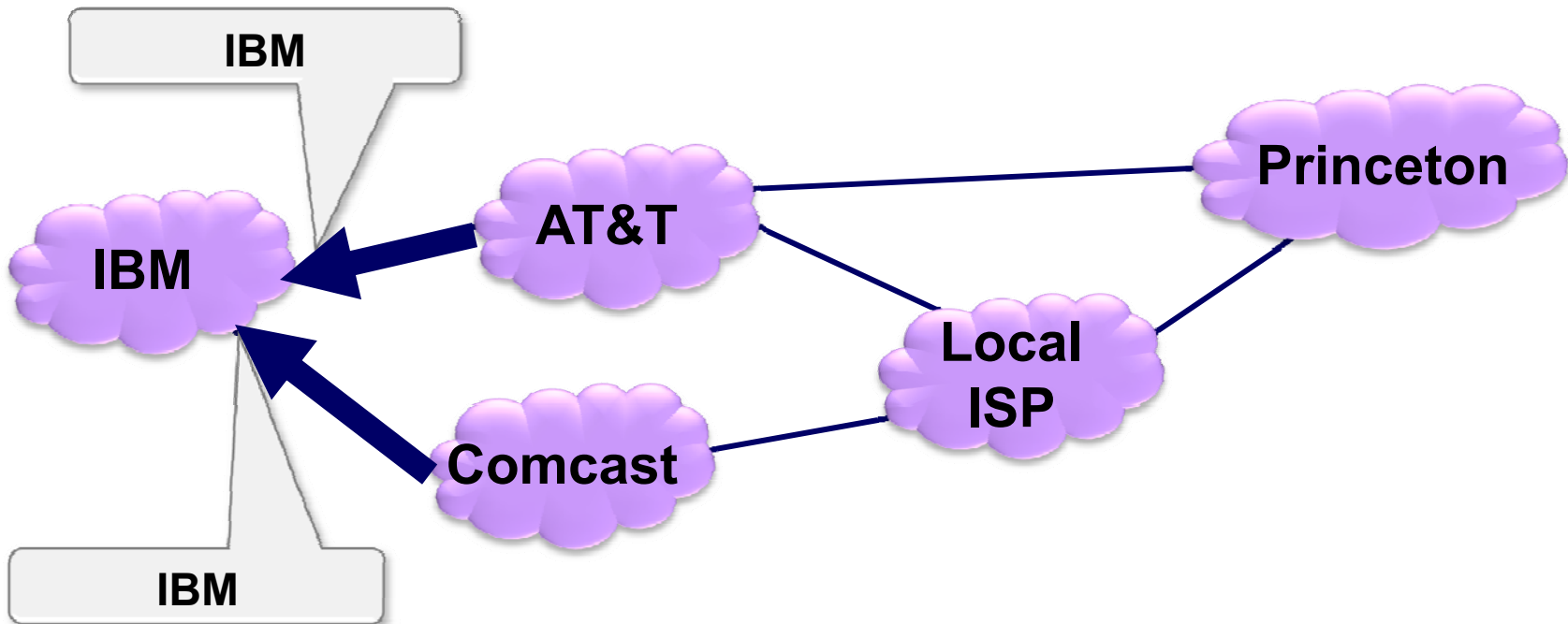
- We find that **cryptographically authenticating routing messages** is not sufficient.
- ... unless we also make **unrealistic assumptions about routing policies**.
- Results are mostly descriptive, not prescriptive





# BGP: The Interdomain Routing Protocol (1a)

The Border Gateway Protocol (BGP) is the routing protocol that sets up paths between Autonomous Systems (ASes).



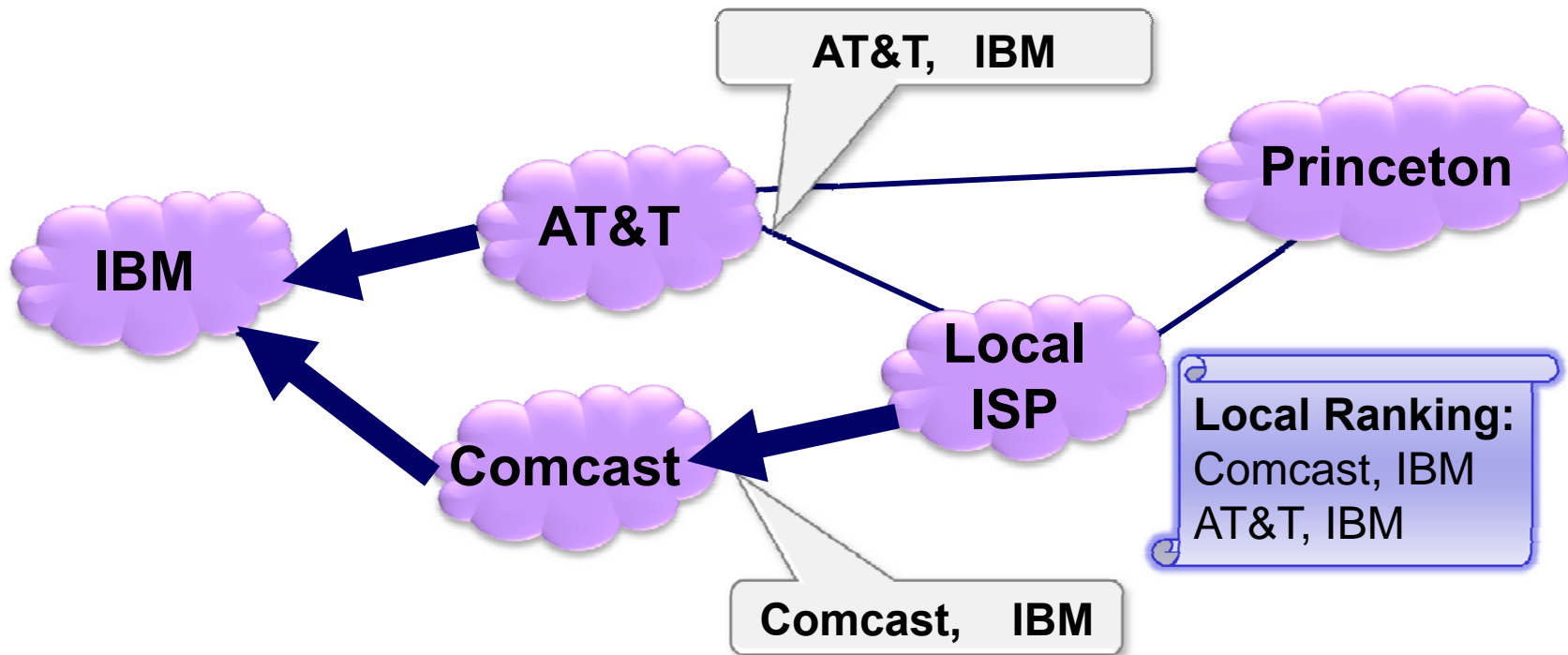
**Forwarding:** Node use **single** outgoing link for all traffic to destination.

**Rankings:** Static and local; usually based on economic relationships.



# BGP: The Interdomain Routing Protocol (1b)

The Border Gateway Protocol (BGP) is the routing protocol that sets up paths between Autonomous Systems (ASes).



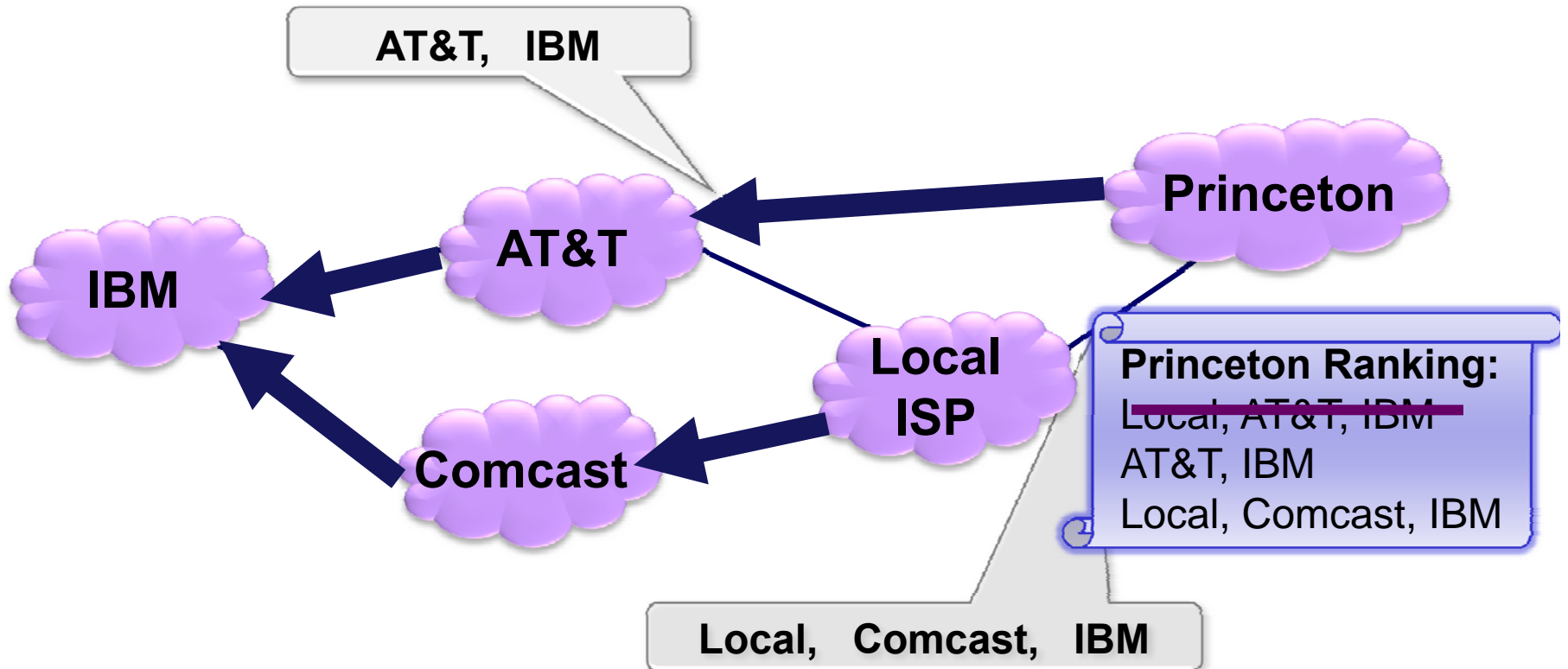
**Forwarding:** Node use **single** outgoing link for all traffic to destination.

**Rankings:** Static and local; usually based on economic relationships.



## BGP: The Interdomain Routing Protocol (2)

The Border Gateway Protocol (BGP) is the routing protocol that sets up paths between Autonomous Systems (ASes).



**Forwarding:** Node use **single** outgoing link for all traffic to destination.

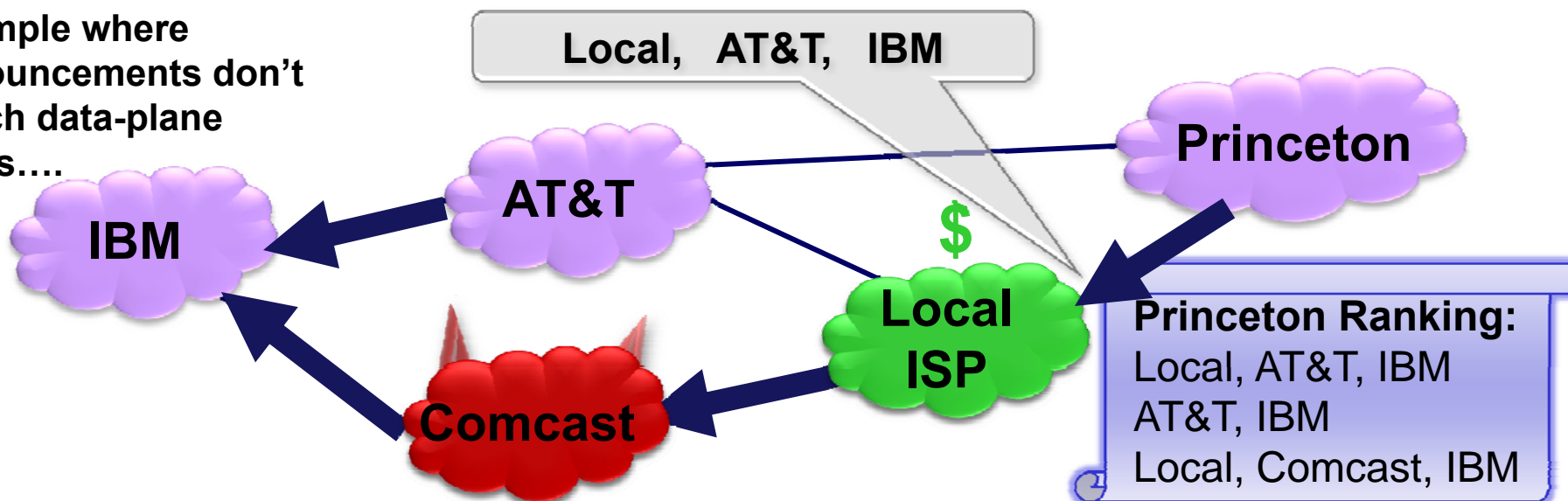
**Rankings:** Static and local; usually based on economic relationships.



## Matching the Data Plane

**BGP announcements match AS-paths packets take in data plane.**

Example where announcements don't match data-plane paths....



**This way, ASes can use BGP messages:**

1. To avoid ASes perceived as adversarial / unreliable
2. To choose high performance paths
3. As part of an accountability framework

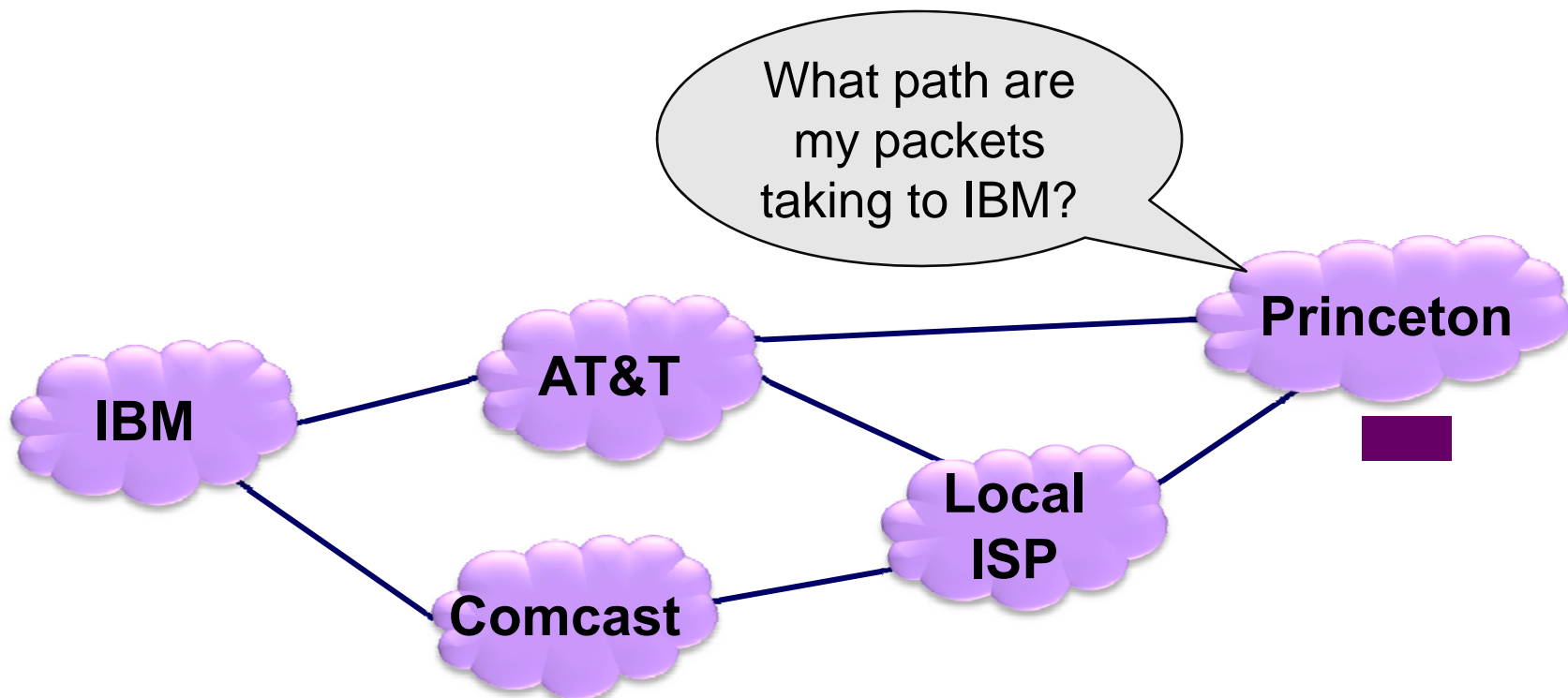


# Data Plane Approaches for Matching BGP Messages

## Secure Data-Plane Protocols:

- Packet Passports [\[LYWA-06\]](#) Packet Obituaries [\[AMISS-07\]](#)  
Truth in advertising [\[WBAGS-07\]](#) Failure Localization [\[BGX-08\]](#)

**✗ Secure AS-path tracing protocols incur high overheads**





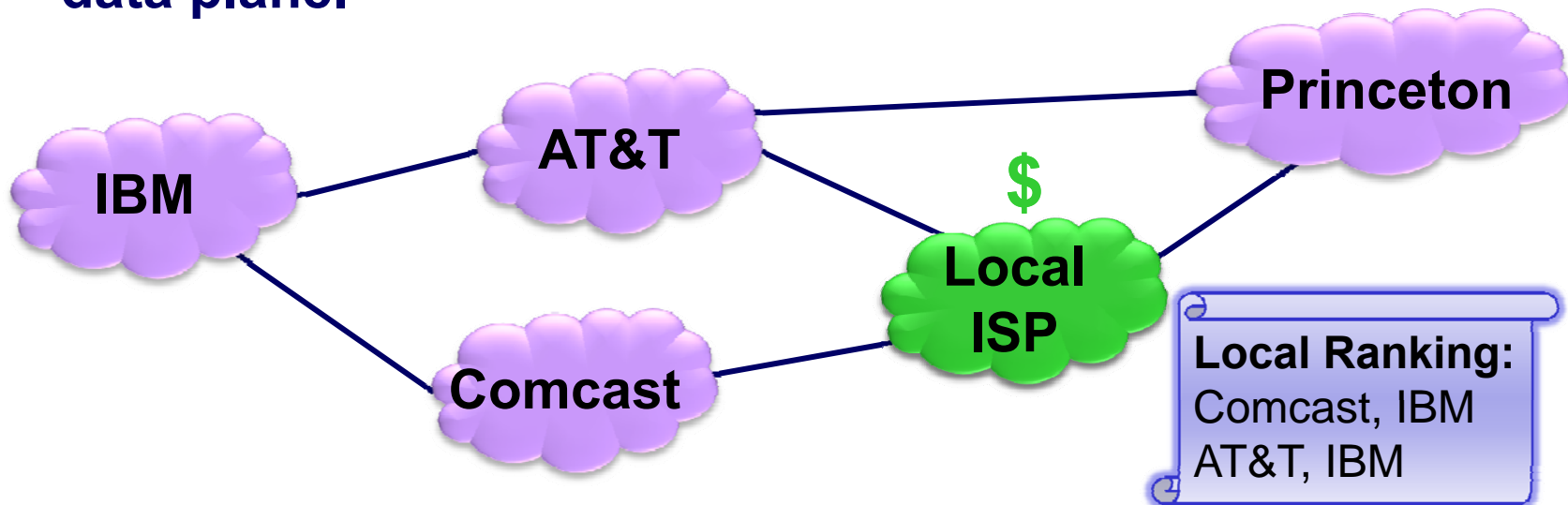
# Routing Protocol Approaches to Match Data Plane

## Routing Protocols + Game Theory:

- [NR-01] [FPS-01] [FPSS-05] [PS-04] [FKMS-05]  
Shortest-path policy / Next-hop policy [FRS-06] [FSS-07]  
Secure BGP [LSZ-08]



- ✓ **Corollary:** If \_\_\_\_\_, rational ASes have no incentive to unilaterally deviate from announcing paths that match data plane.



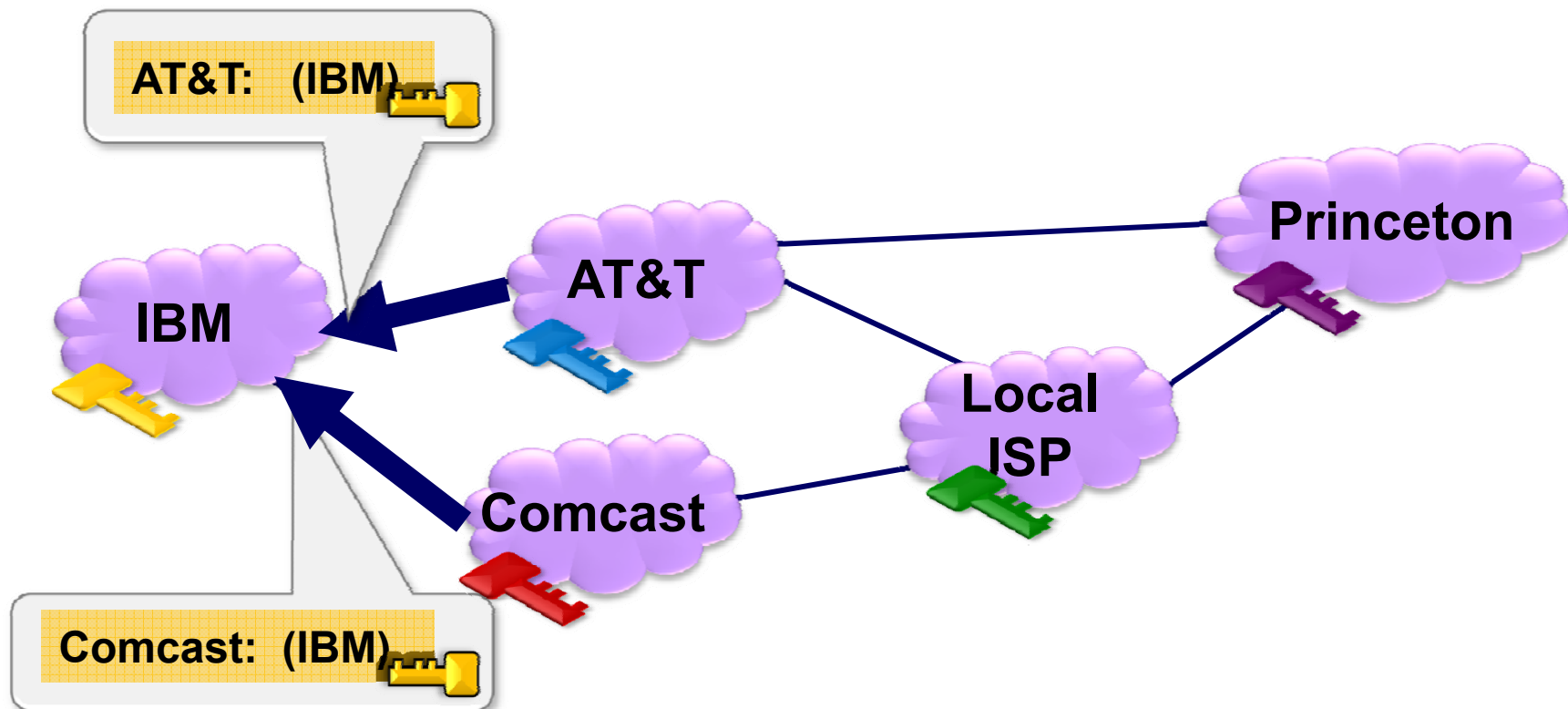




## Secure BGP (1a)

If **a** announced path **abP** then **b** announced **bP** to **a**

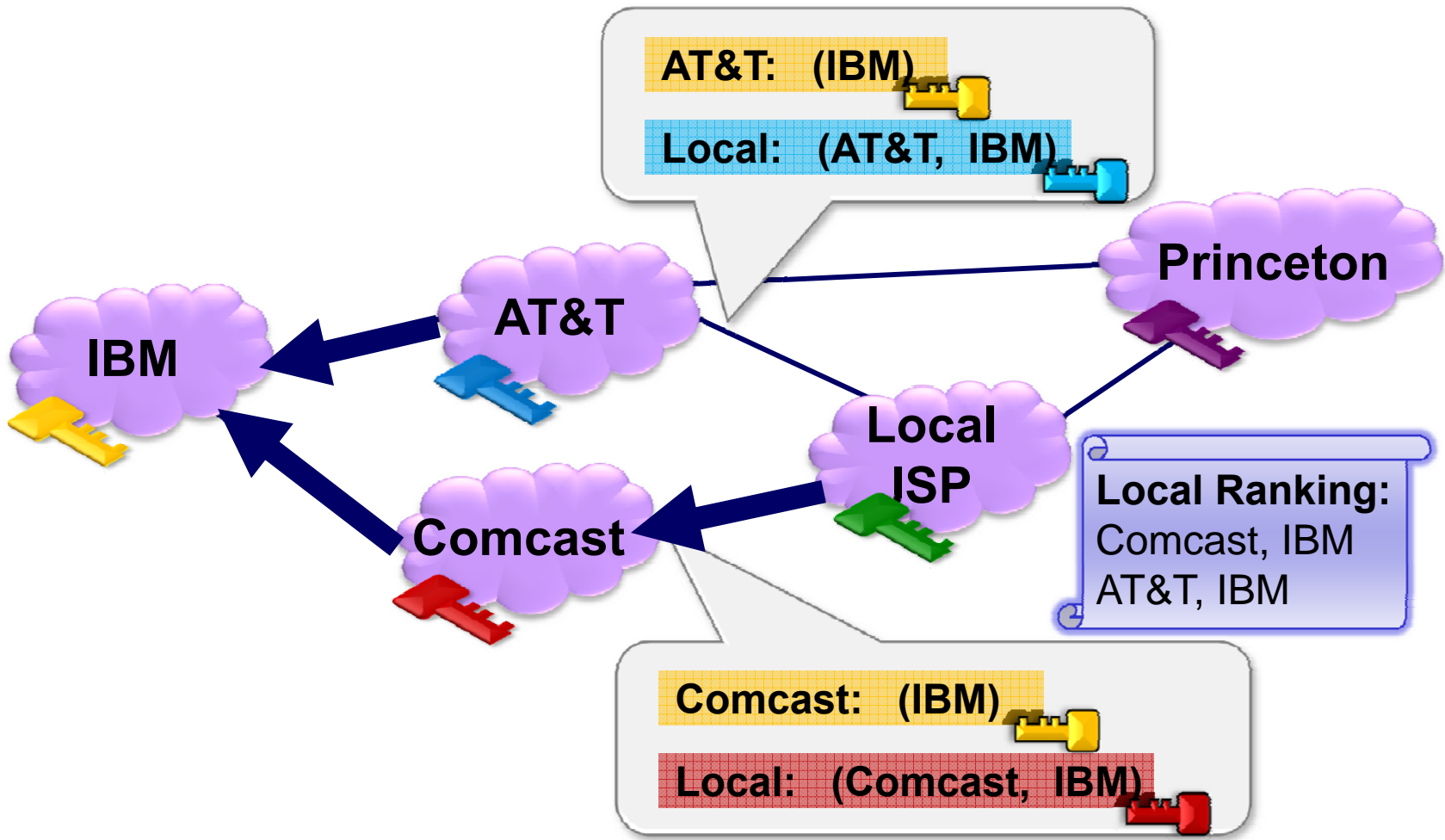
Assumes a public-key infrastructure that, today, we don't have.





# Secure BGP (1b)

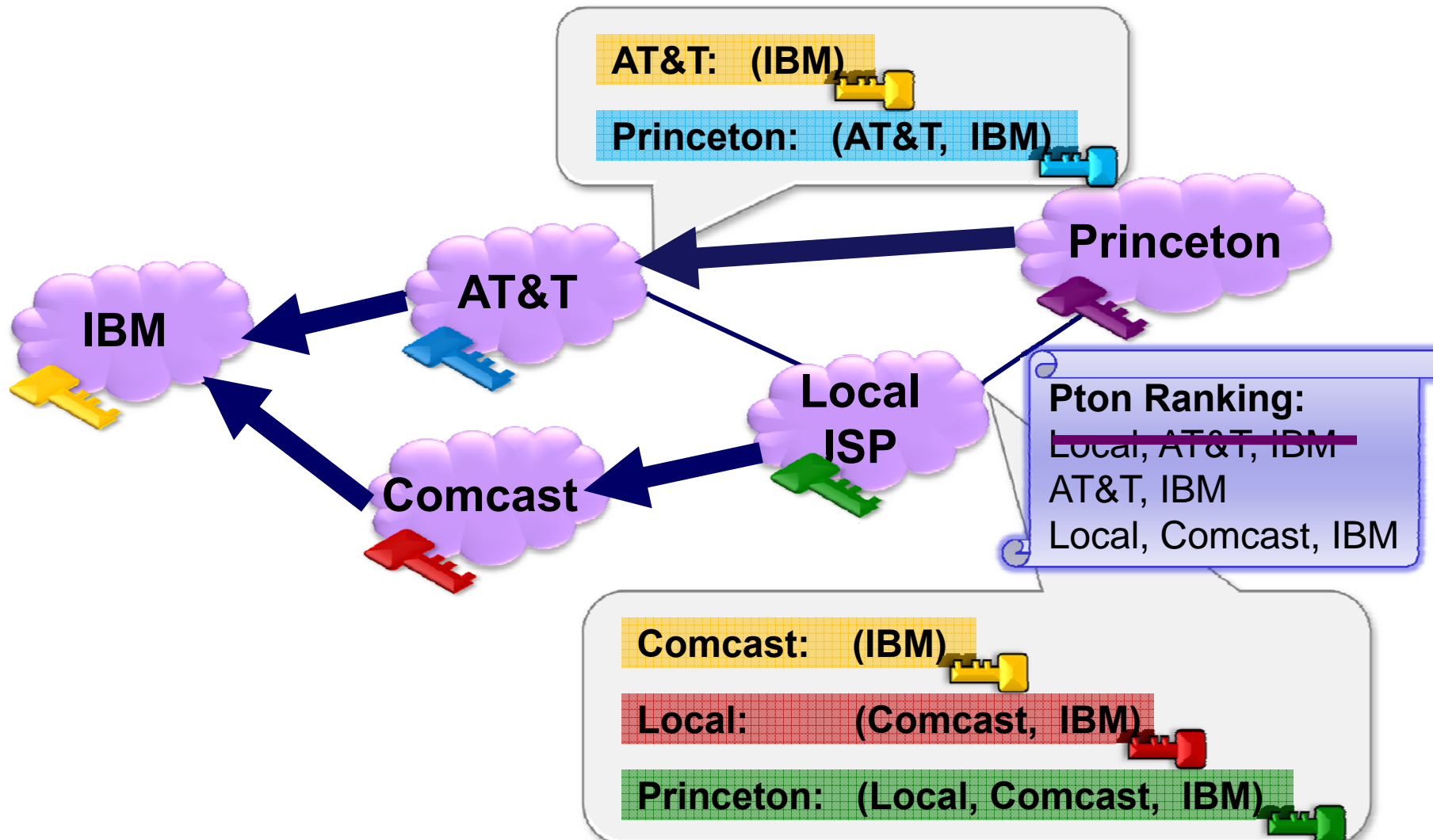
If **a** announced path **abP** then **b** announced **bP** to **a**





## Secure BGP (2)

If **a** announced path **abP** then **b** announced **bP** to **a**

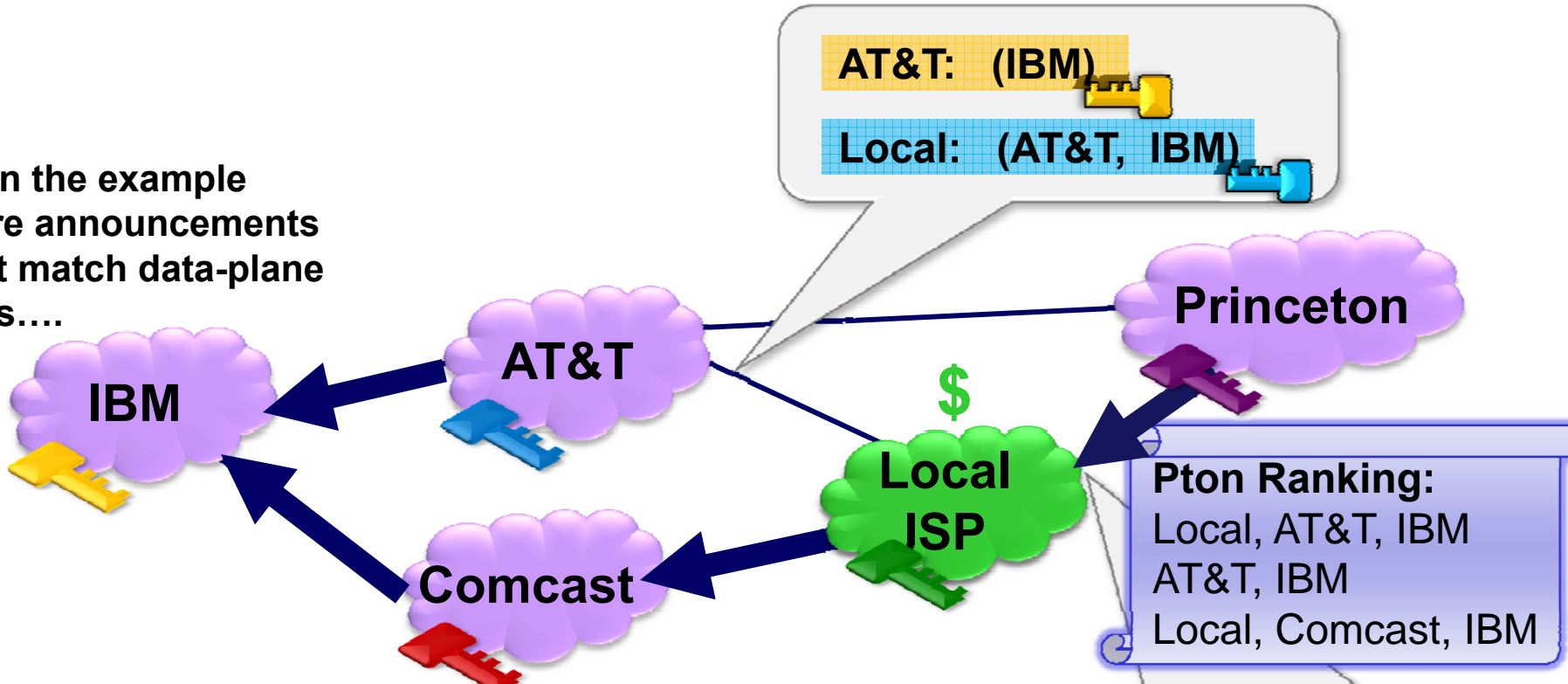




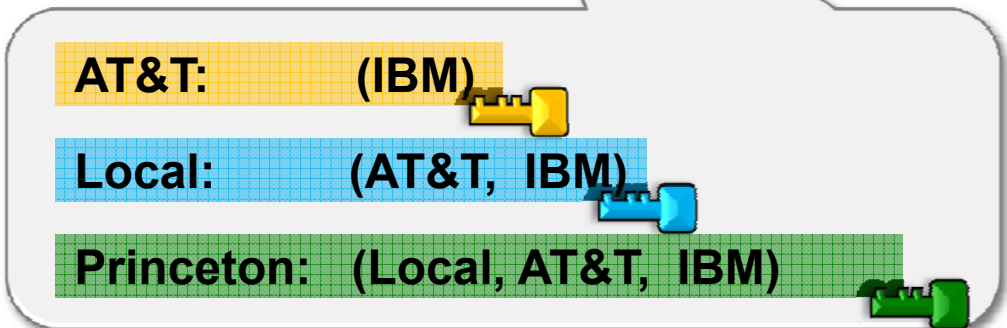
# Secure BGP : Matching the Data Plane ???

If **a** announced path **abP** then **b** announced **bP** to **a**

Again the example where announcements don't match data-plane paths....



Why does Local ISP do this?  
Let's look at utility models.



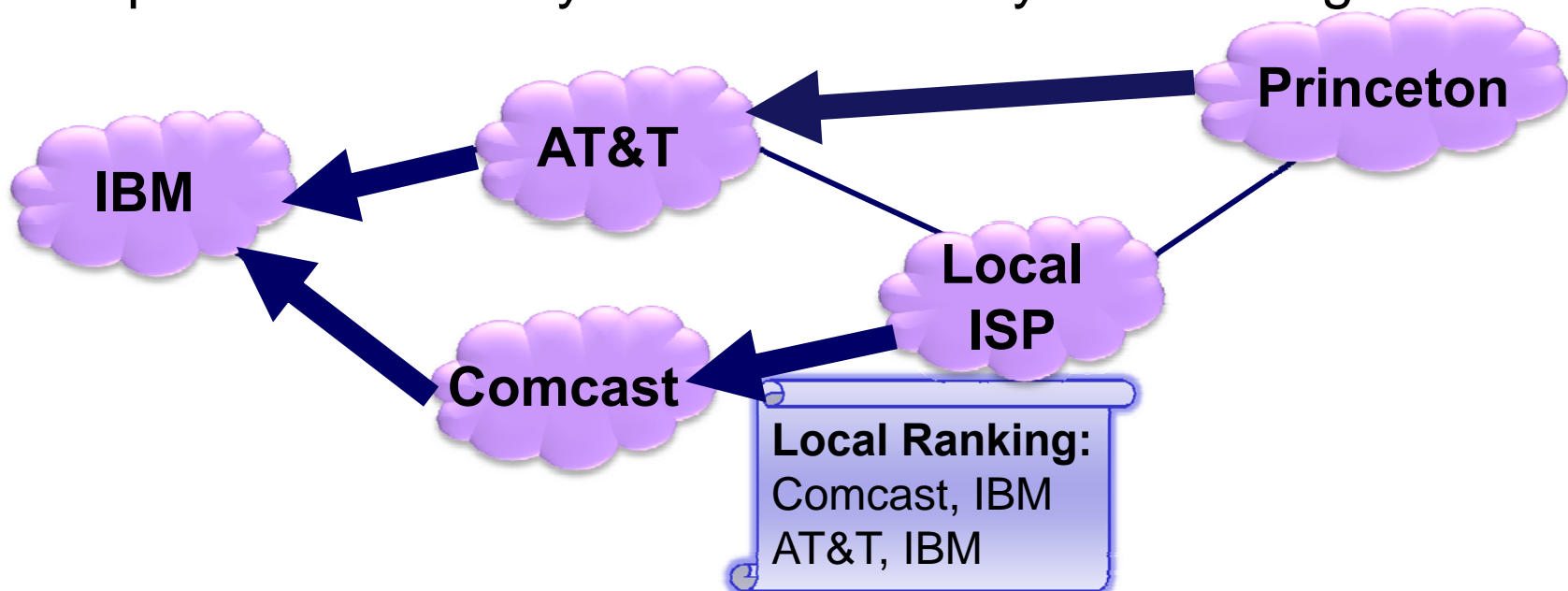


# Modeling Utility (1)

## Our model of utility:

$$\text{Utility of AS} = \text{Utility of outgoing (data-plane) path} + \text{Utility of attracted incoming traffic}$$

In all prior work: Utility is determined by the ranking function



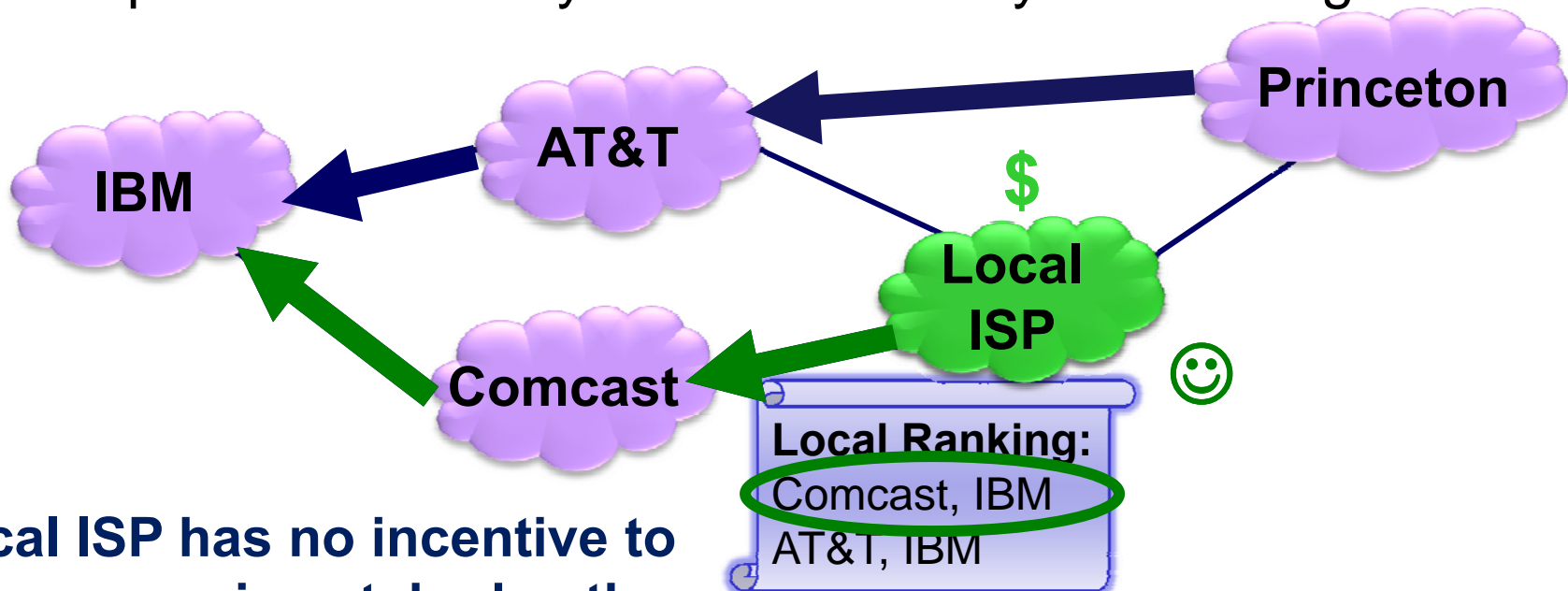


## Modeling Utility (2)

### Our model of utility:

$$\text{Utility of AS} = \text{Utility of outgoing (data-plane) path} + \text{Utility of attracted incoming traffic}$$

In all prior work: Utility is determined by the ranking function



**Local ISP has no incentive to announce mismatched paths.**



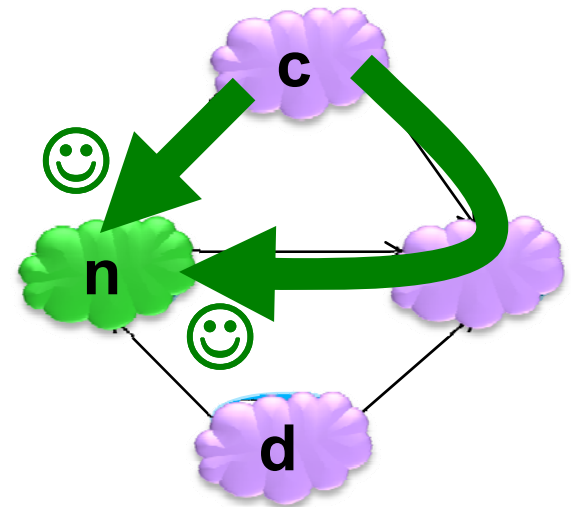
# Modeling Utility with Traffic Attraction (1)

## Our model of utility:

Utility of **AS** = Utility of outgoing (data-plane) path + Utility of attracted incoming traffic

## Traffic-volume attractions:

- **AS** only cares **who** originates traffic
- Models incentive to snoop / tamper
- ... or increase incoming traffic volumes





## Modeling Utility with Traffic Attraction (2)

### Our model of utility:

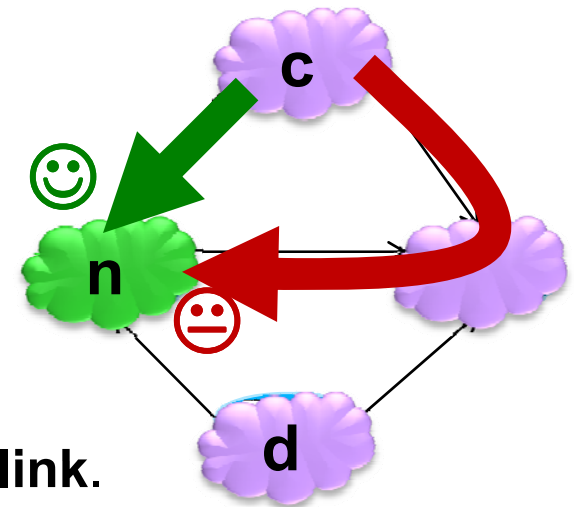
$$\text{Utility of AS} = \text{Utility of outgoing (data-plane) path} + \text{Utility of attracted incoming traffic}$$

### Traffic-volume attractions:

- **AS** only cares **who** originates traffic
- Models incentive to snoop / tamper
- ... or increase incoming traffic volumes

### Customer attractions:

- **AS** wants to attract traffic from customers **via direct link**.
- Models bilateral economic relationships.







## Modeling Utility with Traffic Attraction (3)

### Our model of utility:

$$\text{Utility of AS} = \text{Utility of outgoing (data-plane) path} + \text{Utility of attracted incoming traffic}$$

### Traffic-volume attractions:

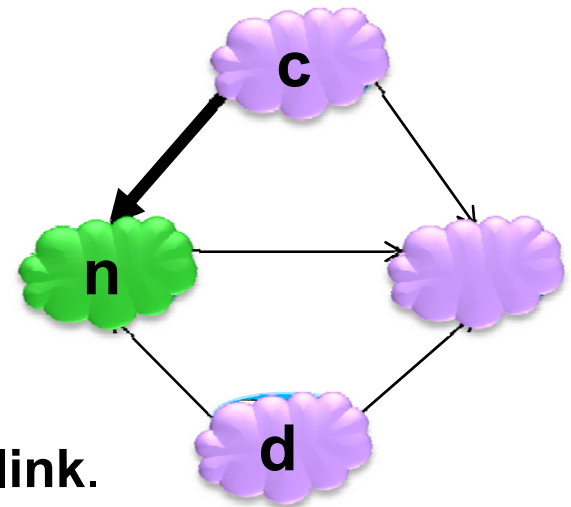
- **AS** only cares **who** originates traffic
- Models incentive to snoop / tamper
- ... or increase incoming traffic volumes

### Customer attractions:

- **AS** wants to attract traffic from customers **via direct link**.
- Models bilateral economic relationships.

### Generic attractions:

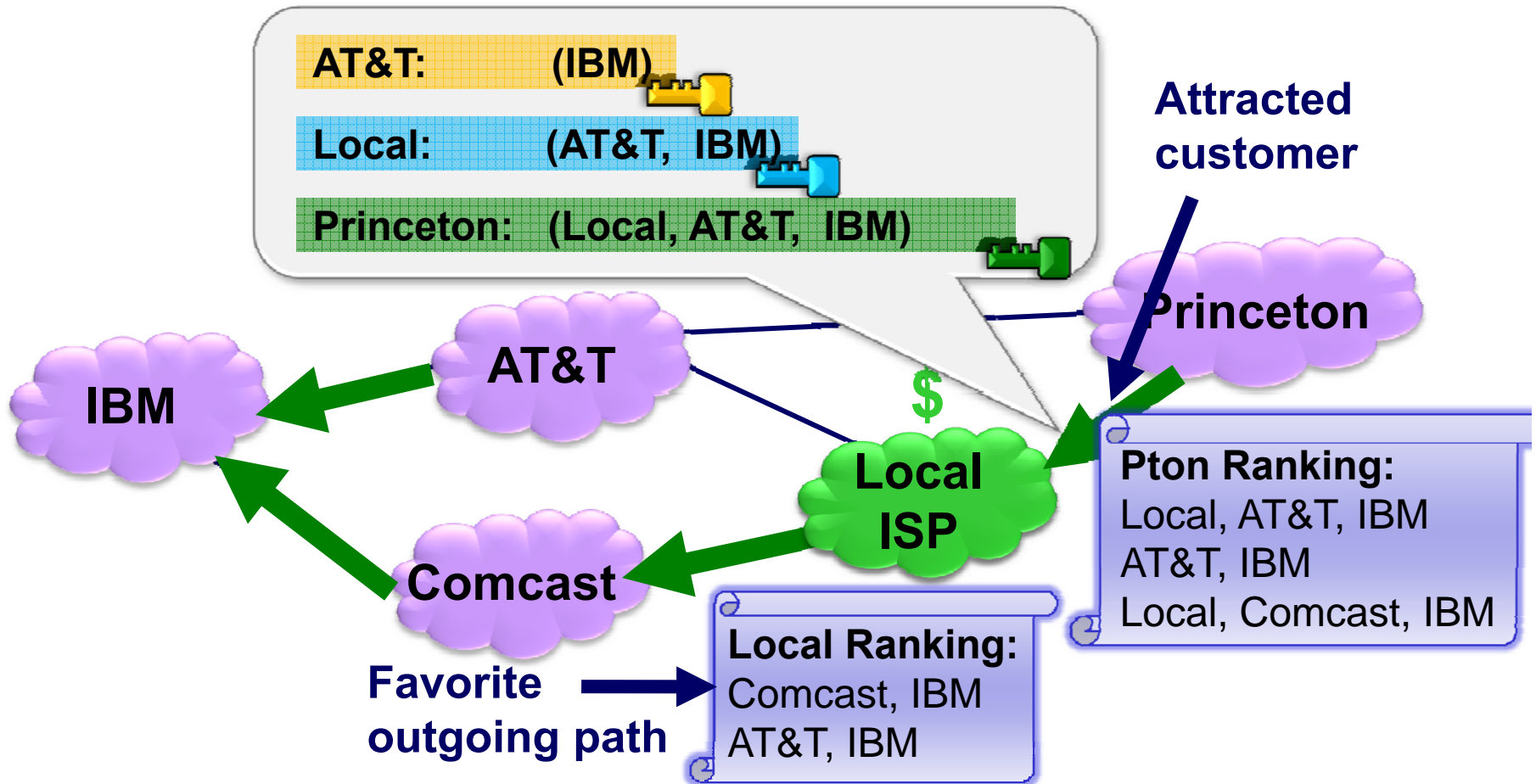
- **AS** wants to attract traffic from specific ASes via a specific path





# Result: Secure BGP is not Sufficient! (1)

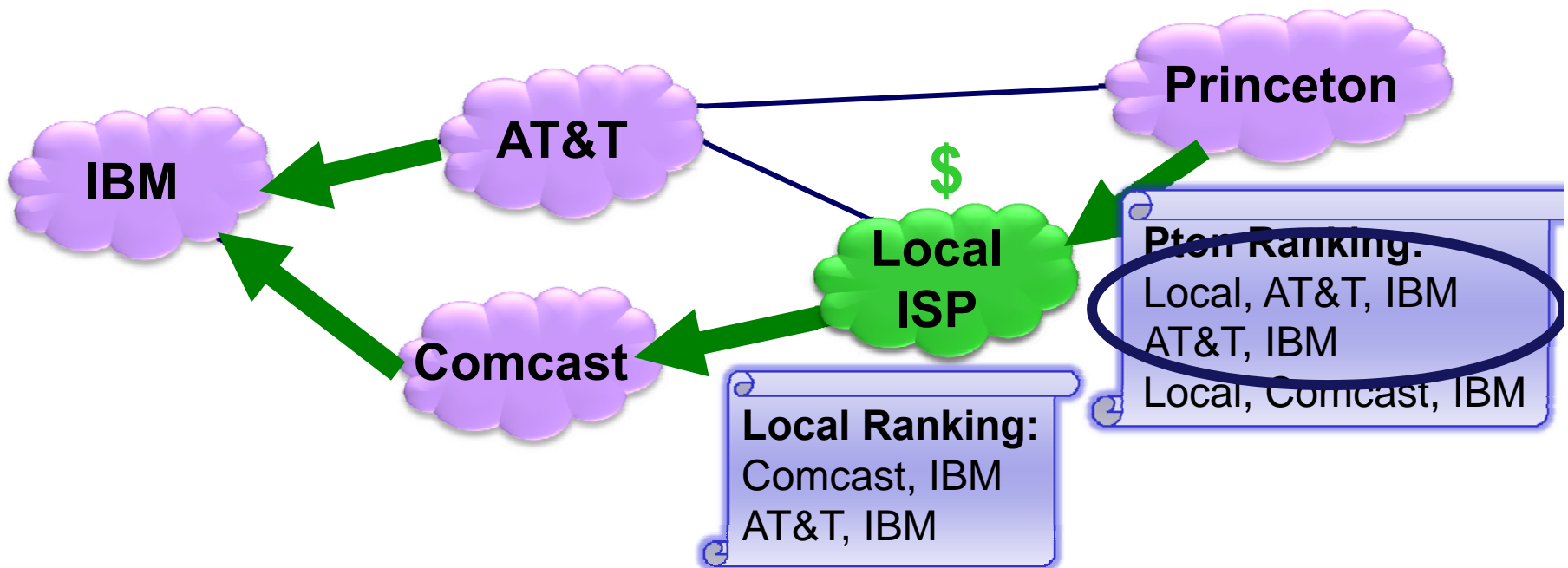
With **traffic-volume** OR **customer** attractions, there can be an incentive to announce mismatched paths, **even with Secure BGP**.





## Result: Secure BGP is not Sufficient! (2)

With **traffic-volume** OR **customer** attractions, there can be an incentive to announce mismatched paths, **even with Secure BGP**.

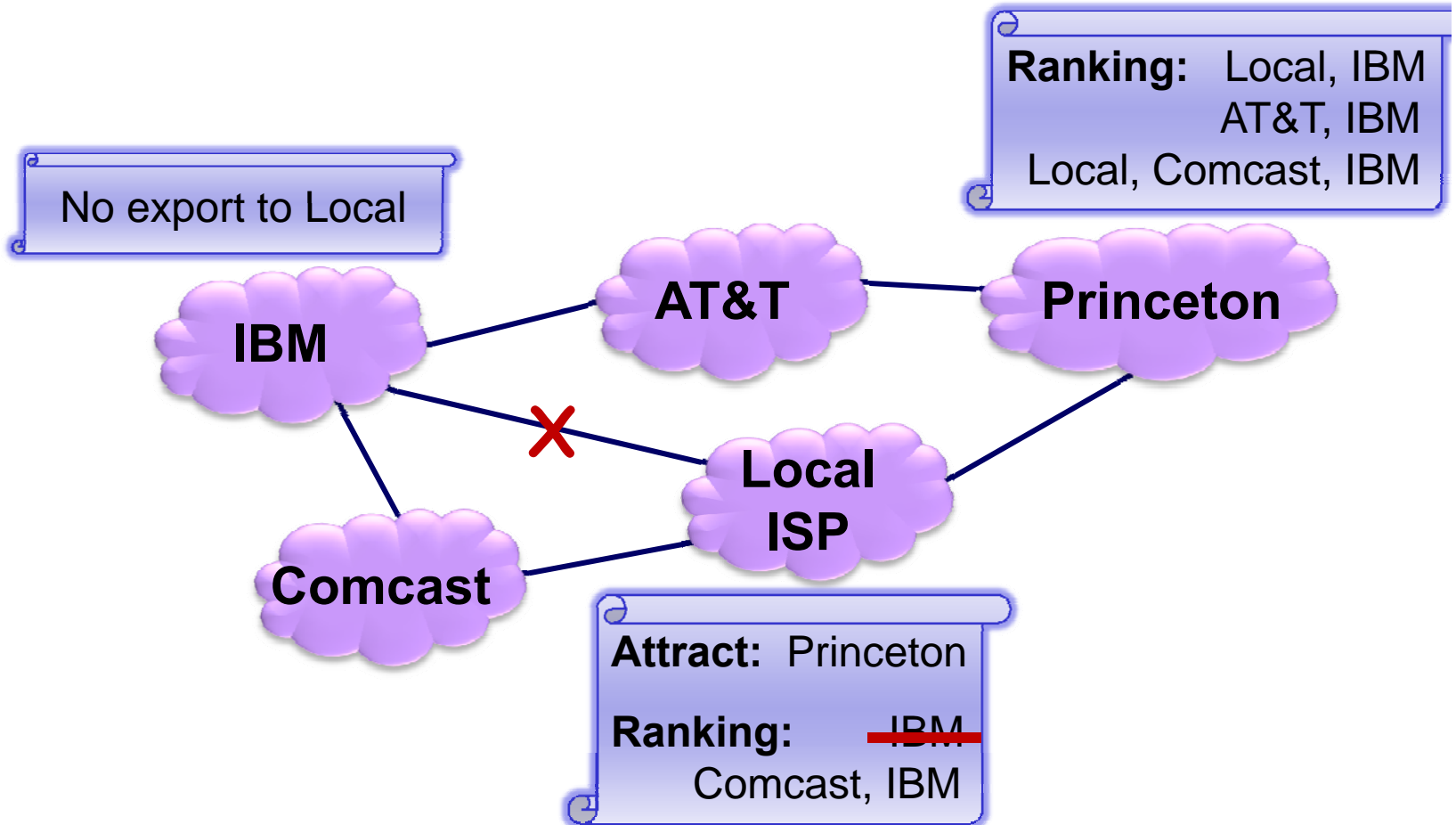


**Observation:** Princeton does not use a shortest-path policy. 19/23



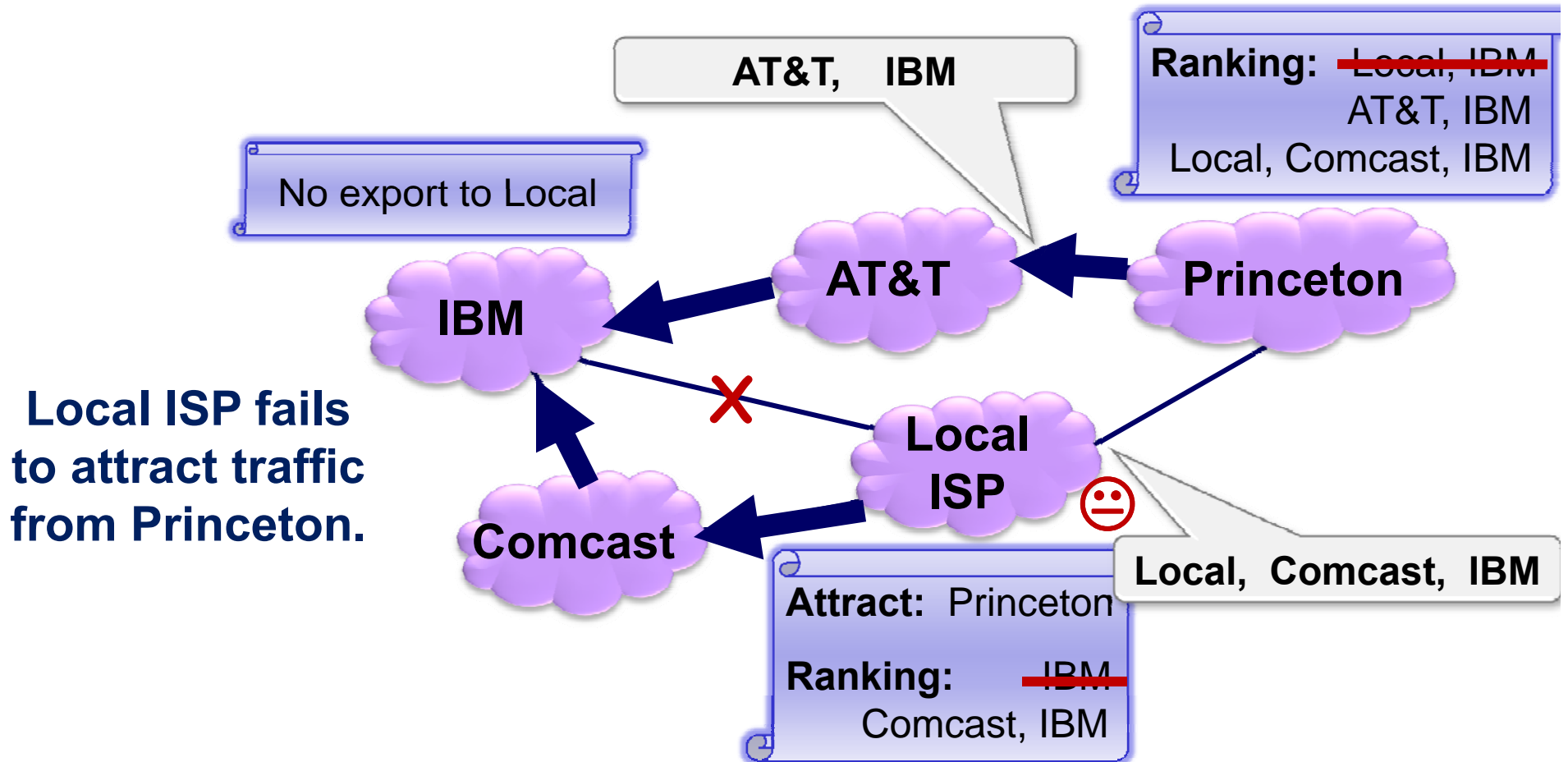
# Result: Shortest-Path Policy is not Sufficient! (1)

With **traffic-volume** OR **customer** attractions, there can be an incentive to mismatch paths, **even with shortest-path policies.**



# Result: Shortest-Path Policy is not Sufficient! (2)

With **traffic-volume** OR **customer** attractions, there can be an incentive to mismatch paths, **even with shortest-path policies.**

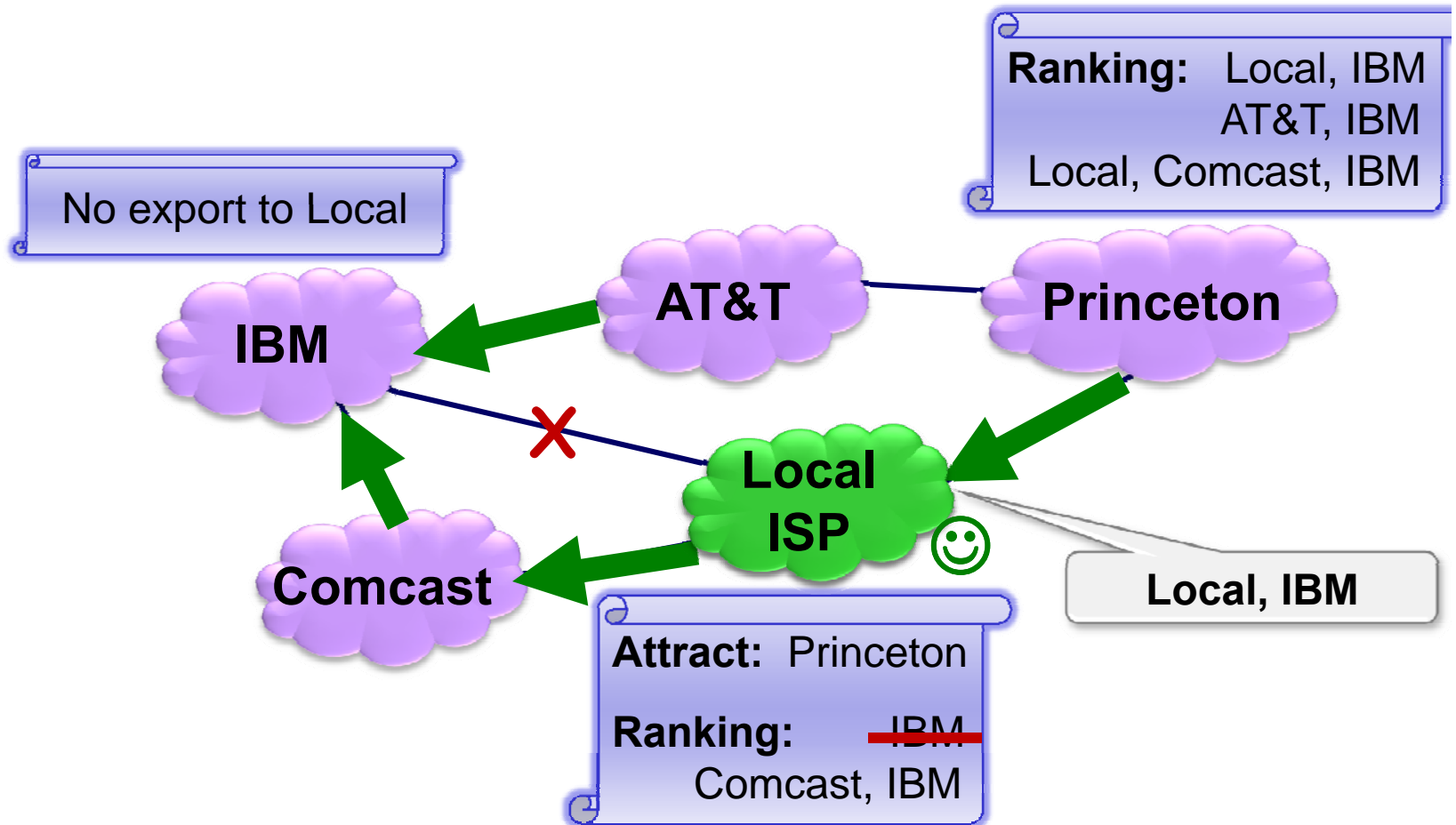


Local ISP fails to attract traffic from Princeton.



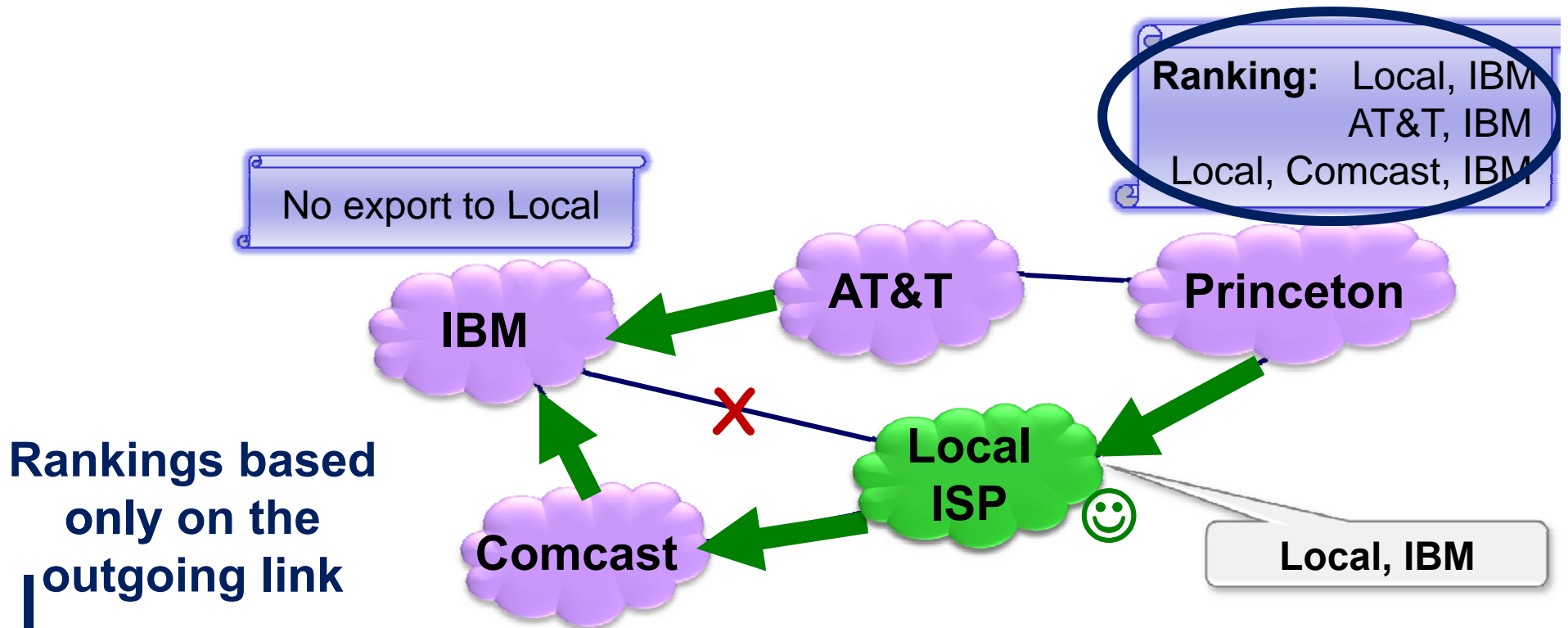
## Result: Shortest-Path Policy is not Sufficient! (3)

With **traffic-volume** OR **customer** attractions, there can be an incentive to mismatch paths, **even with shortest-path policies.**



# Result: Shortest-Path Policy is not Sufficient! (4)

With **traffic-volume** OR **customer** attractions, there can be an incentive to mismatch paths, **even with shortest-path policies.**



**Observation:** Manipulation not possible with Secure BGP.

**Observation:** Princeton does not use a next-hop policy.



# Traffic Volume Attractions

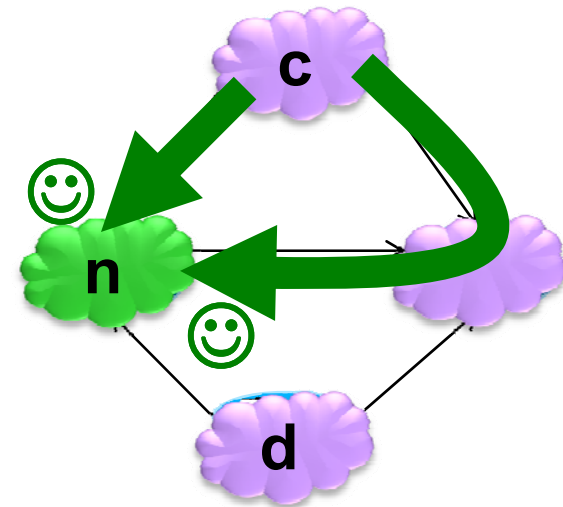
When all attractions are **traffic volume**, nodes have no **incentive** to **unilaterally** announce mismatched paths if all nodes in the network use either:

1. **Secure BGP**, and
2. **Policy consistency**;

OR

1. **Next-hop policies**;

and there is no dispute wheel in the network and there is consistent export (in the first case) or all-or-nothing export (in the second case).







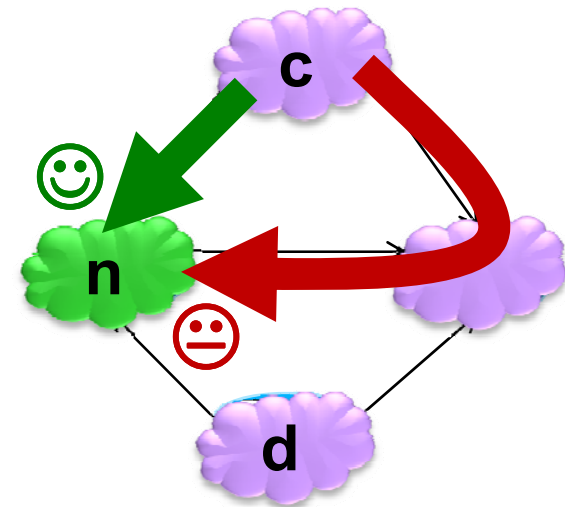
# What about Customer Attractions? (1)

---

Are these sufficient if we have customer attractions?

- X** 1. Secure BGP, and
- X** 2. Policy consistency;

- X** 1. Next-hop policies;



**Customer attractions:**  
Attract customers **via**  
**direct link.**



## What about Customer Attractions? (2)

---

**Are these sufficient if we have customer attractions?**

- 1. Secure BGP, and**
- 2. Next-hop policies;**





## Somewhat More Formally ...

With **generic traffic attraction**, there exists an honest strategy that obtains the best possible stable outcome for each node (*i.e.* that each node has no incentive to **unilaterally** mismatch paths), if every node uses

1. **Secure BGP**, and
2. **Next-hop policies**

**and** there is no dispute wheel in the network  
**and** every node uses all-or nothing export.

Removing any condition gives a counterexample 😞



# Conclusions

---

What conditions ensure BGP messages match data-plane paths?

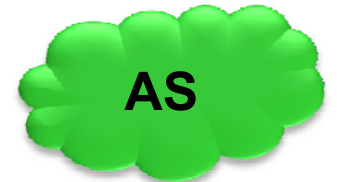
- **Secure BGP** is not sufficient
- ...if it is rational for ASes to want to **attract traffic**.
- Generally, we need **next-hop policy** as well as
- ... other conditions (no dispute wheel, no egress filtering).



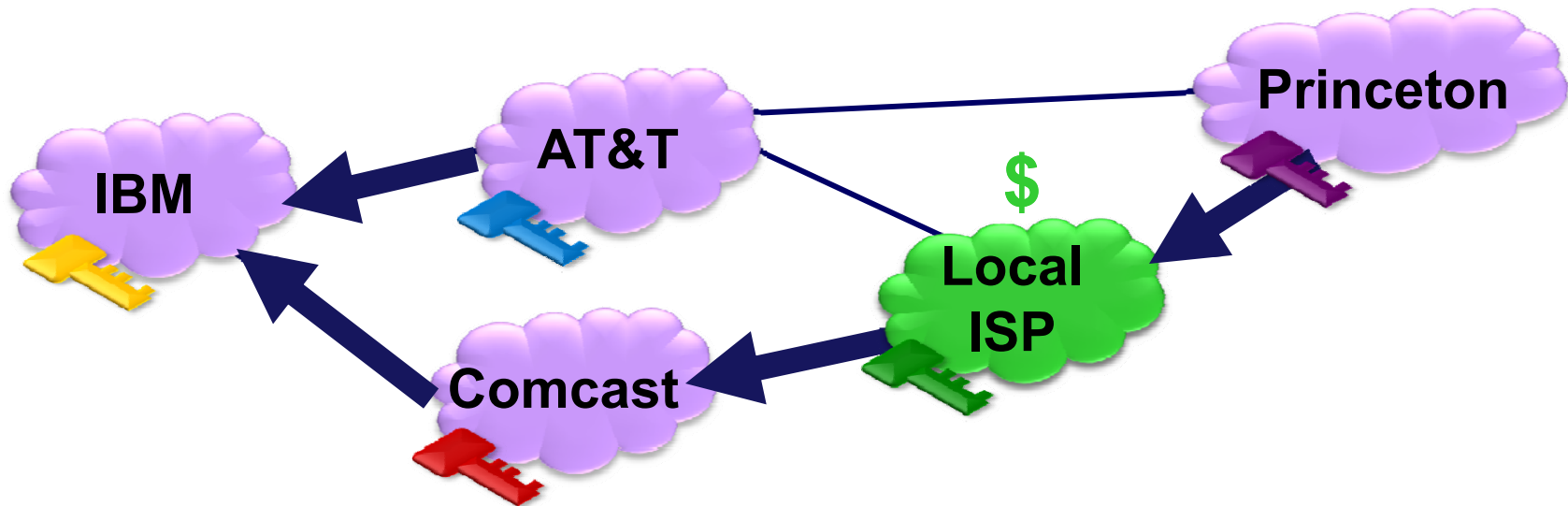
Also, notice how strongly results depend on utility model.

What should we do?

- Use expensive **data-plane protocols**?
- **Forget about** matching BGP messages to data plane?
- Allow ASes to send traffic on **more than one outgoing link**?



# Thanks!



Full version with all proofs and counterexamples available:

[www.princeton.edu/~goldbe/](http://www.princeton.edu/~goldbe/)

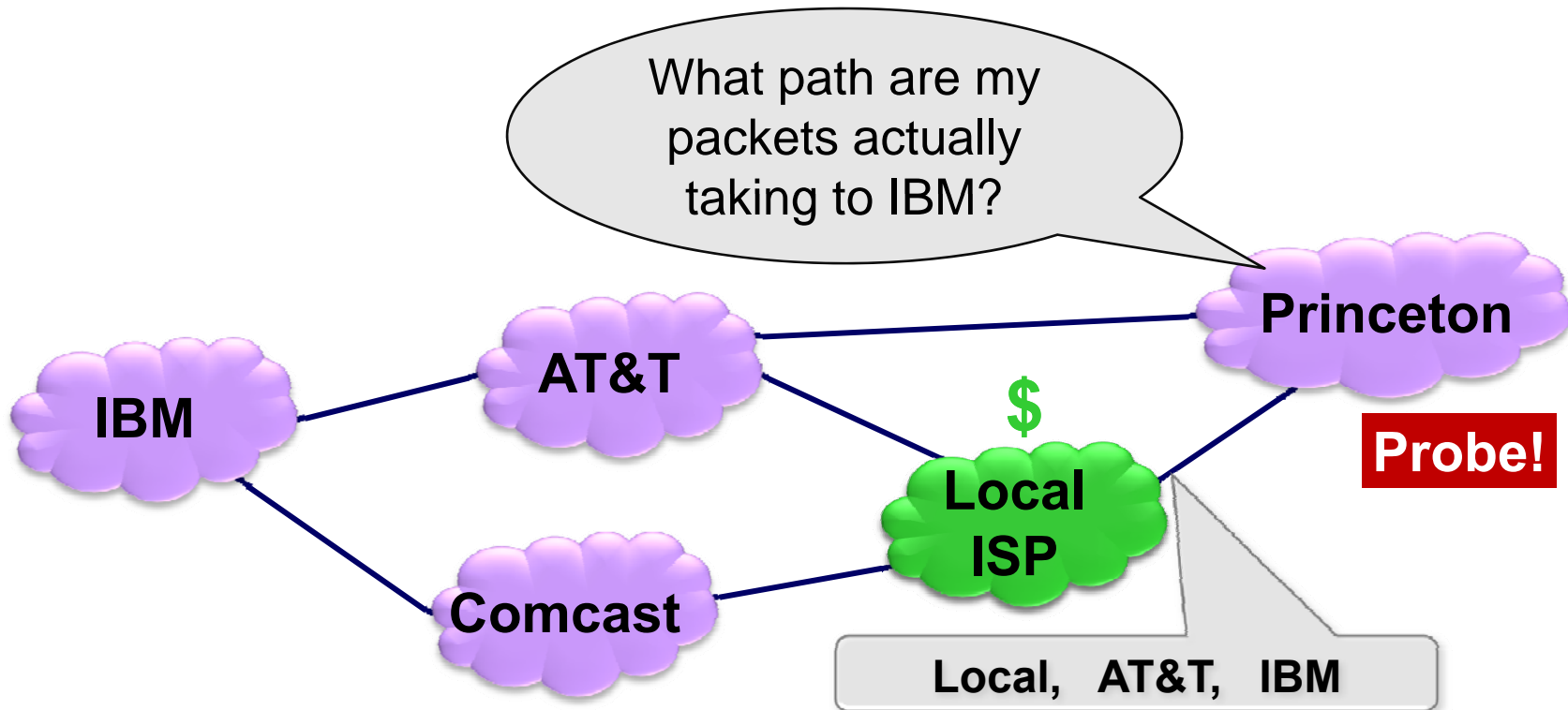




# Matching Data Plane: Related Approaches (Backup)

## Secure Data-Plane Protocols:

- Packet Passports [LYWA-06]      Packet Obituaries [AMISS-07]  
Truth in advertising [WBAGS-07]      Failure Localization [BGX-08]
- ✗ **Secure AS-path tracing protocols incur overheads proportional to the amount of traffic sent in the data plane.**





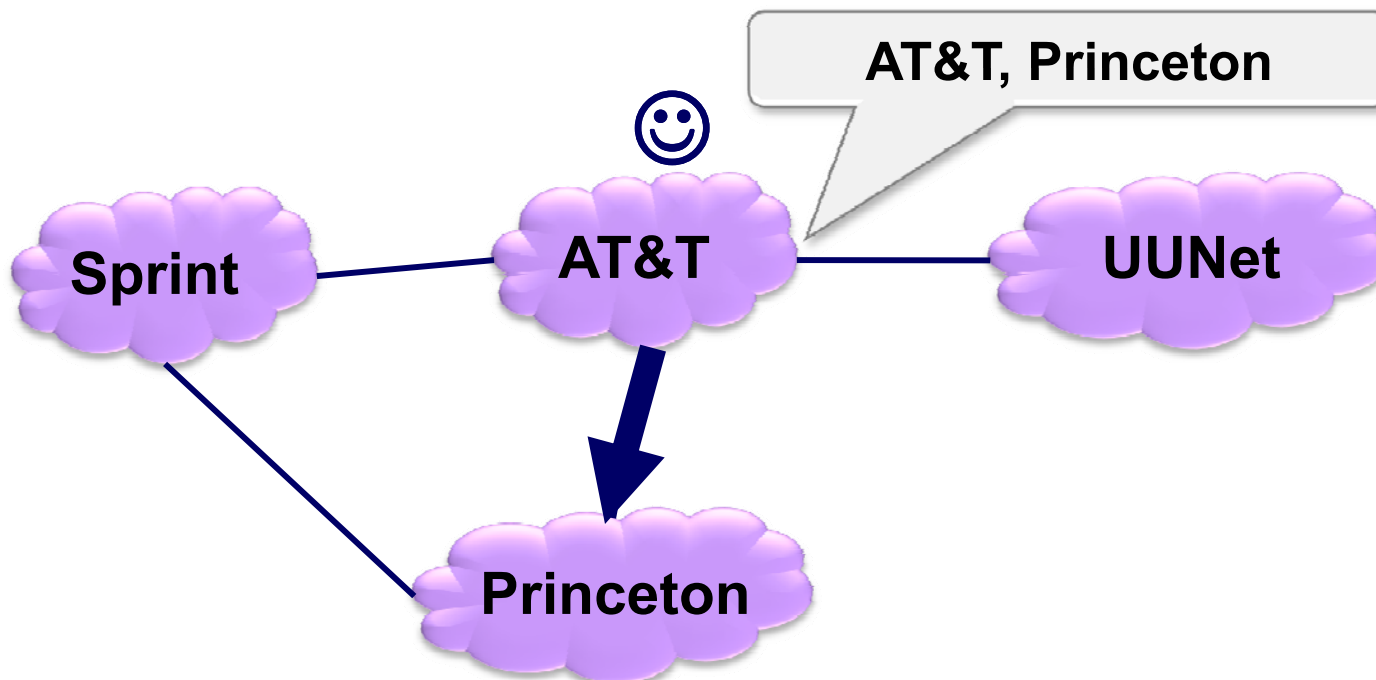
## All-or-Nothing Export (1)

For each neighbor, either export all paths or export none.

**Path-based egress filtering is not allowed!**

**(Incompatible with practice.)**

**IF**



**AT&T makes money because it delivers traffic to a customer.**



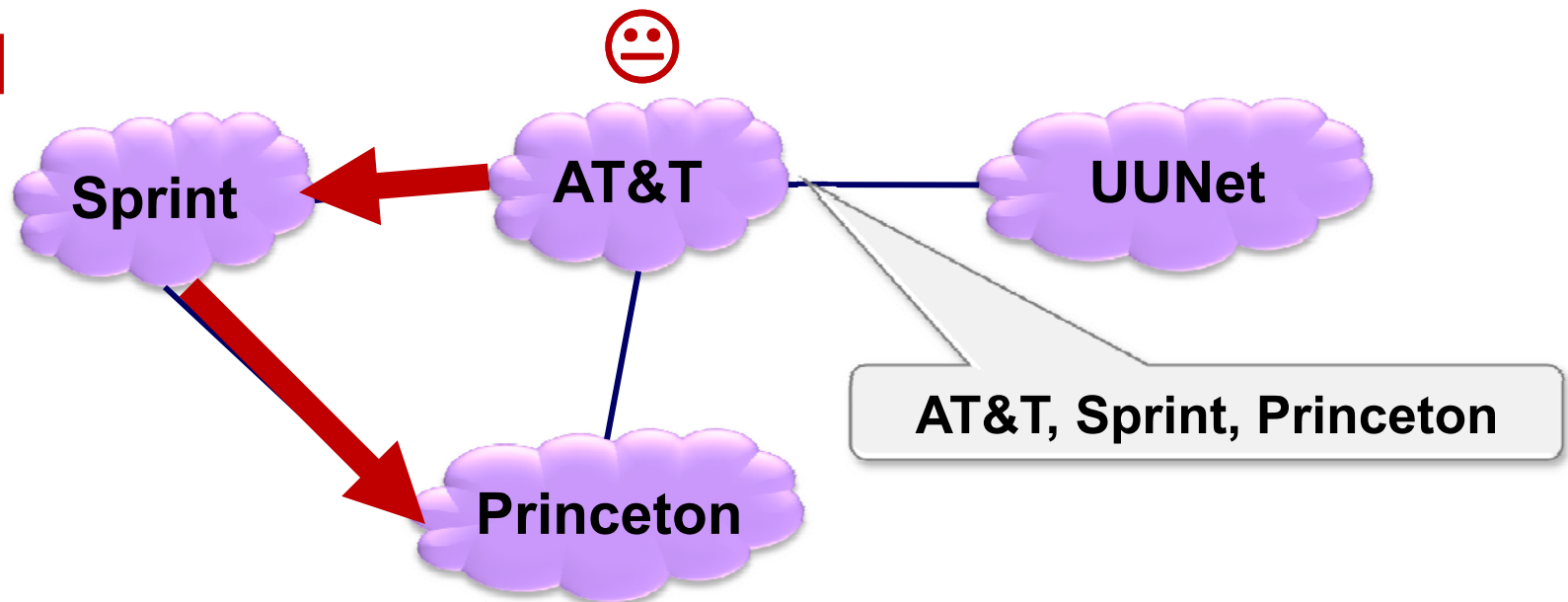
## All-or-Nothing Export (2)

For each neighbor, either export all paths or export none.

**Path-based egress filtering is not allowed!**

**(Incompatible with practice.)**

**THEN**



**AT&T loses money because it transits traffic between its peers.**



# Formalizing the Model

Acts on  
data plane

Acts on  
control plane

## Utility

Satisfaction of node  $n$  with  
a data-plane routing outcome  $T$

$$u_n(T) = v_n(T) + \alpha_n(T)$$

$v_n(T)$  is the **valuation function**  
Satisfaction of  $n$  is with his  
**outgoing path** in  $T$

$\alpha_n(T)$  is the **attraction function**  
Satisfaction of  $n$  with  
**incoming traffic** in  $T$

## Ranking

Ranking of **outgoing paths**  
Used by  $n$  in BGP decision process

$$r_n(T)$$

## Export

The set of neighbours to which  $n$  is  
willing to announce path  $P$

$$e_n(T)$$

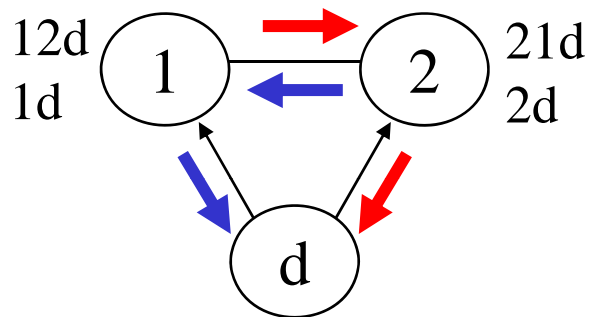
compile

Formal  
model

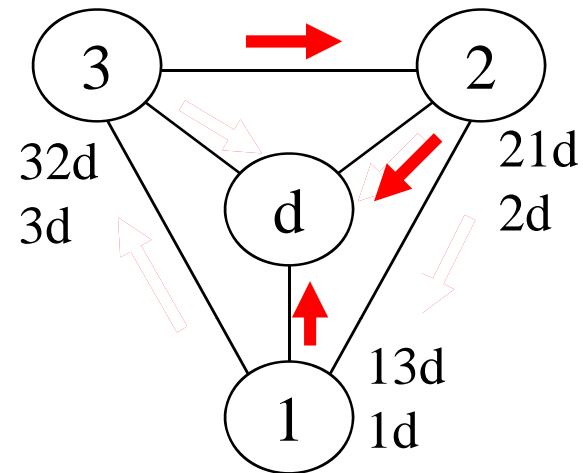


## Stability: No Dispute Wheel

A dispute wheel is a cycle of nodes with rankings that prefer paths through neighbours over direct paths



**Disagree:** 2 stable outcomes



**Bad Gadget:** no stable outcomes