

CS 512, Spring 2018, Handout 22

Counterexamples
and
Probabilistic Model Checking

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April 17, 2018

Counterexamples – in general

(material in this and later slides mostly due to Prof. J-P Katoen of Aachen Univ)

- ▶ Reminder: **model checking** = **bug hunting** ,
bugs are discovered by **counterexamples**,
states that refute a given property (desirable or harmful).
- ▶ **Counterexamples** are (formally expressed) instances of system behavior that contradict a system's (formally expressed) specification.

Counterexamples – in general

- ▶ **Counterexamples in LTL** are typically finite execution paths:
 - ▶ To contradict $(\Box \varphi)$,
we want a finite path ending in a $(\neg\varphi)$ -state.
 - ▶ To contradict $(\Diamond \varphi)$,
we want a finite $(\neg\varphi)$ -path leading to a $(\neg\varphi)$ -cycle.

Methods of LTL model-checkers incorporate forms of **breadth-first search** for generating shortest counterexamples (e.g., see Handout 13).

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Methods of LTL model-checkers incorporate forms of **breadth-first search** for generating shortest counterexamples (*e.g.*, see Handout 13).

- ▶ **Counterexamples in CTL** are typically finite trees of execution paths:
 - ▶ To contradict universal CTL,
we want **all** paths in a tree of execution paths.
 - ▶ To contradict existential CTL,
we want **one** path in a tree of execution paths.

Methods of CTL model-checkers also incorporate some form of **breadth-first search**, combined with more advanced data structures.

Counterexamples – in PCTL (Probabilistic CTL)

► **Problem statement:**

Given a WFF of PCTL of the form $\mathbb{P}_{\leq p}(\varphi)$

– for example, in shorthand, $(p \Updownarrow^{\leq 1/2} q)$ or $(\bigcirc^{\leq 2/3} p)$ –
together with a Markov chain \mathcal{M} and a state s in \mathcal{M} ,
we want to decide whether:

$$\mathcal{M}, s \not\models \mathbb{P}_{\leq p}(\varphi) \quad \text{or, more succinctly,} \quad s \not\models \mathbb{P}_{\leq p}(\varphi)$$

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we want to decide whether:

$$\mathcal{M}, s \not\models \mathbb{P}_{\leq p}(\varphi) \quad \text{or, more succinctly,} \quad s \not\models \mathbb{P}_{\leq p}(\varphi)$$

► **A counterexample C for $\mathbb{P}_{\leq p}(\varphi)$ at state s in \mathcal{M}**

is a set of finite paths (or **evidences**) in \mathcal{M} satisfying:

- if $\pi \in C$, then π starts at s and $\pi \models \varphi$, and
- $\Pr(C) > p$ where $\Pr(C) \triangleq \sum_{\pi \in C} \Pr(\pi)$,
i.e., the sum of the probabilities of the paths in C , exceeds p .

If $\Pr(C) > p$, we conclude that $s \not\models \mathbb{P}_{\leq p}(\varphi)$.

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- In this handout, we limit attention to **discrete-time** Markov chains –
we delay work done on **continuous-time** Markov chains till next year (!).

Counterexamples – in PCTL (Probabilistic CTL)

- ▶ A counterexample C for $\mathbb{P}_{\leq p}(\varphi)$ is **minimal** if $|C| \leq |C'|$ for any counterexample C' for $\mathbb{P}_{\leq p}(\varphi)$.

Counterexamples – in PCTL (Probabilistic CTL)

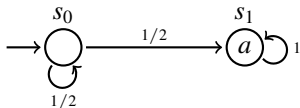
- ▶ A counterexample C for $\mathbb{P}_{\leq p}(\varphi)$ is **minimal** if $|C| \leq |C'|$ for any counterexample C' for $\mathbb{P}_{\leq p}(\varphi)$.
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- ▶ **Fact:** Counterexamples for non-strict probability bounds (*i.e.*, bounds of the form “ $\leq p$ ”, not “ $< p$ ”) are **finite**.

Counterexamples – in PCTL (Probabilistic CTL)

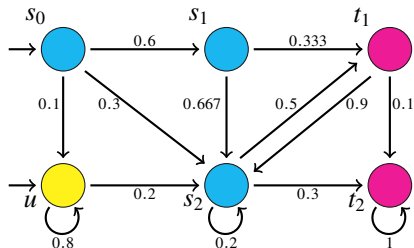
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- ▶ **Fact:** Counterexamples for non-strict probability bounds (*i.e.*, bounds of the form “ $\leq p$ ”, not “ $< p$ ”) are **finite**.
- ▶ **Infinite** counterexamples may be needed for WFF's with strict probability bounds.
- ▶ For example, an **infinite** counterexample is needed for $s_0 \not\models \mathbb{P}_{< 1}(\diamond a)$, *i.e.*, for $s_0 \not\models (\diamond^{< 1} a)$ in the following Markov chain:



Example showing how to handle “until” WFF’s in PCTL¹

¹ Partly inspired by Example 10.41 in [PMC, page 786].

Example showing how to handle “until” WFF’s in PCTL¹



Wanted:

counterexamples for $s_0 \not\models (\varphi \cup^{\leq 1/2} \psi)$

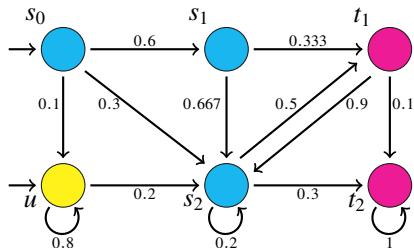
blue states : only prop WFF φ holds,

red states : only prop WFF ψ holds,

yellow states : neither φ nor ψ hold.

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Example showing how to handle “until” WFF’s in PCTL¹



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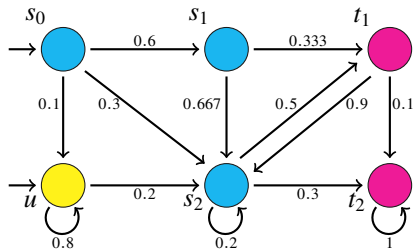
Wanted:

counterexamples for $s_0 \not\models (\varphi \text{ U}^{\leq 1/2} \psi)$

evidence	probability
$\pi_1 \triangleq s_0 s_1 t_1$	0.2
$\pi_2 \triangleq s_0 s_1 s_2 t_1$	0.2
$\pi_3 \triangleq s_0 s_2 t_1$	0.15
$\pi_4 \triangleq s_0 s_1 s_2 t_2$	0.12
$\pi_5 \triangleq s_0 s_2 t_2$	0.09
...	...

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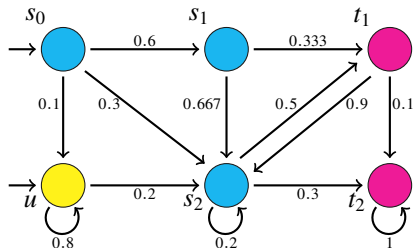
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	counterexample	cardinality	probability
	$\{\pi_1, \pi_2, \pi_3, \pi_4, \pi_5\}$	5	0.76
	$\{\pi_1, \pi_3, \pi_4, \pi_5\}$	4	0.56
	$\{\pi_2, \pi_3, \pi_4, \pi_5\}$	4	0.76
minimal →	$\{\pi_1, \pi_2, \pi_4\}$	3	0.52
minimal →	$\{\pi_1, \pi_2, \pi_3\}$	3	0.55

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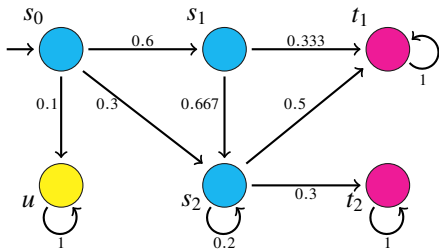
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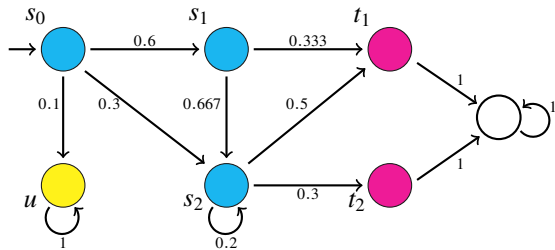
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Obtaining smallest counterexamples



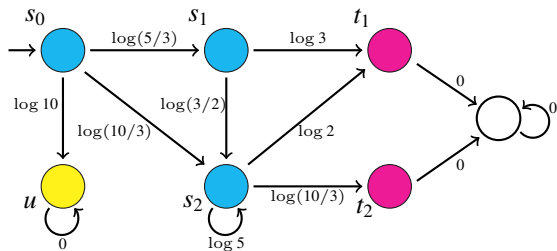
Step 1: Make all ψ -states and all $(\neg\varphi \wedge \neg\psi)$ -states absorbing, which requires eliminating some transitions (e.g., the transitions out of t_1 and u) and making the transition probability = 1 on all self-loops.

Adapting a bit more



Step 2: Insert a sink state and redirect all outgoing edges of ψ -states to it.

A weighted digraph



Step 3: Turn the Markov chain into a weighted digraph (directed graph), where:

$$w(s, s') \triangleq \log \left(\frac{1}{\Pr(s, s')} \right)$$

for every pair of nodes/states s and s' . The logarithm can be base 10, or base e , or base 2 – it does not matter which base we choose.

A simple derivation

Given a finite path $\pi \triangleq s_0 s_1 s_2 \cdots s_n$:

$$\begin{aligned}w(\pi) &= w(s_0, s_1) + w(s_1, s_2) + \cdots + w(s_{n-1}, s_n) \\&= \log \left(\frac{1}{\Pr(s_0, s_1)} \right) + \log \left(\frac{1}{\Pr(s_1, s_2)} \right) + \cdots + \log \left(\frac{1}{\Pr(s_{n-1}, s_n)} \right) \\&= \log \left(\frac{1}{\Pr(s_0, s_1) \cdot \Pr(s_1, s_2) \cdot \cdots \cdot \Pr(s_{n-1}, s_n)} \right) \\&= \log \left(\frac{1}{\Pr(\pi)} \right)\end{aligned}$$

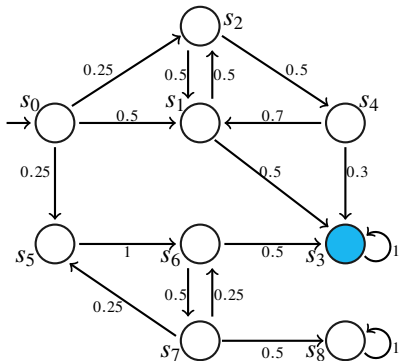
Conclusion 1: For all finite paths π and π' in the Markov chain, we have:

$$\underbrace{\Pr(\pi) \geq \Pr(\pi')}_{\text{in the Markov chain}} \quad \text{if and only if} \quad \underbrace{w(\pi) \leq w(\pi')}_{\text{in the weighted digraph}}$$

Conclusion 2: Finding a **strongest evidence** in the Markov chain is translated to a **shortest path problem** in the weighted digraph.

Another example: How to handle reachability properties

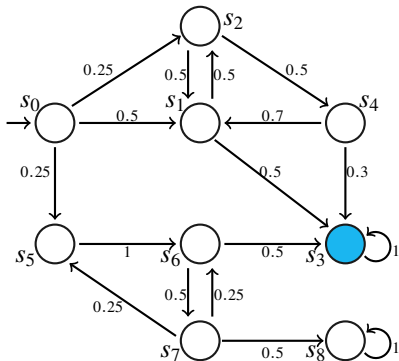
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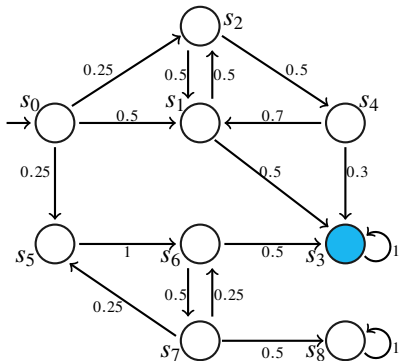
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Approach 1, based on using the transition (right-stochastic) 9×9 matrix A :

$$A = \begin{matrix} & \begin{matrix} s_0 & s_1 & s_2 & s_3 & s_4 & s_5 & s_6 & s_7 & s_8 \end{matrix} \\ \begin{matrix} s_0 \\ s_1 \\ s_2 \\ s_3 \\ s_4 \\ s_5 \\ s_6 \\ s_7 \\ s_8 \end{matrix} & \begin{bmatrix} 0 & .5 & .25 & 0 & 0 & .25 & 0 & 0 & 0 \\ 0 & 0 & .5 & .5 & 0 & 0 & 0 & 0 & 0 \\ 0 & .5 & 0 & 0 & .5 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & .7 & 0 & .3 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & .5 & 0 & 0 & 0 & 0 & .5 \\ 0 & 0 & 0 & 0 & 0 & .25 & .25 & 0 & .5 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix} \end{matrix}$$

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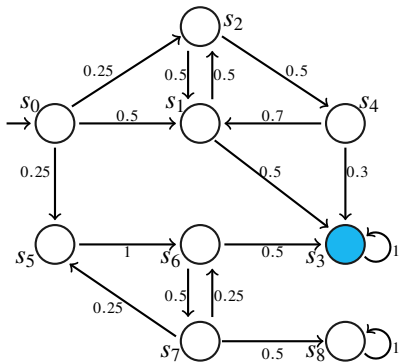
	s_0	s_1	s_2	s_3	s_4	s_5	s_6	s_7	s_8
s_0	0	.5	.25	0	0	.25	0	0	0
s_1	0	0	.5	.5	0	0	0	0	0
s_2	0	.5	0	0	.5	0	0	0	0
s_3	0	0	0	1	0	0	0	0	0
s_4	0	.7	0	.3	0	0	0	0	0
s_5	0	0	0	0	0	0	1	0	0
s_6	0	0	0	.5	0	0	0	.5	0
s_7	0	0	0	0	0	.25	.25	0	.5
s_8	0	0	0	0	0	0	0	0	1

blue state: only one φ -state.

- ▶ initial distribution over 9 states is $\mathbf{d}_0 = (1, 0, 0, 0, 0, 0, 0, 0, 0) = [1\ 0\ 0\ 0\ 0\ 0\ 0\ 0\ 0]$.
- ▶ distribution after 1 transition, 2 transitions, and 3 transitions, respectively:
 $\mathbf{d}_1 = \mathbf{d}_0 \cdot A = (0, .5, .25, \mathbf{0}, 0, .25, 0, 0, 0)$
 $\mathbf{d}_2 = \mathbf{d}_0 \cdot A^2 = (0, .125, .25, \mathbf{.25}, .125, 0, .25, 0, 0)$
 $\mathbf{d}_3 = \mathbf{d}_0 \cdot A^3 = (0, .2125, .0625, \mathbf{.475}, .125, 0, 0, .125, 0)$

Another example: How to handle reachability properties

Wanted: counterexamples for $\mathbb{P}_{\leq 0.4}(\diamond \varphi)$, or, in shorthand, $(\diamond^{\leq 0.4} \varphi)$.



Approach 1, based on using the transition (right-stochastic) 9×9 matrix A :

	s_0	s_1	s_2	s_3	s_4	s_5	s_6	s_7	s_8
s_0	0	.5	.25	0	0	.25	0	0	0
s_1	0	0	.5	.5	0	0	0	0	0
s_2	0	.5	0	0	.5	0	0	0	0
s_3	0	0	0	1	0	0	0	0	0
s_4	0	.7	0	.3	0	0	0	0	0
s_5	0	0	0	0	0	0	1	0	0
s_6	0	0	0	.5	0	0	0	.5	0
s_7	0	0	0	0	0	.25	.25	0	.5
s_8	0	0	0	0	0	0	0	0	1

blue state: only one φ -state.

- ▶ **Conclusion:** Starting from s_0 , state s_3 is reached with probability $.475 > .4$ after 3 transitions.
- ▶ Hence, there is a counterexample C for $s_0 \not\models (\diamond^{\leq 0.4} \varphi)$ consisting of finite paths, each with at most 3 transitions – but we have not determined the members of the counterexample C yet, nor do we know if it is **minimal** or **smallest** (cf. page 8)

Another example: How to handle reachability properties

Wanted: counterexamples for $\mathbb{P}_{\leq 0.4}(\diamond \varphi)$, or, in shorthand, $(\diamond^{\leq 0.4} \varphi)$.

- ▶ **Approach 2:** Let S be the set of states in the Markov chain, $s_0 \in S$ a single initial state, and $\text{Target} \subseteq S$ a non-empty set of target states.

For every state s , we define the probability p_s of reaching the states in Target from s :

$$p_s \triangleq \begin{cases} 1 & \text{if } s \in \text{Target}, \\ 0 & \text{if no state in Target is reachable from } s, \\ \sum_{s' \in S} \text{Pr}(s, s') \cdot p_{s'} & \text{otherwise.} \end{cases}$$

- ▶ This defines a system of linear equations over the variables $V \triangleq \{p_s \mid s \in S\}$ whose unique solution $\sigma : V \rightarrow [0, 1]$ assigns to each p_s the probability of reaching Target from s .
- ▶ Hence, $\mathcal{M} \models \mathbb{P}_{\leq \rho}(\diamond \text{target})$ iff $\sigma(p_{s_0}) \leq \rho$, where “target” is an atomic proposition which labels every state in Target .
- ▶ **Advantage of Approach 2 over Approach 1:** Solving a system of linear equations instead of repeatedly multiplying stochastic matrices.

Another example: How to handle reachability properties

Wanted: counterexamples for $\mathbb{P}_{\leq 0.4}(\diamond \varphi)$, or, in shorthand, $(\diamond^{\leq 0.4} \varphi)$.

- ▶ For the Markov chain \mathcal{M} shown on slide 21, we obtain:

$$\begin{aligned} p_{s_0} &= 0.5 p_{s_1} + 0.25 p_{s_2} + 0.25 p_{s_5} & p_{s_1} &= 0.5 p_{s_2} + 0.5 p_{s_3} \\ p_{s_2} &= 0.5 p_{s_1} + 0.5 p_{s_4} & p_{s_3} &= 1 \\ p_{s_4} &= 0.7 p_{s_1} + 0.3 p_{s_3} & p_{s_5} &= 1 p_{s_6} \\ p_{s_6} &= 0.5 p_{s_3} + 0.5 p_{s_7} & p_{s_7} &= 0.25 p_{s_5} + 0.25 p_{s_6} \end{aligned}$$

We can remove all states from \mathcal{M} which do not reach states in Target. In this example, we remove s_8 , thus also removing equation $p_{s_8} = 0$.

- ▶ Solving the system of linear equations (by hand or by using Matlab or Octave), we obtain a solution $\sigma : \{p_{s_0}, p_{s_1}, \dots, p_{s_7}\} \rightarrow [0, 1]$ such that:

$$\begin{aligned} \sigma(p_{s_0}) &= 11/12 & \sigma(p_{s_1}) &= \sigma(p_{s_2}) = \sigma(p_{s_3}) = \sigma(p_{s_4}) = 1 \\ \sigma(p_{s_5}) &= \sigma(p_{s_6}) = 2/3 & \sigma(p_{s_7}) &= 1/3 \end{aligned}$$

- ▶ **Conclusion:** Starting from s_0 , state s_3 is reached with probability $\frac{11}{12} > .4$. Hence, there is a counterexample C for $s_0 \not\models (\diamond^{\leq .4} \varphi)$, though we do not know the members of C yet!!

Another example: How to handle reachability properties

Wanted: counterexamples for $\mathbb{P}_{\leq 0.4}(\diamond \varphi)$, or, in shorthand, $(\diamond^{\leq 0.4} \varphi)$.

- ▶ **Approach 3**, most efficient and most direct, repeats the steps carried out to find counterexamples for $s_0 \not\models (\varphi \uplus^{\leq 1/2} \psi)$, from slide 12 to slide 20.
- ▶ We obtain, in order of decreasing probabilities:

evidence	weight (rounded)	probability
$\pi_1 \triangleq s_0 s_1 s_3$	1.39	0.25
$\pi_2 \triangleq s_0 s_5 s_6 s_3$	2.08	0.125
$\pi_3 \triangleq s_0 s_2 s_1 s_3$	2.77	0.0625
$\pi_4 \triangleq s_0 s_1 s_2 s_1 s_3$	2.77	0.0625
$\pi_5 \triangleq s_0 s_2 s_4 s_1 s_3$	3.13	0.04375
$\pi_6 \triangleq s_0 s_1 s_2 s_4 s_1 s_3$	3.13	0.04375
$\pi_7 \triangleq s_0 s_2 s_4 s_3$	3.28	0.03750
$\pi_8 \triangleq s_0 s_1 s_2 s_4 s_3$	3.28	0.03750
...

where we take weight $w(s, s') \triangleq -\ln(\Pr(s, s'))$ for all states $s, s' \in S$.

- ▶ $\sum_{i \in \{1,2,3\}} \Pr(\pi_i) = \sum_{i \in \{1,2,4\}} \Pr(\pi_i) = 0.4375 > 0.4$

(but why not $\{\pi_1, \pi_2, s_0 s_2 s_4\}$ or $\{\pi_1, \pi_2, s_0 s_1 s_2 s_4\}$???)

implies both $\{\pi_1, \pi_2, \pi_3\}$ and $\{\pi_1, \pi_2, \pi_4\}$ are smallest counterexamples.

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