CS 512, Spring 2018, Handout 01 Motivating Examples

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example: code fragment written in GCL

taken from Example 18 in *Properties of Transition Systems* (click to retrieve):

```
. . .
```

```
test1 := false; test2 := false;

do true \rightarrow test1 := true; \cdots # process A

| test1 \rightarrow test2 := true; \cdots # process B

| test2 \rightarrow test2 := false; \cdots # process C

od

...
```

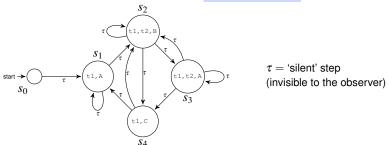
example: code fragment written in GCL

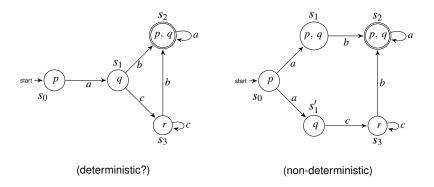
taken from Example 18 in Properties of Transition Systems (click to retrieve):

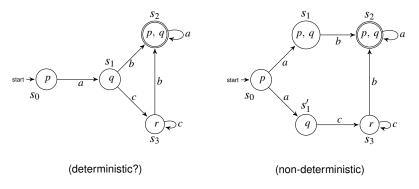
```
. . .
```

```
test1 := false; test2 := false;
do true → test1 := true; ··· # process A
| test1 → test2 := true; ··· # process B
| test2 → test2 := false; ··· # process C
od
...
```

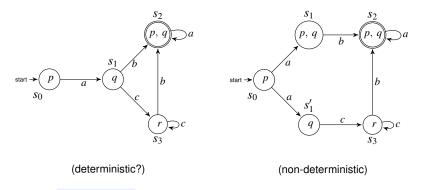
which can be represented by the following **transition system** (actually, its graph):



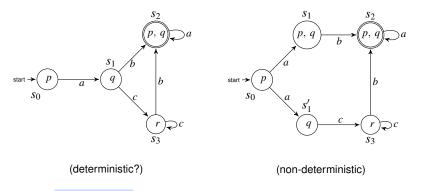




▶ as finite automata , they are equivalent and accept/recognize the set of strings/words defined by the regular expression $a b a^* + a c c^* b a^*$.



- ▶ as finite automata , they are equivalent and *accept/recognize* the set of *strings/words* defined by the regular expression $ab a^* + a c c^* b a^*$.
- as transition systems, they are not equivalent, because the **observable** propositional atoms in state s_1 and s_1' are different.



- ▶ as finite automata , they are equivalent and *accept/recognize* the set of *strings/words* defined by the regular expression $ab a^* + a c c^* b a^*$.
- as transition systems, they are not equivalent, because the **observable** propositional atoms in state s_1 and s'_1 are different.
- as ω-automata, they are
 (stick around, if you want to hear the story later in the semester)

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