# Notes for Lecture 8

## 1 Chinese Remainder Theorem

Let $p \neq q$ be two primes. The Chinese Remainder Theorem (CRT) says that working modulo $n = pq$ is essentially the same as working modulo $p$ and modulo $q$ at the same time: more formally (for those comfortable with abstract algebra), that the ring $\mathbb{Z}_n$ is isomorphic to the product ring $\mathbb{Z}_p \times \mathbb{Z}_q$. (Actually, this is the "light" version of CRT, which is all we need for this course. The full-fledged version says that working modulo $a_1 a_2 \dots a_k$, where $a_i$ are pairwise relatively prime, is the same as working simultaneously modulo $a_1, a_2, \dots, a_k$.)

Here is an example. Consider all the values modulo 35. They are in one-to-one correspondence with values modulo 5 and modulo 7.

| | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 15 | 30 | 10 | 25 | 5 | 20 |
| 1 | 21 | 1 | 16 | 31 | 11 | 26 | 6 |
| 2 | 7 | 22 | 2 | 17 | 32 | 12 | 27 |
| 3 | 28 | 8 | 23 | 3 | 18 | 33 | 13 |
| 4 | 14 | 29 | 9 | 24 | 4 | 19 | 34 |

Observe that if you want to add, say, 17 and 29 (underlined in the table), is the same as adding 3 (which is 17 mod 7) and 1 (which is 29 mod 7) modulo 7 to get 4; adding 2 (which is 17 mod 5) and 4 (which is 29 mod 5) modulo 5 to get 1; and then looking up the value corresponding to coordinates 4 and 1 in the table to get 11 (in a box in the table). Thus, we can do addition coordinatewise. Same for multiplication.

We now formally state and prove the observations above, generalized to $p$ and $q$ instead of 5 and 7.

**Theorem 1.** *Let $p \neq q$ be primes, $n = pq$. For each $a \in \mathbb{Z}_p$, $b \in \mathbb{Z}_q$, there is unique $c$, $0 \leq c < n$ such that $c \equiv a \pmod{p}$ and $c \equiv b \pmod{q}$.*

*Proof.* Let $r = p^{-1} \bmod q$ and $s = q^{-1} \bmod p$. Let $c' = rpb + sqa$. Then $c' \equiv rpb + sqa \equiv r \cdot 0 \cdot b + 1 \cdot a \equiv a \pmod{p}$, and $c' \equiv rpb + sqa \equiv 1 \cdot b + s \cdot 0 \cdot a \equiv b \pmod{q}$. Let $c = c' \bmod pq$. Then $pq|(c - c')$, so $p|(c - c')$, so $c \equiv c' \pmod{p}$. Similarly, $c \equiv c' \pmod{q}$. Hence, $c$ satisfies all the conditions: $0 \leq c < n$, and $c \equiv a \pmod{p}$ (because $c \equiv c' \equiv a \pmod{p}$), and $c \equiv b \pmod{q}$ (because $c \equiv c' \equiv b \pmod{q}$). Thus, for every pair $(a, b)$ there is a $c$. There are $pq = n$ possible pairs, and $n$ possible values of $c$, so for each pair there must be exactly one value of $c$, so it's unique for each $(a, b)$. □

Denote by $\mathrm{crt}(a, b)$ the unique value of $c$ given by the above theorem. Then $\mathrm{crt}(a, b) = c$ if an only if $(a, b) = (c \bmod p, c \bmod q)$. Let $c_1 = \mathrm{crt}(a_1, b_1)$, $c_2 = \mathrm{crt}(a_2, b_2)$, and $c_3 = c_1 + c_2 \bmod n$. Then $c_3 \bmod p = (c_1 + c_2) \bmod p = (a_1 + a_2) \bmod p$ (because $n$ divides $c_3 - c_1 - c_2$, and therefore so does $p$) and similarly $c_3 \bmod q = (b_1 + b_2) \bmod q$. Hence $c_3 = \mathrm{crt}(a_1 + a_2, b_1 + b_2)$. Same for multiplication. Thus, we can look at addition and multiplication modulo $n$ coordinatewise: modulo $p$ and modulo $q$.

We will denote by $\mathbb{Z}_n^*$ the set of values in $\mathbb{Z}_n$ that are relatively prime to $n$. Note that the "coordinates" of $\mathbb{Z}_n^*$ are in $\mathbb{Z}_p^*$ and $\mathbb{Z}_q^*$, and that $\mathbb{Z}_n^*$ has $(p-1)(q-1)$ elements.

Note that the above proof is constructive: that is, $c$ is efficiently (and, in fact, quite easily) computable given $a$ and $b$. Thus, it is often more efficient to work modulo $p$ and $q$ separately and the reconstruct the value modulo $n$ when it is needed.

## 2 Squares and Square Roots

Let $p > 2$ be a prime. Let $QR_p$ denote the set of squares in $\mathbb{Z}_p^*$. Recall from HW2 that for $a \in \mathbb{Z}_p^*$, if $a \in QR_p$, then $a^{(p-1)/2} \equiv 1$, and if $a \notin QR_p$, then $a^{(p-1)/2} \equiv -1$.

Suppose $p \equiv 3 \pmod 4$. Take $s \in \mathbb{Z}_p^*$. It has two roots: $r$ and $-r$. Exactly one of these two roots is itself in $QR_p$. Indeed, consider $r^{(p-1)/2}$ and $(-r)^{(p-1)/2}$. Since $(p-1)/2$ is odd (because $p = 4k+3$ for some $k$), $(-r)^{(p-1)/2} = -\left(r^{(p-1)/2)}\right)$, so one is 1 and the other is $-1$.

Hence, if we let $f_p(x) : QR_p \to QR_p$ be the map $x \mapsto x^2 \bmod p$, we see that for each $s \in QR_p$, there exists a unique inverse $r \in QR_p$ such that $f(r) = s$ (namely, $r$ is the square root of $s$ that is itself a square). So $f_p$ of $x$ is a permutation of $QR_p$. Note that $f_p$ is easy to compute (just squaring) and easy to invert (as shown on HW2, it's easy to compute square roots modulo $p$).

Now let $p \neq q$ be two distinct odd primes, and let $n = pq$. Let $QR_n$ denote the set of squares in $\mathbb{Z}_n^*$. Then if $s$ is a square modulo $n$, it is also a square modulo $p$ and $q$. Since it has two roots $\pm r_1$ modulo $p$ and two roots $\pm r_2$ modulo $q$, it has four roots modulo $n$: $\mathrm{crt}(\pm r_1, \pm r_2)$.

Suppose both $p$ and $q$ are congruent to 3 modulo 4. Then exactly one of $\pm r_1$ is a square modulo $p$, and exactly one of $\pm r_2$ is a square modulo $q$, so exactly one of $\mathrm{crt}(\pm r_1, \pm r_2)$ is a square modulo $n$. Hence, if we let $f_n(x) : QR_n \to QR_n$ be the map $x \mapsto x^2 \bmod n$, we see that $f_n(x)$ is a permutation over $QR_n$. Note that $f_n(x)$ is easy to compute. We will argue below that it is hard to invert—as hard as it is to factor $n$.

## 3 Blum-Blum-Shub Generator

The following construction is due to [BBS86][1]. Starting with a sufficiently long random seed, select two $k$-bit random primes $p, q$ that are 3 modulo 4, let $n = pq$, and let $x$ be random element of $QR_n$ (just select a random element of $\mathbb{Z}_n$, check if it's relatively prime with $n$, and square it). Let $x_1 = x, x_2 = f_n(x), x_3 = f_n(x_2), \ldots, x_l = f_n(x_{l-1})$. Output the least significant bit for each $x_i$.

Note that this looks very much like the Blum-Micali generator, with exponentiation mod $p$ replaced with squaring mod $n$, and $B$ replaced with least significant bit. The proof is very similar, too. We simply need three facts: that the function $f_n$ is a permutation (already shown above), that computing $x$ from $x^2 \bmod n$ is hard (discussed in the next section), and that computing the least significant bit of $x$ from $x^2 \bmod n$ is as hard as computing all of $x$ (shown in [ACGS88]; an alternative proof is given is in [AGS03]; we will not discuss either here). These three facts correspond, in the Blum-Micali case, to the fact that modular exponentiation is a permutation of $\mathbb{Z}_p^*$ (which is used in the reduction because we have to know that the permutation has a unique inverse in order to show that the bits the reduction feeds to the adversary correspond to bits a generator would have generated), to the assumption that discrete logarithm is hard, and the theorem that $B(x)$ is as hard as to compute from $g^x \bmod p$ as $x$ itself.

This generator is more efficient than Blum-Micali: requires only one modular squaring per bit, instead of one one modular exponentiation. It is also based on a different (depending on whom you ask, more or less plausible) assumption: that factoring $n$ is hard. We will show this in the next section.

## 4 Square Roots Modulo a Composite are as Hard as Factoring

We want to justify why we believe it's hard to compute $x$ from $x^2$ modulo $n$. Indeed, let $s = r^2 \bmod n$. Then $s$ has four square roots, as discussed above $\mathrm{crt}(r_1, r_2), \mathrm{crt}(-r_1, -r_2), \mathrm{crt}(r_1, -r_2), \mathrm{crt}(-r_1, r_2)$. Take two of these that are not negatives of each other, e.g., $r = \mathrm{crt}(r_1, r_2)$ and $r' = \mathrm{crt}(r_1, -r_2)$. Add them to get $r + r' = \mathrm{crt}(2r_1, 0)$. Thus, $r + r' \equiv 0 \pmod q$, so $q|(r + r')$. Note also that $r + r' \not\equiv 0 \pmod p$,

[1] Conference version published in Crypto in 1982.

so $p \not| (r + r')$. Hence, $\gcd(r + r', n) = q$. Thus, if you know two such roots, you can factor $n$, by simply computing the greatest common divisor (this can be done quickly with Euclid's algorithm).

Now suppose we have an algorithm $A$ that computes square roots modulo $n$. We will use it to factor $n$ as follows: take a random $r \in \mathbb{Z}_n^*$, compute $s = r^2 \bmod n$, and give $s$ to $A$. $A$ will return some root $r'$ of $s$. Because $s$ has four roots and $r$ was chosen at random (and not given to $A$), no matter how $A$ works, $\Pr[r = \pm r'] = 1/2$. Hence, in half the cases, $\gcd(r + r', n)$ will give you a factor $p$ or $q$ of $n$.

Thus, we just proved (by contradiction and reduction, as usual) that if factoring $n$ is hard, so is computing square roots modulo $n$. Hence, the Blum-Blum-Shub generator is secure based on the following assumption:

**Assumption 1.** For any poly-time algorithm $F$, there exists a negligible function $\eta$ such that, if you generate random $k$-bit primes $p$ and $q$ that are both 3 modulo 4, and let $n = pq$, $\Pr[F(n) = p] \leq \eta(k)$.

## References

[ACGS88] W. Alexi, B. Chor, O. Goldreich, and C. Schnorr. RSA and Rabin functions: Certain parts are as hard as the whole. *SIAM Journal on Computing*, 17(2):194–209, April 1988.

[AGS03] Adi Akavia, Shafi Goldwasser, and Muli Safra. Proving hardcore predicates using list decoding. In *44th Annual Symposium on Foundations of Computer Science*, Cambridge, Massachusetts, October 2003. IEEE.

[BBS86] L. Blum, M. Blum, and M. Shub. A simple unpredictable pseudo-random number generator. *SIAM Journal on Computing*, 15(2):364–383, May 1986.